

THE ARITHMETIC AND GEOMETRIC MEAN INEQUALITY

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ABSTRACT. We provide sketches of proofs of the Arithmetic Mean - Geometric Mean Inequality. These notes are based on discussions with Vitaly Bergelson, Eitan Sayag, and the students of Math 487 (Ohio State, Autumn 2003).

1. INTRODUCTION

Definition 1.1 (Arithmetic Mean). *The Arithmetic Mean of a_1, \dots, a_n is*

$$AM(a_1, \dots, a_n) = \frac{a_1 + \dots + a_n}{n}. \quad (1.1)$$

Definition 1.2 (Geometric Mean). *The Geometric Mean of a_1, \dots, a_n is*

$$GM(a_1, \dots, a_n) = \sqrt[n]{a_1 \cdots a_n}. \quad (1.2)$$

Theorem 1.3 (Arithmetic Mean - Geometric Mean). *Let a_1, \dots, a_n be n positive numbers. Then*

$$AM(a_1, \dots, a_n) \geq GM(a_1, \dots, a_n). \quad (1.3)$$

Remark 1.4. *Note the above trivially holds if $a_1 = \dots = a_n$; in fact, equality holds if and only if all a_i are equal.*

Remark 1.5. *Note that if the Arithmetic Mean - Geometric Mean inequality holds for a_1, \dots, a_n , it holds for $\alpha a_1, \dots, \alpha a_n$ for any $\alpha > 0$. Thus, we can rescale the sum $a_1 + \dots + a_n$ (assuming it is non-zero) to be whatever we want.*

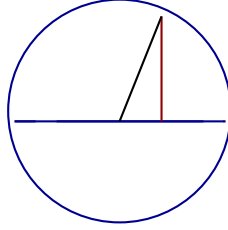
Remark 1.6. *Note the $n = 2$ case follows immediately from $(\sqrt{a_1} - \sqrt{a_2})^2 \geq 0$.*

2. GEOMETRIC PROOF WHEN $n = 2$

Without loss of generality, assume $a_1 > a_2$. Construct a circle with diameter $a_1 + a_2$, hence radius $\frac{a_1 + a_2}{2}$. On the main diagonal, a_1 units from one end (a_2 from the other), draw the perpendicular bisector to the main diagonal, which hits the circle at some point, say P .

The length of the perpendicular bisector from the main diagonal to the circle is $\sqrt{a_1 a_2}$ – this can be shown by applications of the Pythagorean Theorem.

Form a triangle using this as one side, and with hypotenuse from the center of the circle to P . The hypotenuse will have length $\frac{a_1+a_2}{2}$, which must therefore be larger than the side $\sqrt{a_1a_2}$.



Blue Line: Diameter of the Circle, Length $a_1 + a_2$
 Black Line: Radius (Arithmetic Mean), Length $\frac{a_1+a_2}{2}$
 Red Line: Altitude (Geometric Mean), Length $\sqrt{a_1a_2}$

3. MULTIVARIABLE CALCULUS PROOF

Use Lagrange Multipliers, with

$$\begin{aligned} f(a_1, \dots, a_n) &= (a_1 \cdots a_n)^{\frac{1}{n}} \\ g(a_1, \dots, a_n) &= \frac{a_1 + \cdots + a_n}{n} - c. \end{aligned} \quad (3.4)$$

To find if the Lagrange Multipliers give a maximum or minimum, check at $1, 1, 1, \dots, 1, 2^n$.

4. STANDARD INDUCTION PROOF

We proceed by induction, the $n = 1$ and $n = 2$ cases already handled above.

We must show

$$\frac{a_1 + \cdots + a_n}{n} \geq \sqrt[n]{a_1 \cdots a_n}. \quad (4.5)$$

Without loss of generality, we may rescale the a_i so that $a_1 \cdots a_n = 1$. If all $a_i = 1$, the proof is trivial. Thus, assume at least one $a_i > 1$ and one $a_i < 1$; we assume $a_1 > 1, a_2 < 1$.

Thus, by the inductive assumption, we have

$$\frac{a_1a_2 + a_3 + \cdots + a_n}{n-1} \geq \sqrt[n-1]{(a_1a_2)a_3 \cdots a_n} = 1. \quad (4.6)$$

Thus, we have

$$a_1a_2 + a_3 + \cdots + a_n \geq n-1. \quad (4.7)$$

We need to show

$$a_1 + a_2 + \cdots + a_n \geq n. \quad (4.8)$$

This would follow if $a_1 + a_2 - (a_1a_2 + 1) \geq 0$. But

$$a_1 + a_2 - (a_1a_2 + 1) = (a_1 - 1)(1 - a_2) \geq 0, \quad (4.9)$$

proving the claim. \square

5. INDUCTION BY POWERS OF 2

We first show if the Arithmetic Mean - Geometric Mean Inequality holds for $n = 2^{k-1}$, then it holds for $n = 2^k$. We then show how to handle n that are not powers of 2.

Lemma 5.1. *If the AM - GM Inequality holds for $n = 2^{k-1}$, it holds for $n = 2^k$.*

Proof. We assume the case $n = 2$ has already been done, and is available for use below. namely for any $c_1, c_2 > 0$,

$$\frac{c_1 + c_2}{2} \geq \sqrt{c_1 c_2} = 1. \quad (5.10)$$

Without loss of generality, rescale so that $a_1 \cdots a_{2^k} = 1$. Let $b_1 = a_1 + \cdots + a_{2^{k-1}}$ and $b_2 = a_{2^{k-1}+1} + \cdots + a_{2^k}$.

By induction, we can apply the AM-GM to b_1 and b_2 and we find

$$\frac{b_1}{2^{k-1}} = \frac{a_1 + \cdots + a_{2^{k-1}}}{2^{k-1}} \geq \sqrt[2^{k-1}]{a_1 \cdots a_{2^{k-1}}} \quad (5.11)$$

and

$$\frac{b_2}{2^{k-1}} = \frac{a_{2^{k-1}+1} + \cdots + a_{2^k}}{2^{k-1}} \geq \sqrt[2^{k-1}]{a_{2^{k-1}+1} \cdots a_{2^k}}. \quad (5.12)$$

Combining yields

$$\frac{b_1 + b_2}{2^k} = \frac{a_1 + \cdots + a_{2^k}}{2^k} \geq \frac{\sqrt[2^{k-1}]{a_1 \cdots a_{2^{k-1}}} + \sqrt[2^{k-1}]{a_{2^{k-1}+1} \cdots a_{2^k}}}{2}. \quad (5.13)$$

Applying the $n = 2$ case to the right hand side yields

$$\frac{a_1 + \cdots + a_{2^k}}{2^k} \geq \sqrt{\sqrt[2^{k-1}]{a_1 \cdots a_{2^{k-1}}} \cdot \sqrt[2^{k-1}]{a_{2^{k-1}+1} \cdots a_{2^k}}} = 1, \quad (5.14)$$

as the right hand side is now just $\sqrt[2^k]{a_1 \cdots a_n} = 1$. \square

We now prove the AM - GM Inequality for any n . Choose k so that $2^{k-1} < n < 2^k$. Then we need to add $2^k - n$ terms to have a power of 2. As always, we may assume $a_1 \cdots a_n = 1$.

Let $A_n = \frac{a_1 + \cdots + a_n}{n}$. Consider the sequence $a_1, \dots, a_n, A_n, \dots, A_n$, where we have A_n a total of $2^k - n$ times.

Exercise 5.2. *Show the Arithmetic Mean of these 2^k numbers is still A_n .*

Exercise 5.3. *Show the Geometric Mean of these 2^k numbers is*

$$\sqrt[2^k]{a_1 \cdots a_n A_n^{2^k - n}} = \left(\frac{a_1 + \cdots + a_n}{n} \right)^{1 - \frac{n}{2^k}}; \quad (5.15)$$

remember we have rescaled so that $a_1 \cdots a_n = 1$.

As we have 2^k numbers, we may apply the AM - GM Inequality, and we obtain

$$\begin{aligned} \frac{a_1 + \cdots + a_n + (2^k - n) \cdot \frac{a_1 + \cdots + a_n}{n}}{2^k} &\geq \sqrt[2^k]{a_1 \cdots a_n \left(\frac{a_1 + \cdots + a_n}{n} \right)^{2^k - n}} \\ \frac{a_1 + \cdots + a_n}{n} &\geq \left(\frac{a_1 + \cdots + a_n}{n} \right)^{1 - \frac{n}{2^k}} \\ \left(\frac{a_1 + \cdots + a_n}{n} \right)^{\frac{2^k}{n}} &\geq 1 \\ \frac{a_1 + \cdots + a_n}{n} &\geq 1, \end{aligned} \tag{5.16}$$

as claimed, completing the proof.

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