(Note: The final is a 24 hour take home exam. We can pick it up at any point starting the Saturday after classes end and ending the Sunday after the Sunday after classes end. However, Professor Goldmakher strongly encourages not taking the exam the first Saturday or Sunday, as he's planning a review session for Monday is morning.)

Let's talk about the sudoku rule from last time!

Sudoky Rule: Given p a prime, fix any a £ \(\frac{2}{1},2,\theta, p-1\)\). The at row of the multiplication table (mod p) consists of all the numbers in \(\frac{1}{2},-\theta, p-1\)\)\, each appearing exactly once. (The rule also and covers columns, but since multiplication is commutative, if \(\text{commutative}\) it's true for the rows, it's true

for the columns.)

That's a relatively long definition of the rule. Let's tighten it up by writing it with mathematical symbols. But how? We need a way to say that any two elements of the ath row are distinct. Note that every element of the ath row is a times some number between 1 and p-1, mod p. in If we want to say two elements are unique, then, we get:

We don't even need to restrict k and l to being in this set! k and l can be any integers—They just have to be distinct (mod p). Thus, we can condense this further:

k 丰 l (mod p) → ak ‡ al (mod p).

If this is true, every element in the ath row must be distinct. However, there are p-1 possible elements and p-1 spaces, so each is used exactly once.

	This means:
	k ≠ l (mod p) = ak ≠ al (mod p)
	<b>\$</b>
	∀k, l∈ {1,2,, p-13 5.+. k+l, ak (mod p) ≠ al (mod p)
	The Sudoku Rule.
	That one line is equivalent to the sudoku rule, and it's much
	faster to write! Note that that statement is equivalent to
	its contrapositive:
	$k \neq l \pmod{p} \Rightarrow ak \neq al \pmod{p} \iff Given a \neq 0 \pmod{p}$ , if $ak \equiv al \pmod{p}$ ,
	then $k \equiv l \pmod{p}$ .
()	To prove this, it'd be rather nice to just say ak=al(mod P)=ak-al=0(mod P)=a(k-l)=0(mod P)
	ak=bl(mod P) $\Rightarrow$ ak-al=0(mod P) $\Rightarrow$ a(k-l)=0(mod P) $\Rightarrow$ a=0 (mod p) or k-l=0 (mod P). How can we show show that last step?
	we show that last step?
	All we need to do to prove that workers contrapositive
	is show the last step in the wishful thinking cloud,
Proposition:	given $a \neq 0 \pmod{p}$ , if $ak \equiv al \pmod{p}$ , then $k \equiv l \pmod{p}$ $\iff$
Proposition:	If $wq \equiv 0 \pmod{p}$ , then $w \equiv 0 \pmod{p}$ or $q \equiv 0 \pmod{p}$ .
	Exercise for you: why are the previous two propositions equivalent?
	Next, we observe that the latter proposition is equivalent to the following result
	about prime numbers:
Proposition:	plug ⇒ plw or plg.
	Let's prove this! (The proof is wild!)
	Proof: Suppose plug, but prw. Then gld(pw)=1(since
	$\rho$ only has 2 factors). By Bézont's theorem (which is on the book homework, this means $\exists x,y \in \mathbb{Z} \text{ s.t. } \rho x + wy = 1$ .
	Bone homework, this means Ix, y EZ s.t. px+ wy = 1.

Multiply both sides by q: we get
This means pxq+ wqy=q. plpxq and plwqy (since
plug), so pl(pxq+wqy). Thus, plq.
Hooray! The Sudoka rule is true!
Awesome consequence of this:
What's the product of all the elements appearing in the ath
on ore hand, it's (p-1)! "Since the Endoku rule is the.
On the other hand, by definition, it's a (2a)(3a)(p-1)a (modp).
Thus, (p-1)! = a (2a)(3a) (p-1)a man = a [p-1)! (mod p) =
$\exists 1 \equiv a \pmod{p}$ .
We've proved Fermat's Little Theorem: $a^{p-1} \equiv 1 \pmod{p} \ \forall a \equiv 0 \pmod{p}.$
$a^{p-1} \equiv 1 \pmod{p}  \forall \alpha \equiv 0 \pmod{p}$ .