# Benford's law, or: Why the IRS should care about number theory!

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Central Connecticut State University September 25th, 2009

# **Summary**

- Review Benford's Law.
- Discuss examples and applications.
- Sketch proofs.
- Describe open problems.

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  - $\diamond$  Example:  $\pi$  mod 1 is about .14159.

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  - ♦ Long street [1, L]: L = 199 versus L = 999.
  - ♦ Oscillates between 1/9 and 5/9 with first digit 1.
  - Many streets of different sizes: close to Benford.

# **Examples**

- recurrence relations
- special functions (such as n!)
- iterates of power, exponential, rational maps
- products of random variables
- L-functions, characteristic polynomials
- iterates of the 3x + 1 map
- differences of order statistics
- hydrology and financial data
- many hierarchical Bayesian models

# **Applications**

- analyzing round-off errors
- determining the optimal way to store numbers

detecting tax and image fraud, and data integrity

# **General Theory**

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Key observation:  $\log_{10}(x) = \log_{10}(\tilde{x}) \mod 1$  if and only if x and  $\tilde{x}$  have the same leading digits. Thus often study  $y = \log_{10} x$ .

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# **Equidistribution**

 $\{y_n\}_{n=1}^{\infty}$  is equidistributed modulo 1 if probability  $y_n \mod 1 \in [a, b]$  tends to b - a:

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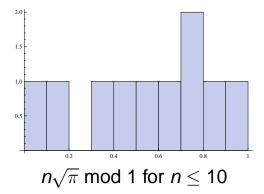
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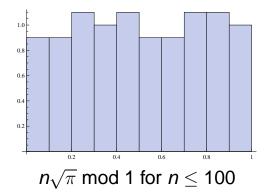
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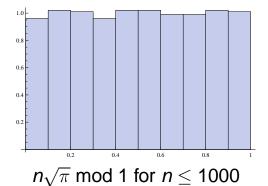
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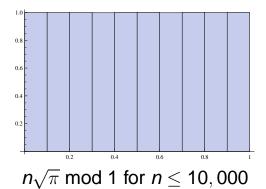
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- Examples:  $\log_{10} 2$ ,  $\log_{10} \left( \frac{1+\sqrt{5}}{2} \right) \notin \mathbb{Q}$ . *Proof:* if rational:  $2 = 10^{p/q}$ Thus  $2^q = 10^p$  or  $2^{q-p} = 5^p$ , impossible.









#### **Denseness**

# **Dense**

A sequence  $\{z_n\}_{n=1}^{\infty}$  of numbers in [0,1] is dense if for any interval [a,b] there are infinitely many  $z_n$  in [a,b].

- Dirichlet's Box (or Pigeonhole) Principle:
   If n + 1 objects are placed in n boxes, at least one box has two objects.
- Denseness of  $n\alpha$ : Thm: If  $\alpha \notin \mathbb{Q}$  then  $z_n = n\alpha \mod 1$  is dense.

# **Proof** $n\alpha \mod 1$ dense if $\alpha \notin \mathbb{Q}$

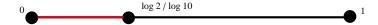
- Enough to show in [0, b] infinitely often for any b.
- Choose any integer Q > 1/b.
- Q bins:  $\left[0, \frac{1}{Q}\right], \left[\frac{1}{Q}, \frac{2}{Q}\right], \ldots, \left[\frac{Q-1}{Q}, Q\right].$
- Q + 1 objects:  $\{\alpha \mod 1, 2\alpha \mod 1, \dots, (Q+1)\alpha \mod 1\}.$
- Two in same bin, say  $q_1\alpha \mod 1$  and  $q_2\alpha \mod 1$ .
- Exists integer p with  $0 < q_2 \alpha q_1 \alpha p < \frac{1}{Q}$ .
- Get  $(q_2 q_1)\alpha \mod 1 \in [0, b]$ .

# **Fundamental Equivalence**

Data set  $\{x_i\}$  is Benford base B if  $\{y_i\}$  is equidistributed mod 1, where  $y_i = \log_B x_i$ .

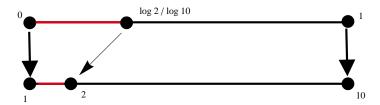
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# **Proof:**

- $x = M_B(x) \cdot B^k$  for some  $k \in \mathbb{Z}$ .
- $FD_B(x) = d \text{ iff } d \le M_B(x) < d + 1.$
- $\log_B d \le y < \log_B (d+1)$ ,  $y = \log_B x \mod 1$ .
- If  $Y \sim \text{Unif}(0, 1)$  then above probability is  $\log_{\mathbb{R}}(\frac{d+1}{d})$ .

•  $2^n$  is Benford base 10 as  $\log_{10} 2 \notin \mathbb{Q}$ .

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 Guess  $a_n = n^r$ :  $r^{n+1} = r^n + r^{n-1}$  or  $r^2 = r+1.$  Roots  $r = (1 \pm \sqrt{5})/2.$  General solution:  $a_n = c_1 r_1^n + c_2 r_2^n.$  Binet:  $a_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2}\right)^n.$ 

Fibonacci numbers are Benford base 10.

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$$\diamond a_{n+1} = 2a_n - a_{n-1}$$

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$$\diamond$$
 take  $a_0 = a_1 = 1$  or  $a_0 = 0$ ,  $a_1 = 1$ .

# Digits of $2^n$

First 60 values of 2<sup>n</sup> (only displaying 30)

1	1024	1048576	digit	#	Obs Prob	Benf Prob
2	2048	2097152	1	18	.300	.301
4	4096	4194304	2	12	.200	.176
8	8192	8388608	3	6	.100	.125
16	16384	16777216	4	6	.100	.097
32	32768	33554432	5	6	.100	.079
64	65536	67108864	6	4	.067	.067
128	131072	134217728	7	2	.033	.058
256	262144	268435456	8	5	.083	.051
512	524288	536870912	9	1	.017	.046

# **Data Analysis**

Introduction

- $\chi^2$ -Tests: Test if theory describes data
  - ⋄ Expected probability:  $p_d = \log_{10} \left( \frac{d+1}{d} \right)$ .
  - $\diamond$  Expect about  $Np_d$  will have first digit d.
  - $\diamond$  Observe Obs(d) with first digit d.

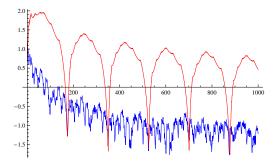
$$\phi \chi^2 = \sum_{d=1}^9 \frac{(\mathrm{Obs}(d) - Np_d)^2}{Np_d}$$
.

- $\diamond$  Smaller  $\chi^2$ , more likely correct model.
- Will study  $\gamma^n$ ,  $e^n$ ,  $\pi^n$ .

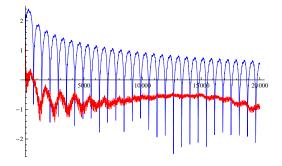
 $\chi^2$  values for  $\alpha^n$ ,  $1 \le n \le N$  (5% 15.5).

N	$\chi^2(\gamma)$	$\chi^2(e)$	$\chi^2(\pi)$
100	0.72	0.30	46.65
200	0.24	0.30	8.58
400	0.14	0.10	10.55
500	0.08	0.07	2.69
700	0.19	0.04	0.05
800	0.04	0.03	6.19
900	0.09	0.09	1.71
1000	0.02	0.06	2.90

 $\log(\chi^2)$  vs N for  $\pi^n$  (red) and  $e^n$  (blue),  $n \in \{1, ..., N\}$ . Note  $\pi^{175} \approx 1.0028 \cdot 10^{87}$ , (5%,  $\log(\chi^2) \approx 2.74$ ).



 $\log(\chi^2)$  vs N for  $\pi^n$  (red) and  $e^n$  (blue),  $n \in \{1, ..., N\}$ . Note  $e^3 \approx 20.0855$ , (5%,  $\log(\chi^2) \approx 2.74$ ).

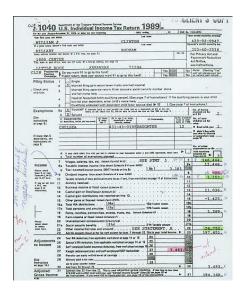


# **Applications**

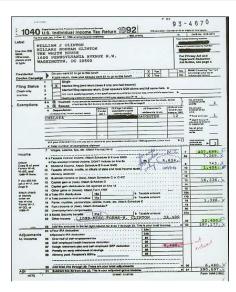
### **Stock Market**

Milestone	Date	Effective Rate from last milestone
108.35	Jan 12, 1906	
500.24	Mar 12, 1956	3.0%
1003.16	Nov 14, 1972	4.2%
2002.25	Jan 8, 1987	4.9%
3004.46	Apr 17, 1991	9.5%
4003.33	Feb 23, 1995	7.4%
5023.55	Nov 21, 1995	30.6%
6010.00	Oct 14, 1996	20.0%
7022.44	Feb 13, 1997	46.6%
8038.88	Jul 16, 1997	32.3%
9033.23	Apr 6, 1998	16.1%
10006.78	Mar 29, 1999	10.5%
11209.84	Jul 16, 1999	38.0%
12011.73	Oct 19, 2006	1.0%
13089.89	Apr 25, 2007	16.7%
14000.41	Jul 19, 2007	28.9%

# **Applications for the IRS: Detecting Fraud**



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#### Exhibit 3: Check Fraud in Arizona

The table lists the checks that a manager in the office of the Arizona State Treasurer wrote to divert funds for his own use. The vendors to whom the checks were issued were fictitious.

Date of Check	Amount
October 9, 1992	\$ 1,927.48
+	27,902.31
October 14, 1992	86,241.90
	72,117.46
	81,321.75
+	97,473.96
October 19, 1992	93,249.11
	89,658.17
	87,776.89
	92,105.83
ACCEPTAGE OF THE PROPERTY OF THE PERSON OF T	79,949.16
	87,602.93
	96,879.27
	91,806.47
	84,991.67
	90,831.83
	93,766.67
	88,338.72
	94,639.49
STREET,	83,709.28
	96,412.21
	88,432.86
<b>*</b>	71,552.16
TOTAL	\$ 1,878,687.58

# **Applications for the IRS: Detecting Fraud (cont)**

- Embezzler started small and then increased dollar amounts.
- Most amounts below \$100,000 (critical threshold for data requiring additional scrutiny).
- Over 90% had first digit of 7, 8 or 9.

# **Detecting Fraud**

#### **Bank Fraud**

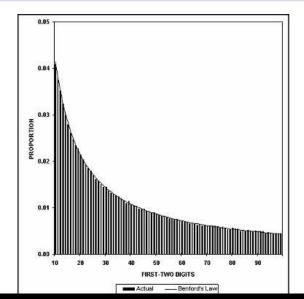
- Audit of a bank revealed huge spike of numbers starting with 48 and 49, most due to one person.
- Write-off limit of \$5,000. Officer had friends applying for credit cards, ran up balances just under \$5,000 then he would write the debts off.

# **Detecting Fraud**

#### **Enron**

- Benford's Law detected manipulation of revenue numbers.
- Results showed a tendency towards round Earnings Per Share (0.10, 0.20, etc.).
   Consistent with a small but noticeable increase in earnings management in 2002.

# Data Integrity: Stream Flow Statistics: 130 years, 457,440 records



#### **Election Fraud: Iran 2009**

Numerous protests and complaints over Iran's 2009 elections.

Lot of analysis done; data is moderately suspicious. Tests done include

- First and second leading digits;
- Last two digits (should almost be uniform);
- Last two digits differing by at least 2.

Warning: do enough tests, even if nothing is wrong will find a suspicious result, but when all tests are on the boundary....

### **Benford Good Processes**

#### Poisson Summation and Benford's Law: Definitions

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Poisson Summation Formula: f nice:

$$\sum_{\ell=-\infty}^{\infty} f(\ell) = \sum_{\ell=-\infty}^{\infty} \widehat{f}(\ell),$$

Fourier transform 
$$\hat{f}(\xi) = \int_{-\infty}^{\infty} f(x) e^{-2\pi i x \xi} dx$$
.

 $X_T$  is Benford Good if there is a nice f st

$$\mathrm{CDF}_{\overrightarrow{Y}_{T,B}}(y) = \int_{-\infty}^{y} \frac{1}{T} f\left(\frac{t}{T}\right) dt + E_{T}(y) := G_{T}(y)$$

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$$G_T(\infty) - G_T(Th(T)) = o(1)$$
,  $G_T(-Th(T)) - G_T(-\infty) = o(1)$ .

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- Small translated error:  $\mathcal{E}(a, b, T) = \sum_{|\ell| \leq Th(T)} [E_T(b + \ell) E_T(a + \ell)] = o(1).$

#### **Main Theorem**

# Theorem (Kontorovich and M-, 2005)

 $X_T$  converging to X as  $T \to \infty$  (think spreading Gaussian). If  $X_T$  is Benford good, then X is Benford.

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 $X_T$  converging to X as  $T \to \infty$  (think spreading Gaussian). If  $X_T$  is Benford good, then X is Benford.

- Examples
  - ♦ L-functions
  - characteristic polynomials (RMT)
  - $\diamond$  3x + 1 problem
  - geometric Brownian motion.

## Sketch of the proof

- Structure Theorem:
  - main term is something nice spreading out
  - ⋄ apply Poisson summation

# Sketch of the proof

- Structure Theorem:
  - main term is something nice spreading out
  - apply Poisson summation
- Control translated errors:
  - ♦ hardest step
  - techniques problem specific

$$\sum_{\ell=-\infty}^{\infty} \mathbb{P}\left(\boldsymbol{a} + \ell \leq \overrightarrow{\mathbf{Y}}_{T,B} \leq \boldsymbol{b} + \ell\right)$$

$$\sum_{\ell=-\infty}^{\infty} \mathbb{P}\left(\mathbf{a}+\ell \leq \overrightarrow{\mathbf{Y}}_{T,B} \leq \mathbf{b}+\ell
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ight] + \mathrm{o}(1)$$

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$$= \widehat{f}(0) \cdot (b-a) + \sum_{\ell \neq 0} \widehat{f}(T\ell) \frac{e^{2\pi i b\ell} - e^{2\pi i a\ell}}{2\pi i \ell} + o(1).$$

$$\zeta(s) = \sum_{n=1}^{\infty} \frac{1}{n^s} = \prod_{p \text{ prime}} \left(1 - \frac{1}{p^s}\right)^{-1}.$$

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$$\prod_{p \text{ prime}} \left( 1 - \frac{1}{p^{s}} \right)^{-1} = \prod_{p \text{ prime}} \left( 1 + \frac{1}{p^{s}} + \frac{1}{p^{2s}} + \cdots \right)$$

$$= \left( 1 + \frac{1}{2^{s}} + \frac{1}{2^{2s}} + \cdots \right) \left( 1 + \frac{1}{3^{s}} + \frac{1}{3^{2s}} + \cdots \right)$$

$$= 1 + \frac{1}{2^{s}} + \frac{1}{3^{s}} + \frac{1}{4^{s}} + \frac{1}{5^{s}} + \frac{1}{(2 \cdot 3)^{s}} + \cdots$$

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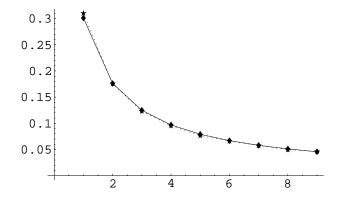
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 $\zeta(2) = \pi^2/6$  implies infinitely many primes.

$$\left|\zeta\left(\frac{1}{2}+i\frac{k}{4}\right)\right|, \ k \in \{0, 1, \dots, 65535\}.$$



The 3x + 1 Problem and Benford's Law

- Kakutani (conspiracy), Erdös (not ready).
- x odd,  $T(x) = \frac{3x+1}{2^k}$ ,  $2^k ||3x+1$ .

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- 7  $\rightarrow_1$  11  $\rightarrow_1$  17  $\rightarrow_2$  13  $\rightarrow_3$  5  $\rightarrow_4$  1  $\rightarrow_2$  1, 2-path (1,1), 5-path (1,1,2,3,4). m-path:  $(k_1,\ldots,k_m)$ .

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$$\begin{aligned} a_{n+1} &= & \mathcal{T}(a_n) \\ \mathbb{E}[\log a_{n+1}] &\approx & \sum_{k=1}^{\infty} \frac{1}{2^k} \log \left( \frac{3a_n}{2^k} \right) \\ &= & \log a_n + \log 3 - \log 2 \sum_{k=1}^{\infty} \frac{k}{2^k} \\ &= & \log a_n + \log \left( \frac{3}{4} \right). \end{aligned}$$

Geometric Brownian Motion, drift log(3/4) < 1.

$$\mathbb{P}(A) = \lim_{N \to \infty} \frac{\#\{n \le N: n \equiv 1, 5 \mod 6, n \in A\}}{\#\{n \le N: n \equiv 1, 5 \mod 6\}}.$$

$$\begin{split} \mathbb{P}(A) &= \lim_{N \to \infty} \frac{\#\{n \leq N: n \equiv 1, 5 \bmod 6, n \in A\}}{\#\{n \leq N: n \equiv 1, 5 \bmod 6\}}. \\ (k_1, \ldots, k_m): \text{ two full arithm progressions:} \\ 6 \cdot 2^{k_1 + \cdots + k_m} p + q. \end{split}$$

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# Theorem (Sinai, Kontorovich-Sinai)

 $k_i$ -values are i.i.d.r.v. (geometric, 1/2):

$$\mathbb{P}\left(\frac{\log_2\left[\frac{\mathsf{x}_m}{\left(\frac{3}{4}\right)^m\mathsf{x}_0}\right]}{\sqrt{2m}} \leq a\right) = \mathbb{P}\left(\frac{\mathsf{S}_m - 2m}{\sqrt{2m}} \leq a\right)$$

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### Structure Theorem: Sinai, Kontorovich-Sinai

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#### 3x + 1 and Benford

## Theorem (Kontorovich and M-, 2005)

As  $m \to \infty$ ,  $x_m/(3/4)^m x_0$  is Benford.

# Theorem (Lagarias-Soundararajan 2006)

 $X \ge 2^N$ , for all but at most  $c(B)N^{-1/36}X$  initial seeds the distribution of the first N iterates of the 3x + 1 map are within  $2N^{-1/36}$  of the Benford probabilities.

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Quantified Equidistribution:

$$I_{\ell} = \{\ell M, \dots, (\ell+1)M-1\}, M = m^{c}, c < 1/2$$

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$$\begin{split} &I_{\ell} = \{\ell M, \dots, (\ell+1)M-1\}, \, M = m^c, \, c < 1/2 \\ &k_1, k_2 \in I_{\ell} \colon \left| \eta \left( \frac{k_1}{\sqrt{m}} \right) - \eta \left( \frac{k_2}{\sqrt{m}} \right) \right| \, \text{small} \\ &C = \log_B 2 \, \text{of irrationality type} \, \kappa < \infty ; \end{split}$$

$$\#\{k\in I_\ell: \overline{kC}\in [a,b]\}=M(b-a)+O(M^{1+\epsilon-1/\kappa}).$$

### **Irrationality Type**

## Irrationality type

 $\alpha$  has irrationality type  $\kappa$  if  $\kappa$  is the supremum of all  $\gamma$  with

$$\underline{\lim}_{q\to\infty}q^{\gamma+1}\min_{p}\left|\alpha-\frac{p}{q}\right|=0.$$

- Algebraic irrationals: type 1 (Roth's Thm).
- Theory of Linear Forms: log<sub>8</sub> 2 of finite type.

#### **Linear Forms**

# Theorem (Baker)

 $\alpha_1, \ldots, \alpha_n$  algebraic numbers height  $A_j \geq 4$ ,  $\beta_1, \ldots, \beta_n \in \mathbb{Q}$  with height at most  $B \geq 4$ ,

$$\Lambda = \beta_1 \log \alpha_1 + \cdots + \beta_n \log \alpha_n.$$

If 
$$\Lambda \neq 0$$
 then  $|\Lambda| > B^{-C\Omega \log \Omega'}$ , with  $d = [\mathbb{Q}(\alpha_i, \beta_j) : \mathbb{Q}]$ ,  $C = (16nd)^{200n}$ ,  $\Omega = \prod_j \log A_j$ ,  $\Omega' = \Omega/\log A_n$ .

Gives  $\log_{10} 2$  of finite type, with  $\kappa < 1.2 \cdot 10^{602}$ :

$$|\log_{10} 2 - p/q| = |q \log 2 - p \log 10|/q \log 10.$$

## **Quantified Equidistribution**

## Theorem (Erdös-Turan)

$$D_{N} = \frac{\sup_{[a,b]} |N(b-a) - \#\{n \leq N : x_{n} \in [a,b]\}|}{N}$$

There is a C such that for all m:

$$D_N \leq C \cdot \left( \frac{1}{m} + \sum_{h=1}^m \frac{1}{h} \left| \frac{1}{N} \sum_{n=1}^N e^{2\pi i h x_n} \right| \right)$$

#### **Proof of Erdös-Turan**

Consider special case  $x_n = n\alpha$ ,  $\alpha \notin \mathbb{Q}$ .

- Exponential sum  $\leq \frac{1}{|\sin(\pi h\alpha)|} \leq \frac{1}{2||h\alpha||}$ .
- Must control  $\sum_{h=1}^{m} \frac{1}{h||h\alpha||}$ , see irrationality type enter.
- type  $\kappa$ ,  $\sum_{h=1}^{m} \frac{1}{h||h\alpha||} = O(m^{\kappa-1+\epsilon})$ , take  $m = \lfloor N^{1/\kappa} \rfloor$ .

## 3x + 1 Data: random 10,000 digit number, $2^k ||3x + 1|$

80,514 iterations  $((4/3)^n = a_0 \text{ predicts } 80,319);$   $\chi^2 = 13.5 \text{ (5\% } 15.5).$ 

Digit	Number	Observed	Benford
1	24251	0.301	0.301
2	14156	0.176	0.176
3	10227	0.127	0.125
4	7931	0.099	0.097
5	6359	0.079	0.079
6	5372	0.067	0.067
7	4476	0.056	0.058
8	4092	0.051	0.051
9	3650	0.045	0.046

### 3x + 1 Data: random 10,000 digit number, 2|3x + 1

241,344 iterations,  $\chi^2 = 11.4$  (5% 15.5).

Digit	Number	Observed	Benford
1	72924	0.302	0.301
2	42357	0.176	0.176
3	30201	0.125	0.125
4	23507	0.097	0.097
5	18928	0.078	0.079
6	16296	0.068	0.067
7	13702	0.057	0.058
8	12356	0.051	0.051
9	11073	0.046	0.046

## 5x + 1 Data: random 10,000 digit number, $2^{k}||5x + 1|$

27,004 iterations,  $\chi^2 = 1.8$  (5% 15.5).

Digit	Number	Observed	Benford
1	8154	0.302	0.301
2	4770	0.177	0.176
3	3405	0.126	0.125
4	2634	0.098	0.097
5	2105	0.078	0.079
6	1787	0.066	0.067
7	1568	0.058	0.058
8	1357	0.050	0.051
9	1224	0.045	0.046

### 5x + 1 Data: random 10,000 digit number, 2|5x + 1

241,344 iterations,  $\chi^2 = 3 \cdot 10^{-4}$  (5% 15.5).

Digit	Number	Observed	Benford
1	72652	0.301	0.301
2	42499	0.176	0.176
3	30153	0.125	0.125
4	23388	0.097	0.097
5	19110	0.079	0.079
6	16159	0.067	0.067
7	13995	0.058	0.058
8	12345	0.051	0.051
9	11043	0.046	0.046

# Products and Chains of Random Variables

### **Key Ingredients**

- Mellin transform and Fourier transform related by logarithmic change of variable.
- Poisson summation from collapsing to modulo 1 random variables.

#### **Preliminaries**

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♦ Proof: Prob( $\Xi_1 \cdot \Xi_2 \in [0, x]$ ):

$$\int_{t=0}^{\infty} \operatorname{Prob}\left(\Xi_{2} \in \left[0, \frac{x}{t}\right]\right) f_{1}(t) dt$$

$$= \int_{t=0}^{\infty} F_{2}\left(\frac{x}{t}\right) f_{1}(t) dt,$$

differentiate.

### **Mellin Transform**

$$(\mathcal{M}f)(s) = \int_0^\infty f(x)x^s \frac{dx}{x}$$
  
 $(\mathcal{M}^{-1}g)(x) = \frac{1}{2\pi i} \int_{c-i\infty}^{c+i\infty} g(s)x^{-s}ds$   
 $g(s) = (\mathcal{M}f)(s), f(x) = (\mathcal{M}^{-1}g)(x).$ 

$$(f_1 \star f_2)(x) = \int_0^\infty f_2\left(\frac{x}{t}\right) f_1(t) \frac{dt}{t}$$
  
$$(\mathcal{M}(f_1 \star f_2))(s) = (\mathcal{M}f_1)(s) \cdot (\mathcal{M}f_2)(s).$$

#### **Mellin Transform Formulation: Products Random Variables**

#### **Theorem**

 $X_i$ 's independent, densities  $f_i$ .  $\Xi_n = X_1 \cdots X_n$ ,

$$h_n(x_n) = (f_1 \star \cdots \star f_n)(x_n)$$
  
 $(\mathcal{M}h_n)(s) = \prod_{m=1}^n (\mathcal{M}f_m)(s).$ 

As  $n \to \infty$ ,  $\Xi_n$  becomes Benford:  $Y_n = \log_B \Xi_n$ ,  $|\operatorname{Prob}(Y_n \bmod 1 \in [a,b]) - (b-a)| \le$ 

$$(b-a)\cdot \sum_{\ell\neq 0,\ell=-\infty}^{\infty} \prod_{m=1}^{n} (\mathcal{M} \mathit{f}_{i}) \left(1-\frac{2\pi i\ell}{\log B}\right).$$

#### Conditions

•  $\{\mathcal{D}_i(\theta)\}_{i\in I}$ : one-parameter distributions, densities  $f_{\mathcal{D}_i(\theta)}$  on  $[0,\infty)$ .

#### Conditions

- $\{\mathcal{D}_i(\theta)\}_{i\in I}$ : one-parameter distributions, densities  $f_{\mathcal{D}_i(\theta)}$  on  $[0,\infty)$ .
- $\bullet \ p: \mathbb{N} \to I, \ X_1 \sim \mathcal{D}_{p(1)}(1), \ X_m \sim \mathcal{D}_{p(m)}(X_{m-1}).$

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$$f_m(x_m) = \int_0^\infty f_{\mathcal{D}_{p(m)}(1)}\left(\frac{x_m}{x_{m-1}}\right) f_{m-1}(x_{m-1}) \frac{dx_{m-1}}{x_{m-1}}$$

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$$\lim_{n\to\infty}\sum_{\ell=-\infty}^{\infty}\,\prod_{m=1}^{n}(\mathcal{M}f_{\mathcal{D}_{p(m)}(1)})\left(1-\frac{2\pi i\ell}{\log B}\right)\;=\;0$$

# Theorem (JKKKM)

- If conditions hold, as  $n \to \infty$  the distribution of leading digits of  $X_n$  tends to Benford's law.
- The error is a nice function of the Mellin transforms: if  $Y_n = \log_B X_n$ , then

$$|\operatorname{Prob}(Y_n \mod 1 \in [a, b]) - (b + a)| \le$$

$$\left| (b - a) \cdot \sum_{\substack{\ell = -\infty \ \ell \neq 0}}^{\infty} \prod_{m=1}^{n} (\mathcal{M} f_{\mathcal{D}_{p(m)}(1)}) \left( 1 - \frac{2\pi i \ell}{\log B} \right) \right|$$

### **Example:** All $X_i \sim \text{Exp}(1)$

- $X_i \sim \operatorname{Exp}(1), \ Y_n = \log_B \Xi_n$ .
- Needed ingredients:

• 
$$|P_n(s) - \log_{10}(s)| \le$$

$$\log_B s \sum_{\ell=1}^{\infty} \left( \frac{2\pi^2 \ell / \log B}{\sinh(2\pi^2 \ell / \log B)} \right)^{n/2}.$$

### **Example:** All $X_i \sim \text{Exp}(1)$

# Bounds on the error

- $|P_n(s) \log_{10} s| \le$ •  $3.3 \cdot 10^{-3} \log_B s$  if n = 2, •  $1.9 \cdot 10^{-4} \log_B s$  if n = 3,
  - $\diamond 1.1 \cdot 10^{-5} \log_B s$  if n = 5, and
  - $\diamond 3.6 \cdot 10^{-13} \log_B s \text{ if } n = 10.$
- Error at most

$$\log_{10} s \sum_{\ell=1}^{\infty} \left( \frac{17.148\ell}{\exp(8.5726\ell)} \right)^{n/2} \le .057^n \log_{10} s$$

### Conclusions

See many different systems exhibit Benford behavior.

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- Ingredients of proofs (logarithms, equidistribution).
- Applications to fraud detection / data integrity.
- Future work:
  - Study digits of other systems.
  - Develop more sophisticated tests for fraud.

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