

# From Cookie Monster to the IRS: Some Fruitful Interactions between Probability, Combinatorics and Number Theory

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## Outline

Summarize some of my interests, highlighting interplays across fields.  
Joint with many faculty, [grad students](#) and [undergrads](#).

- **Gaussian behavior in Zeckendorf decompositions:** [Gene Kopp](#), [Murat Koloğlu](#), [Yinghui Wang](#).
- **More Sums Than Differences Sets:** Peter Hegarty, [Brooke Orosz](#), [Dan Scheinerman](#).
- **Classical Random Matrix Theory:** Eduardo Dueñez, Chris Hughes, Jon Keating, Nina Snaith, [Duc Khiem Huynh](#), [Tim Novikoff](#), [Chris Hammond](#), [Steven Jackson](#), [Gene Kopp](#), [Murat Koloğlu](#), [Adam Massey](#), [Thuy Pham](#), [Anthony Sabelli](#), [John Sinsheimer](#).
- **Benford's Law:** Chaouki Abdallah, Gregory Heileman, Mark Nigrini, Fernando Perez-Gonzalez, Tu-Thach Quach, [Alex Kontorovich](#), [Dennis Jang](#), [Jung Uk Kang](#), [Alex Kruckman](#), [Jun Kudo](#).

## Fibonacci Numbers

## Previous Results

**Fibonacci Numbers:**  $F_{n+1} = F_n + F_{n-1}$ ;  
 $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, \dots$

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### Lekkerkerker's Theorem (1952)

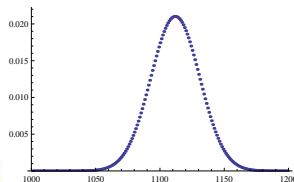
The average number of summands in the Zeckendorf decomposition for integers in  $[F_n, F_{n+1})$  tends to

$$\frac{n}{\varphi^2 + 1} \approx .276n, \text{ where } \varphi = \frac{1 + \sqrt{5}}{2} \text{ is the golden mean.}$$

# Central Limit Type Theorem

## Central Limit Type Theorem

As  $n \rightarrow \infty$ , the distribution of the number of summands, i.e.,  $a_1 + a_2 + \cdots + a_m$  in the generalized Zeckendorf decomposition  $\sum_{i=1}^m a_i H_i$  for integers in  $[H_n, H_{n+1})$  is Gaussian.



## Far-difference Representation

### Theorem (Alpert, 2009) (Analogue to Zeckendorf)

Every integer can be written uniquely as a sum of the  $\pm F_n$ 's, such that every two terms of the same (opposite) sign differ in index by at least 4 (3).

**Example:**  $1900 = F_{17} - F_{14} - F_{10} + F_6 + F_2$ .

$K$ : # of positive terms,  $L$ : # of negative terms.

### Generalized Lekkerkerker's Theorem

As  $n \rightarrow \infty$ ,  $E[K]$  and  $E[L] \rightarrow n/10$ ,  $E[K] - E[L] = \varphi/2 \approx .809$ .

### Central Limit Type Theorem

As  $n \rightarrow \infty$ ,  $K$  and  $L$  converges to a bivariate Gaussian.

- $\text{corr}(K, L) = -(21 - 2\varphi)/(29 + 2\varphi) \approx -.551$ .
- $K + L$  and  $K - L$  are independent.

## New Approach: Case of Fibonacci Numbers

$p_{n,k} = \# \{N \in [F_n, F_{n+1}): \text{the Zeckendorf decomposition of } N \text{ has exactly } k \text{ summands}\}.$

- Recurrence relation:**

$$N \in [F_{n+1}, F_{n+2}): N = F_{n+1} + F_t + \cdots, t \leq n-1.$$

$$p_{n+1,k+1} = p_{n-1,k} + p_{n-2,k} + \cdots$$

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## New Approach: Case of Fibonacci Numbers

$\rho_{n,k} = \# \{N \in [F_n, F_{n+1}): \text{the Zeckendorf decomposition of } N \text{ has exactly } k \text{ summands}\}.$

- Recurrence relation:**

$$N \in [F_{n+1}, F_{n+2}): N = F_{n+1} + F_t + \cdots, t \leq n-1.$$

$$\rho_{n+1,k+1} = \rho_{n-1,k} + \rho_{n-2,k} + \cdots$$

$$\rho_{n,k+1} = \rho_{n-2,k} + \rho_{n-3,k} + \cdots$$

$$\Rightarrow \rho_{n+1,k+1} = \rho_{n,k+1} + \rho_{n-1,k}.$$

- Generating function**  $\sum_{n,k>0} \rho_{n,k} x^k y^n = \frac{y}{1-y-xy^2}.$

- Partial fraction expansion:**

$$\frac{y}{1-y-xy^2} = \frac{-y}{y_1(x) - y_2(x)} \left( \frac{1}{y - y_1(x)} - \frac{1}{y - y_2(x)} \right)$$

where  $y_1(x)$  and  $y_2(x)$  are the roots of  $1 - y - xy^2 = 0$ .

**Coefficient of  $y^n$ :**  $g(x) = \sum_{n,k>0} \rho_{n,k} x^k.$

## New Approach: Case of Fibonacci Numbers (Continued)

$K_n$ : random variable associated with  $k$ .

$$g(x) = \sum_{n,k>0} p_{n,k} x^k.$$

- Differentiating identities:**

$$g(1) = \sum_{n,k>0} p_{n,k} = F_{n+1} - F_n,$$

$$g'(x) = \sum_{n,k>0} k p_{n,k} x^{k-1}, \quad g'(1) = g(1) E[K_n],$$

$$(xg'(x))' = \sum_{n,k>0} k^2 p_{n,k} x^{k-1},$$

$$(xg'(x))' |_{x=1} = g(1) E[K_n^2], \dots$$

- Method of moments** (for normalized  $K'_n$ ):

$$E[(K'_n)^{2m}] / (SD(K'_n))^{2m} \rightarrow (2m-1)!!,$$

$$E[(K'_n)^{2m-1}] / (SD(K'_n))^{2m-1} \rightarrow 0.$$

## New Approach: General Case

Let  $p_{n,k} = \# \{N \in [H_n, H_{n+1}) : \text{the generalized Zeckendorf decomposition of } N \text{ has exactly } k \text{ summands}\}.$

- Recurrence relation:**

Fibonacci:  $p_{n+1,k+1} = p_{n,k+1} + p_{n,k}.$

**General:**  $p_{n+1,k} = \sum_{m=0}^{L-1} \sum_{j=s_m}^{s_{m+1}-1} p_{n-m,k-j}.$   
 where  $s_0 = 0, s_m = c_1 + c_2 + \cdots + c_m.$

- Generating function:**

Fibonacci:  $\frac{y}{1-y-xy^2}.$

**General:**

$$\frac{\sum_{n \leq L} p_{n,k} x^k y^n - \sum_{m=0}^{L-1} \sum_{j=s_m}^{s_{m+1}-1} x^j y^{m+1} \sum_{n < L-m} p_{n,k} x^k y^n}{1 - \sum_{m=0}^{L-1} \sum_{j=s_m}^{s_{m+1}-1} x^j y^{m+1}}.$$

## New Approach: General Case (Continued)

- Partial fraction expansion:

Fibonacci:  $-\frac{y}{y_1(x)-y_2(x)} \left( \frac{1}{y-y_1(x)} - \frac{1}{y-y_2(x)} \right).$

General:

$$-\frac{1}{\sum_{j=s_L-1}^{s_L-1} x^j} \sum_{i=1}^L \frac{B(x, y)}{(y - y_i(x)) \prod_{j \neq i} (y_j(x) - y_i(x))}.$$

$$B(x, y) = \sum_{n \leq L} p_{n,k} x^k y^n - \sum_{m=0}^{L-1} \sum_{j=s_m}^{s_{m+1}-1} x^j y^{m+1} \sum_{n < L-m} p_{n,k} x^k y^n,$$

$$y_i(x): \text{root of } 1 - \sum_{m=0}^{L-1} \sum_{j=s_m}^{s_{m+1}-1} x^j y^{m+1} = 0.$$

**Coefficient of  $y^n$ :**  $g(x) = \sum_{n,k > 0} p_{n,k} x^k.$

- Differentiating identities
- Method of moments:  $\Rightarrow K_n \rightarrow \text{Gaussian}.$

## Random Matrix Theory

## Fundamental Problem: Spacing Between Events

**General Formulation:** Studying system, observe values at  $t_1, t_2, t_3, \dots$

**Question:** What rules govern the spacings between the  $t_i$ ?

**Examples:**

- Spacings b/w Energy Levels of Nuclei.
- Spacings b/w Eigenvalues of Matrices.
- Spacings b/w Primes.
- Spacings b/w  $n^k \alpha \bmod 1$ .
- Spacings b/w Zeros of  $L$ -functions.

## Origins of Random Matrix Theory

Classical Mechanics: 3 Body Problem Intractable.

Heavy nuclei (Uranium: 200+ protons / neutrons) worse!

Get some info by shooting high-energy neutrons into nucleus, see what comes out.

**Fundamental Equation:**

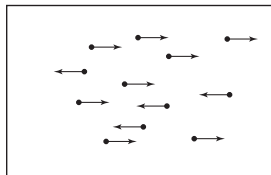
$$H\psi_n = E_n\psi_n$$

$H$  : matrix, entries depend on system

$E_n$  : energy levels

$\psi_n$  : energy eigenfunctions

## Origins of Random Matrix Theory



- Statistical Mechanics: for each configuration, calculate quantity (say pressure).
- Average over all configurations – most configurations close to system average.
- Nuclear physics: choose matrix at random, calculate eigenvalues, average over matrices (real Symmetric  $A = A^T$ , complex Hermitian  $\overline{A}^T = A$ ).

## Random Matrix Ensembles

$$A = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \cdots & a_{1N} \\ a_{12} & a_{22} & a_{23} & \cdots & a_{2N} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{1N} & a_{2N} & a_{3N} & \cdots & a_{NN} \end{pmatrix} = A^T, \quad a_{ij} = a_{ji}$$

Fix  $p$ , define

$$\text{Prob}(A) = \prod_{1 \leq i \leq j \leq N} p(a_{ij}).$$

This means

$$\text{Prob}(A : a_{ij} \in [\alpha_{ij}, \beta_{ij}]) = \prod_{1 \leq i \leq j \leq N} \int_{\alpha_{ij}}^{\beta_{ij}} p(x_{ij}) dx_{ij}.$$

Want to understand eigenvalues of  $A$ .

## Eigenvalue Distribution

$\delta(\mathbf{x} - \mathbf{x}_0)$  is a unit point mass at  $\mathbf{x}_0$ :

$$\int f(\mathbf{x})\delta(\mathbf{x} - \mathbf{x}_0)d\mathbf{x} = f(\mathbf{x}_0).$$

To each  $A$ , attach a probability measure:

$$\begin{aligned}\mu_{A,N}(\mathbf{x}) &= \frac{1}{N} \sum_{i=1}^N \delta\left(\mathbf{x} - \frac{\lambda_i(A)}{2\sqrt{N}}\right) \\ \int_a^b \mu_{A,N}(\mathbf{x}) d\mathbf{x} &= \frac{\#\left\{\lambda_i : \frac{\lambda_i(A)}{2\sqrt{N}} \in [a, b]\right\}}{N} \\ \text{k}^{\text{th}} \text{ moment} &= \frac{\sum_{i=1}^N \lambda_i(A)^k}{2^k N^{\frac{k}{2}+1}}.\end{aligned}$$

## SKETCH OF PROOF: Eigenvalue Trace Lemma

Want to understand  $A$ 's eigenvalues, but it's  $A$ 's elements that are chosen randomly and independently.

### Eigenvalue Trace Lemma

Let  $A$  be an  $N \times N$  matrix with eigenvalues  $\lambda_i(A)$ . Then

$$\text{Trace}(A^k) = \sum_{n=1}^N \lambda_i(A)^k,$$

where

$$\text{Trace}(A^k) = \sum_{i_1=1}^N \cdots \sum_{i_k=1}^N a_{i_1 i_2} a_{i_2 i_3} \cdots a_{i_N i_1}.$$

## SKETCH OF PROOF: Correct Scale

$$\text{Trace}(A^2) = \sum_{i=1}^N \lambda_i(A)^2.$$

By the Central Limit Theorem:

$$\text{Trace}(A^2) = \sum_{i=1}^N \sum_{j=1}^N a_{ij} a_{ji} = \sum_{i=1}^N \sum_{j=1}^N a_{ij}^2 \sim N^2$$

$$\sum_{i=1}^N \lambda_i(A)^2 \sim N^2$$

Gives  $N \text{Ave}(\lambda_i(A)^2) \sim N^2$  or  $\text{Ave}(\lambda_i(A)) \sim \sqrt{N}$ .

## SKETCH OF PROOF: Averaging Formula

Recall  $k$ -th moment of  $\mu_{A,N}(x)$  is  $\text{Trace}(A^k)/2^k N^{k/2+1}$ .

Average  $k$ -th moment is

$$\int \cdots \int \frac{\text{Trace}(A^k)}{2^k N^{k/2+1}} \prod_{i \leq j} p(a_{ij}) da_{ij}.$$

Proof by method of moments: Two steps

- Show average of  $k$ -th moments converge to moments of semi-circle as  $N \rightarrow \infty$ ;
- Control variance (show it tends to zero as  $N \rightarrow \infty$ ).

## SKETCH OF PROOF: Averaging Formula for Second Moment

Substituting into expansion gives

$$\frac{1}{2^2 N^2} \int_{-\infty}^{\infty} \cdots \int_{-\infty}^{\infty} \sum_{i=1}^N \sum_{j=1}^N a_{ij}^2 \cdot p(a_{11}) da_{11} \cdots p(a_{NN}) da_{NN}$$

Integration factors as

$$\int_{a_{ij}=-\infty}^{\infty} a_{ij}^2 p(a_{ij}) da_{ij} \cdot \prod_{\substack{(k,l) \neq (i,j) \\ k < l}} \int_{a_{kl}=-\infty}^{\infty} p(a_{kl}) da_{kl} = 1.$$

Higher moments involve more advanced combinatorics (Catalan numbers).

## SKETCH OF PROOF: Averaging Formula for Higher Moments

Higher moments involve more advanced combinatorics (Catalan numbers).

$$\frac{1}{2^k N^{k/2+1}} \int_{-\infty}^{\infty} \cdots \int_{-\infty}^{\infty} \sum_{i_1=1}^N \cdots \sum_{i_k=1}^N a_{i_1 i_2} \cdots a_{i_k i_1} \cdot \prod_{i \leq j} p(a_{ij}) da_{ij}.$$

Main contribution when the  $a_{i_\ell i_{\ell+1}}$ 's matched in pairs, not all matchings contribute equally (if did would get a Gaussian and not a semi-circle; this is seen in Real Symmetric Palindromic Toeplitz matrices).

## Real Symmetric *m*-Circulant Ensemble

## Circulant Matrices

Study circulant matrices, period  $m$  on diagonals.

6-by-6 real symmetric period 2-circulant matrix:

$$\begin{pmatrix} c_0 & c_1 & c_2 & c_3 & c_2 & d_1 \\ c_1 & d_0 & d_1 & d_2 & c_3 & d_2 \\ c_2 & d_1 & c_0 & c_1 & c_2 & c_3 \\ c_3 & d_2 & c_1 & d_0 & d_1 & d_2 \\ c_2 & c_3 & c_2 & d_1 & c_0 & c_1 \\ d_1 & d_2 & c_3 & d_2 & c_1 & d_0 \end{pmatrix}.$$

Look at the *expected value* for the moments:

$$\begin{aligned} M_n(N) &:= \mathbb{E}(M_n(A, N)) \\ &= \frac{1}{N^{\frac{n}{2}+1}} \sum_{1 \leq i_1, \dots, i_n \leq N} \mathbb{E}(a_{i_1 i_2} a_{i_2 i_3} \cdots a_{i_n i_1}). \end{aligned}$$

## Matchings

Rewrite:

$$M_n(N) = \frac{1}{N^{\frac{n}{2}+1}} \sum_{\sim} \eta(\sim) m_{d_1(\sim)} \cdots m_{d_l(\sim)}.$$

where the sum is over equivalence relations on  $\{(1, 2), (2, 3), \dots, (n, 1)\}$ . The  $d_j(\sim)$  denote the sizes of the equivalence classes, and the  $m_d$  the moments of  $p$ .

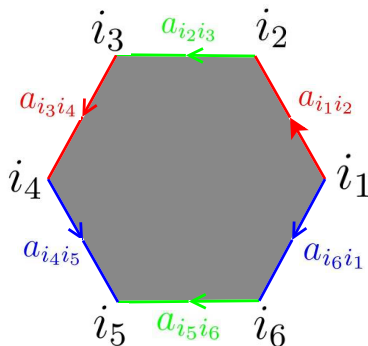
Finally, the coefficient  $\eta(\sim)$  is the number of solutions to the system of Diophantine equations:

Whenever  $(s, s+1) \sim (t, t+1)$ ,

- $i_{s+1} - i_s \equiv i_{t+1} - i_t \pmod{N}$  and  $i_s \equiv i_t \pmod{m}$ , or
- $i_{s+1} - i_s \equiv -(i_{t+1} - i_t) \pmod{N}$  and  $i_s \equiv i_{t+1} \pmod{m}$ .

## Matchings

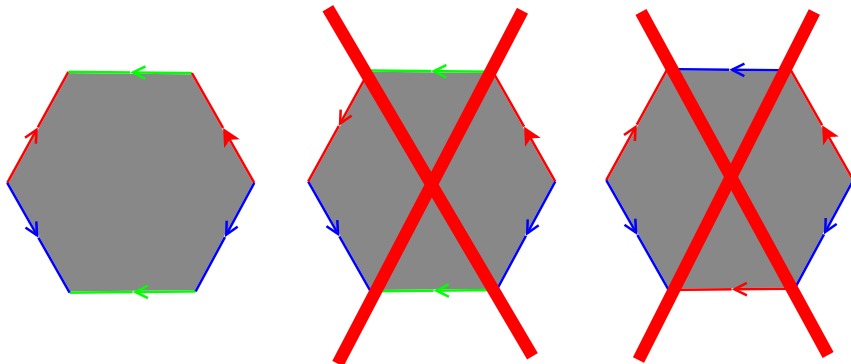
- $i_{s+1} - i_s \equiv i_{t+1} - i_t \pmod{N}$  and  $i_s \equiv i_t \pmod{m}$ , or
- $i_{s+1} - i_s \equiv -(i_{t+1} - i_t) \pmod{N}$  and  $i_s \equiv i_{t+1} \pmod{m}$ .



**Figure:** Red edges same orientation and blue, green opposite.

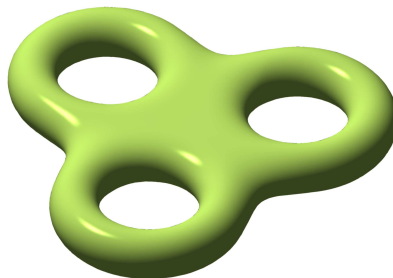
## Contributing Terms

As  $N \rightarrow \infty$ , the only terms that contribute to this sum are those in which the entries are matched in pairs and with opposite orientation.



## Contributing Terms: Algebraic Topology

Think of pairings as topological identifications, the contributing ones give rise to orientable surfaces.



Contribution from such a pairing is  $m^{-2g}$ , where  $g$  is the genus (number of holes) of the surface. Proof: combinatorial argument involving Euler characteristic.

## Computing the Even Moments

### Theorem: Even Moment Formula

$$M_{2k} = \sum_{g=0}^{\lfloor k/2 \rfloor} \varepsilon_g(k) m^{-2g} + O_k \left( \frac{1}{N} \right),$$

with  $\varepsilon_g(k)$  the number of pairings of the edges of a  $(2k)$ -gon giving rise to a genus  $g$  surface.

J. Harer and D. Zagier (1986) gave generating functions for the  $\varepsilon_g(k)$ .

## Computing the Even Moments

### Harer and Zagier

$$\sum_{g=0}^{\lfloor k/2 \rfloor} \varepsilon_g(k) r^{k+1-2g} = (2k-1)!! c(k, r)$$

where

$$1 + 2 \sum_{k=0}^{\infty} c(k, r) x^{k+1} = \left( \frac{1+x}{1-x} \right)^r.$$

Thus, we write

$$M_{2k} = m^{-(k+1)} (2k-1)!! c(k, m).$$

## Computing the Even Moments

A multiplicative convolution and Cauchy's residue formula yields the *characteristic function* of the distribution (inverse Fourier transform of the density).

$$\begin{aligned}
 \phi(t) &= \sum_{k=0}^{\infty} \frac{(it)^{2k} M_{2k}}{(2k)!} \\
 &= \frac{1}{2\pi im} \oint_{|z|=2} \frac{1}{2z^{-1}} \left( \left( \frac{1+z^{-1}}{1-z^{-1}} \right)^m - 1 \right) e^{-t^2 z/2m} \frac{dz}{z} \\
 &= \frac{1}{m} e^{\frac{-t^2}{2m}} \sum_{l=1}^m \binom{m}{l} \frac{1}{(l-1)!} \left( \frac{-t^2}{m} \right)^{l-1}.
 \end{aligned}$$

## Results

Fourier transform and algebra yields

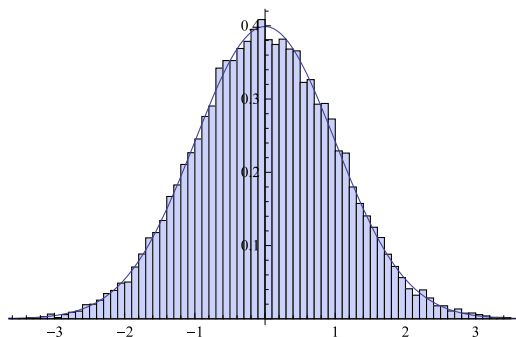
### Theorem: Kopp, Koloğlu and M–

The limiting spectral density function  $f_m(x)$  of the real symmetric  $m$ -circulant ensemble is given by the formula

$$f_m(x) = \frac{e^{-\frac{mx^2}{2}}}{\sqrt{2\pi m}} \sum_{r=0}^m \frac{1}{(2r)!} \sum_{s=0}^{m-r} \binom{m}{r+s+1} \frac{(2r+2s)!}{(r+s)!s!} \left(-\frac{1}{2}\right)^s (mx^2)^r.$$

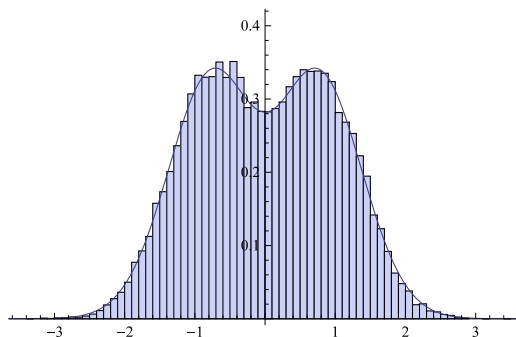
As  $m \rightarrow \infty$ , the limiting spectral densities approach the semicircle distribution.

## Results (continued)



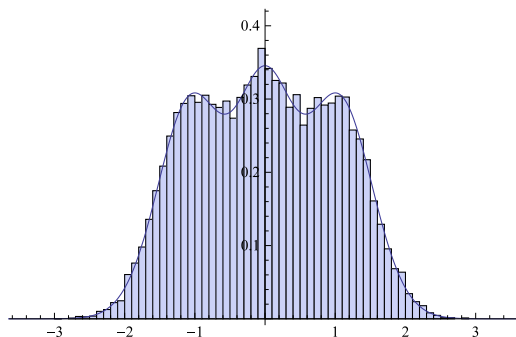
**Figure:** Plot for  $f_1$  and histogram of eigenvalues of 100 circulant matrices of size  $400 \times 400$ .

## Results (continued)



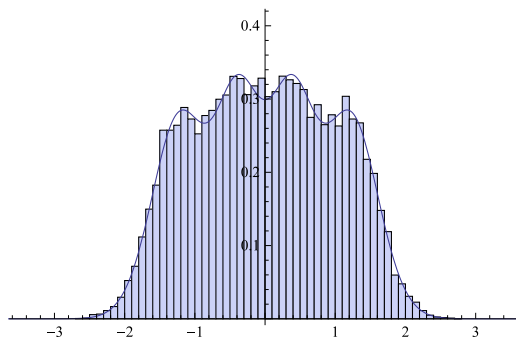
**Figure:** Plot for  $f_2$  and histogram of eigenvalues of 100 2-circulant matrices of size  $400 \times 400$ .

## Results (continued)



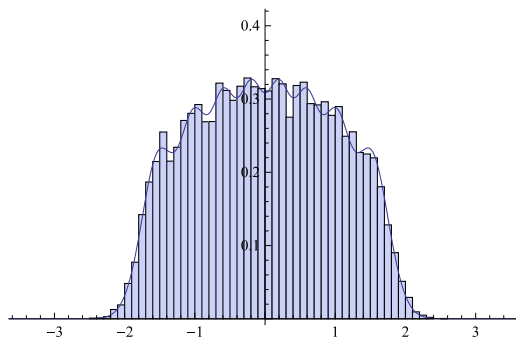
**Figure:** Plot for  $f_3$  and histogram of eigenvalues of 100 3-circulant matrices of size  $402 \times 402$ .

## Results (continued)



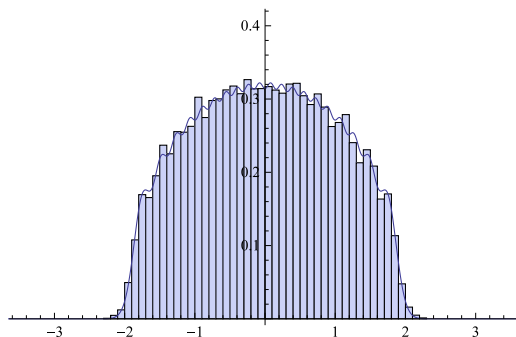
**Figure:** Plot for  $f_4$  and histogram of eigenvalues of 100 4-circulant matrices of size  $400 \times 400$ .

## Results (continued)



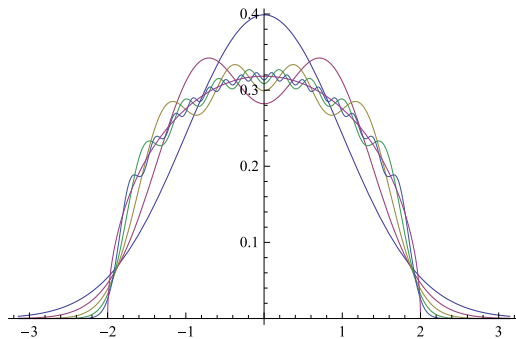
**Figure:** Plot for  $f_8$  and histogram of eigenvalues of 100 8-circulant matrices of size  $400 \times 400$ .

## Results (continued)



**Figure:** Plot for  $f_{20}$  and histogram of eigenvalues of 100 20-circulant matrices of size  $400 \times 400$ .

## Results (continued)



**Figure:** Plot of convergence to the semi-circle.

## Benford History and Applications

## Benford's Law: Newcomb (1881), Benford (1938)

### Statement

For many data sets, probability of observing a first digit of  $d$  base  $B$  is  $\log_B \left( \frac{d+1}{d} \right)$ .

First 60 values of  $2^n$  (only displaying 30)

			digit	#	Obs Prob	Benf Prob
1	1024	1048576				
2	2048	2097152	1	18	.300	.301
4	4096	4194304	2	12	.200	.176
8	8192	8388608	3	6	.100	.125
16	16384	16777216	4	6	.100	.097
32	32768	33554432	5	6	.100	.079
64	65536	67108864	6	4	.067	.067
128	131072	134217728	7	2	.033	.058
256	262144	268435456	8	5	.083	.051
512	524288	536870912	9	1	.017	.046

## Examples

- recurrence relations
- special functions (such as  $n!$ )
- iterates of power, exponential, rational maps
- products of random variables
- $L$ -functions, characteristic polynomials
- iterates of the  $3x + 1$  map
- differences of order statistics
- hydrology and financial data

## Applications

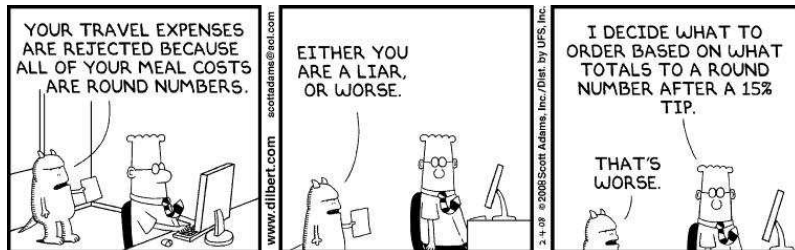
- analyzing round-off errors
- determining the optimal way to store numbers
- detecting tax fraud and data integrity

## Caveats!

- A math test indicating fraud is *not* proof of fraud: unlikely events, alternate reasons.

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## Detecting Fraud

### Bank Fraud

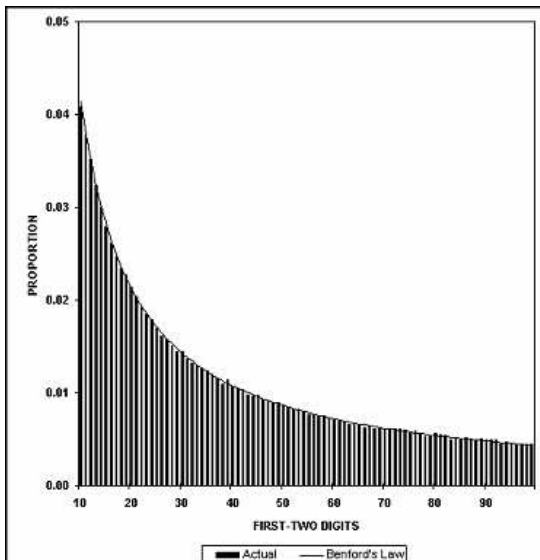
- Bank audit: huge spike of numbers starting with 48 and 49, most due to one person.

## Detecting Fraud

### Bank Fraud

- Bank audit: huge spike of numbers starting with 48 and 49, most due to one person.
- Write-off limit of \$5,000. Officer had friends applying, run up balances just under \$5,000....

# Data Integrity: Stream Flow Statistics: 130 years, 457,440 records



## Election Fraud: Iran 2009

Numerous protests/complaints over Iran's 2009 elections.

Lot of analysis; data moderately suspicious:

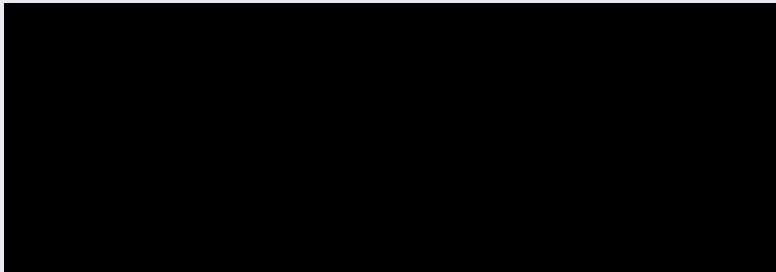
- First and second leading digits;
- Last two digits (should almost be uniform);
- Last two digits differing by at least 2.

Warning: enough tests, even if nothing wrong will find a suspicious result (but when all tests are on the boundary...).

## New Test for Fraud

## New Test for Fraud

**Victoria Cuff, Allie Lewis, M– (2010)**



## Image Analysis

- Pictures aren't Benford's law, but coefficients of Discrete Cosine Transform (DCT) very close (slightly modified law).
- Analysis of coefficients, from Generalized Gaussian Distributions:  

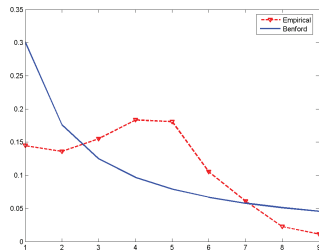
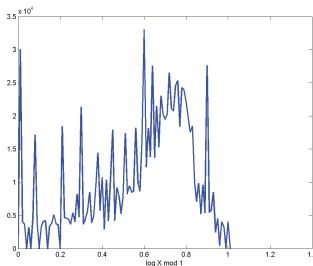
$$f_X(x) = A \exp(-\|\beta x\|^c).$$
- Application: detect compression, steganography (hidden message in picture by modifying least significant bit in pixels).

## Image Analysis (continued)



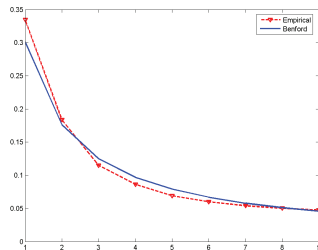
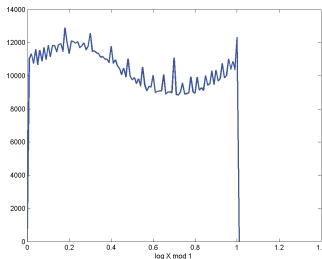
Figure 'Man' used in the experiments.

## Image Analysis (continued)



(a) Histogram of the luminance values of 'Man' in Benford ( $\log_{10} \bmod 1$ ) domain; (b) Distribution of first digits from 'Man'.

## Image Analysis (continued)



Histogram of the DCT values of 'Man' in Benford ( $\log_{10} \bmod 1$ ) domain; (b) Distribution of first digits from 'Man'.

## Logarithms and Benford's Law

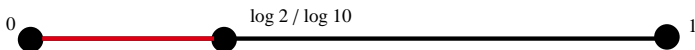
### Fundamental Equivalence

Data set  $\{x_i\}$  is Benford base  $B$  if  $\{y_i\}$  is equidistributed mod 1, where  $y_i = \log_B x_i$ .

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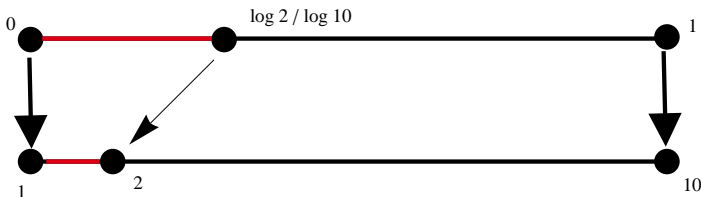
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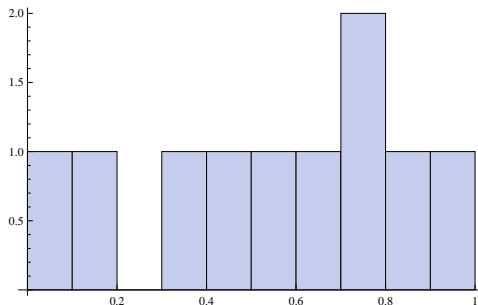
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### Kronecker-Weyl Theorem

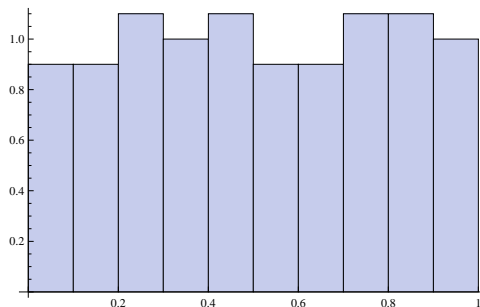
If  $\beta \notin \mathbb{Q}$  then  $n\beta \bmod 1$  is equidistributed.  
(Thus if  $\log_B \alpha \notin \mathbb{Q}$ , then  $\alpha^n$  is Benford.)

## Example of Equidistribution: $n\sqrt{\pi} \bmod 1$



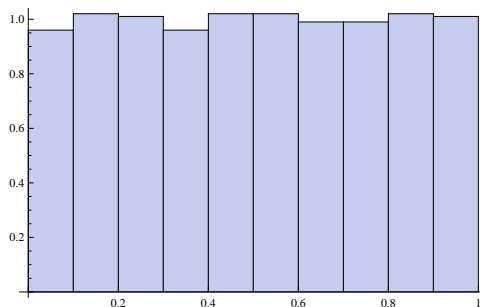
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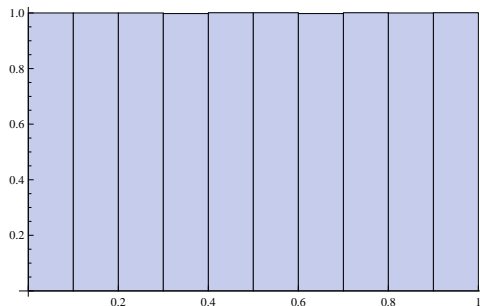
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$n\sqrt{\pi} \bmod 1$  for  $n \leq 1000$

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$n\sqrt{\pi} \bmod 1$  for  $n \leq 10,000$

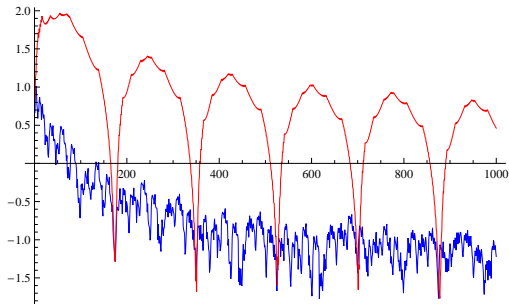
## Logarithms and Benford's Law

$\chi^2$  values for  $\alpha^n$ ,  $1 \leq n \leq N$  (5% 15.5).

$N$	$\chi^2(\gamma)$	$\chi^2(e)$	$\chi^2(\pi)$
100	0.72	0.30	46.65
200	0.24	0.30	8.58
400	0.14	0.10	10.55
500	0.08	0.07	2.69
700	0.19	0.04	0.05
800	0.04	0.03	6.19
900	0.09	0.09	1.71
1000	0.02	0.06	2.90

## Logarithms and Benford's Law: Base 10

$\log(\chi^2)$  vs  $N$  for  $\pi^n$  (red) and  $e^n$  (blue),  
 $n \in \{1, \dots, N\}$ . Note  $\pi^{175} \approx 1.0028 \cdot 10^{87}$ , (5%,  
 $\log(\chi^2) \approx 2.74$ ).



## Benford Good Processes

## Poisson Summation and Benford's Law: Definitions

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- Poisson Summation Formula:  $f$  nice:

$$\sum_{\ell=-\infty}^{\infty} f(\ell) = \sum_{\ell=-\infty}^{\infty} \hat{f}(\ell),$$

$$\text{Fourier transform } \hat{f}(\xi) = \int_{-\infty}^{\infty} f(x) e^{-2\pi i x \xi} dx.$$

## Benford Good Process

$X_T$  is **Benford Good** if there is a nice  $f$  st

$$\text{CDF}_{\vec{Y}_{T,B}}(y) = \int_{-\infty}^y \frac{1}{T} f\left(\frac{t}{T}\right) dt + E_T(y) := G_T(y)$$

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 $\sum_{|\ell| \leq Th(T)} [E_T(b + \ell) - E_T(a + \ell)] = o(1).$

## Main Theorem

### Theorem (Kontorovich and M–, 2005)

$X_T$  converging to  $X$  as  $T \rightarrow \infty$  (think spreading Gaussian). If  $X_T$  is Benford good, then  $X$  is Benford.

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- Examples

- ◇  $L$ -functions
- ◇ characteristic polynomials (RMT)
- ◇  $3x + 1$  problem
- ◇ geometric Brownian motion.

## Sketch of the proof

- **Structure Theorem:**
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  - ◇ hardest step
  - ◇ techniques problem specific

## Sketch of the proof (continued)

$$\sum_{\ell=-\infty}^{\infty} \mathbb{P} \left( \mathbf{a} + \ell \leq \vec{Y}_{T,B} \leq \mathbf{b} + \ell \right)$$

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 &= \widehat{f}(0) \cdot (\mathbf{b} - \mathbf{a}) + \sum_{\ell \neq 0} \widehat{f}(T\ell) \frac{e^{2\pi i b \ell} - e^{2\pi i a \ell}}{2\pi i \ell} + o(1).
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## Some Results

### Theorem (Kontorovich and M–, 2005)

*$L(s, f)$  a good  $L$ -function, as  $T \rightarrow \infty$ ,  
 $L(\sigma_T + it, f)$  is Benford.*

### Theorem (Kontorovich and M–, 2005)

*As  $N \rightarrow \infty$ , the distribution of digits of the  
absolute values of the characteristic  
polynomials of  $N \times N$  unitary matrices (with  
respect to Haar measure) converges to the  
Benford probabilities.*

# The $3x + 1$ Problem and Benford's Law

## 3x + 1 Problem

- Kakutani (conspiracy), Erdős (not ready).
- $x$  odd,  $T(x) = \frac{3x+1}{2^k}$ ,  $2^k || 3x + 1$ .

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 2-path (1, 1), 5-path (1, 1, 2, 3, 4).  
*m*-path:  $(k_1, \dots, k_m)$ .

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 &= \log a_n + \log \left( \frac{3}{4} \right).
 \end{aligned}$$

Geometric Brownian Motion, drift  $\log(3/4) < 1$ .

## Structure Theorem: Sinai, Kontorovich-Sinai

$$\mathbb{P}(A) = \lim_{N \rightarrow \infty} \frac{\#\{n \leq N: n \equiv 1, 5 \pmod{6}, n \in A\}}{\#\{n \leq N: n \equiv 1, 5 \pmod{6}\}}.$$

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## 3x + 1 and Benford

### Theorem (Kontorovich and M–, 2005)

*As  $m \rightarrow \infty$ ,  $x_m / (3/4)^m x_0$  is Benford.*

### Theorem (Lagarias-Soundararajan 2006)

*$X \geq 2^N$ , for all but at most  $c(B)N^{-1/36}X$  initial seeds the distribution of the first  $N$  iterates of the  $3x + 1$  map are within  $2N^{-1/36}$  of the Benford probabilities.*

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- **Quantified Equidistribution:**  $I_\ell = \{\ell M, \dots, (\ell + 1)M - 1\}$ ,  
 $M = m^c$ ,  $c < 1/2$   
 $k_1, k_2 \in I_\ell$ :  $\left| \eta\left(\frac{k_1}{\sqrt{m}}\right) - \eta\left(\frac{k_2}{\sqrt{m}}\right) \right|$  small  
 $C = \log_B 2$  of irrationality type  $\kappa < \infty$ :

$$\#\{k \in I_\ell : \overline{kC} \in [a, b]\} = M(b - a) + O(M^{1+\epsilon-1/\kappa}).$$

# Irrationality Type

## Irrationality type

$\alpha$  has irrationality type  $\kappa$  if  $\kappa$  is the supremum of all  $\gamma$  with

$$\lim_{q \rightarrow \infty} q^{\gamma+1} \min_p \left| \alpha - \frac{p}{q} \right| = 0.$$

- Algebraic irrationals: type 1 (Roth's Thm).
- Theory of Linear Forms:  $\log_B 2$  of finite type.

## Linear Forms

### Theorem (Baker)

$\alpha_1, \dots, \alpha_n$  algebraic numbers height  $A_j \geq 4$ ,  
 $\beta_1, \dots, \beta_n \in \mathbb{Q}$  with height at most  $B \geq 4$ ,

$$\Lambda = \beta_1 \log \alpha_1 + \dots + \beta_n \log \alpha_n.$$

If  $\Lambda \neq 0$  then  $|\Lambda| > B^{-C\Omega \log \Omega'}$ , with  
 $d = [\mathbb{Q}(\alpha_i, \beta_j) : \mathbb{Q}]$ ,  $C = (16nd)^{200n}$ ,  
 $\Omega = \prod_j \log A_j$ ,  $\Omega' = \Omega / \log A_n$ .

Gives  $\log_{10} 2$  of finite type, with  $\kappa < 1.2 \cdot 10^{602}$ :

$$|\log_{10} 2 - p/q| = |q \log 2 - p \log 10| / q \log 10.$$

## Quantified Equidistribution

### Theorem (Erdős-Turan)

$$D_N = \frac{\sup_{[a,b]} |N(b-a) - \#\{n \leq N : x_n \in [a,b]\}|}{N}$$

*There is a  $C$  such that for all  $m$ :*

$$D_N \leq C \cdot \left( \frac{1}{m} + \sum_{h=1}^m \frac{1}{h} \left| \frac{1}{N} \sum_{n=1}^N e^{2\pi i h x_n} \right| \right)$$

## Proof of Erdős-Turan

Consider special case  $x_n = n\alpha$ ,  $\alpha \notin \mathbb{Q}$ .

- Exponential sum  $\leq \frac{1}{|\sin(\pi h\alpha)|} \leq \frac{1}{2||h\alpha||}$ .
- Must control  $\sum_{h=1}^m \frac{1}{h||h\alpha||}$ , see irrationality type enter.
- type  $\kappa$ ,  $\sum_{h=1}^m \frac{1}{h||h\alpha||} = O(m^{\kappa-1+\epsilon})$ , take  $m = \lfloor N^{1/\kappa} \rfloor$ .

# 3x + 1 Data: random 10,000 digit number, $2^k || 3x + 1$

80,514 iterations ( $(4/3)^n = a_0$  predicts 80,319);  
 $\chi^2 = 13.5$  (5% 15.5).

Digit	Number	Observed	Benford
1	24251	0.301	0.301
2	14156	0.176	0.176
3	10227	0.127	0.125
4	7931	0.099	0.097
5	6359	0.079	0.079
6	5372	0.067	0.067
7	4476	0.056	0.058
8	4092	0.051	0.051
9	3650	0.045	0.046

## 3x + 1 Data: random 10,000 digit number, 2|3x + 1

241,344 iterations,  $\chi^2 = 11.4$  (5% 15.5).

Digit	Number	Observed	Benford
1	72924	0.302	0.301
2	42357	0.176	0.176
3	30201	0.125	0.125
4	23507	0.097	0.097
5	18928	0.078	0.079
6	16296	0.068	0.067
7	13702	0.057	0.058
8	12356	0.051	0.051
9	11073	0.046	0.046

**5x + 1 Data: random 10,000 digit number,  $2^k \parallel 5x + 1$**

27,004 iterations,  $\chi^2 = 1.8$  (5% 15.5).

Digit	Number	Observed	Benford
1	8154	0.302	0.301
2	4770	0.177	0.176
3	3405	0.126	0.125
4	2634	0.098	0.097
5	2105	0.078	0.079
6	1787	0.066	0.067
7	1568	0.058	0.058
8	1357	0.050	0.051
9	1224	0.045	0.046

## 5x + 1 Data: random 10,000 digit number, 2|5x + 1

241,344 iterations,  $\chi^2 = 3 \cdot 10^{-4}$  (5% 15.5).

Digit	Number	Observed	Benford
1	72652	0.301	0.301
2	42499	0.176	0.176
3	30153	0.125	0.125
4	23388	0.097	0.097
5	19110	0.079	0.079
6	16159	0.067	0.067
7	13995	0.058	0.058
8	12345	0.051	0.051
9	11043	0.046	0.046

## Products and Chains of Random Variables

## Preliminaries

- $X_1 \cdots X_n \Leftrightarrow Y_1 + \cdots + Y_n \bmod 1$ ,  $Y_i = \log_B X_i$
- Density  $Y_i$  is  $g_i$ , density  $Y_i + Y_j$  is

$$(g_i * g_j)(y) = \int_0^1 g_i(t)g_j(y - t)dt.$$

- $h_n = g_1 * \cdots * g_n$ ,  $\widehat{g}(\xi) = \widehat{g}_1(\xi) \cdots \widehat{g}_n(\xi)$ .
- Dirac delta functional:  $\int \delta_\alpha(y)g(y)dy = g(\alpha)$ .

## Fourier input

- Fejér kernel:

$$F_N(x) = \sum_{n=-N}^N \left(1 - \frac{|n|}{N}\right) e^{2\pi i n x}.$$

- Fejér series:  $T_N f(x) =$

$$(f * F_N)(x) = \sum_{n=-N}^N \left(1 - \frac{|n|}{N}\right) \widehat{f}(n) e^{2\pi i n x}.$$

- Lebesgue's Theorem:  $f \in L^1([0, 1])$ . As  $N \rightarrow \infty$ ,  $T_N f$  converges to  $f$  in  $L^1([0, 1])$ .
- $T_N(f * g) = (T_N f) * g$ : convolution assoc.

## Modulo 1 Central Limit Theorem

### Theorem (M– and Nigrini 2007)

$\{Y_m\}$  independent continuous random variables on  $[0, 1)$  (not necc. i.i.d.), densities  $\{g_m\}$ .

$Y_1 + \cdots + Y_M \bmod 1$  converges to the uniform distribution as  $M \rightarrow \infty$  in  $L^1([0, 1])$  iff  $\forall n \neq 0$ ,  
 $\lim_{M \rightarrow \infty} \widehat{g}_1(n) \cdots \widehat{g}_M(n) = 0$ .

## Generalizations

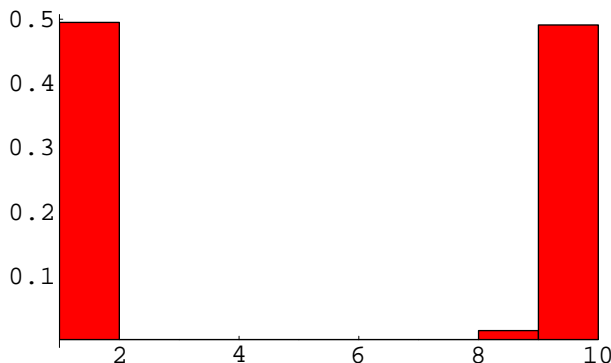
- Levy proved for i.i.d.r.v. just one year after Benford's paper.
- Generalized to other compact groups, with estimates on the rate of convergence.
  - ◇ Stromberg:  $n$ -fold convolution of a regular probability measure on a compact Hausdorff group  $G$  converges to normalized Haar measure in weak-star topology iff support of the distribution not contained in a coset of a proper normal closed subgroup of  $G$ .

## Non-Benford Product

Distribution of digits (base 10) of 1000 products

$X_1 \cdots X_{1000}$ , where  $g_{10,m} = \phi_{11^m}$ .

$\phi_m(x) = m$  if  $|x - 1/8| \leq 1/2m$  (0 otherwise).



## Proof of Modulo 1 CLT

- Density of sum is  $h_\ell = g_1 * \cdots * g_\ell$ .

- Suffices show  $\forall \epsilon$ :

$$\lim_{M \rightarrow \infty} \int_0^1 |h_M(x) - 1| dx < \epsilon.$$

- Lebesgue's Theorem:  $N$  large,

$$\|h_1 - T_N h_1\|_1 =$$

$$\int_0^1 |h_1(x) - T_N h_1(x)| dx < \frac{\epsilon}{2}.$$

- Claim: above holds for  $h_M$  for all  $M$ .

## Proof of Modulo 1 CLT

Show  $\lim_{M \rightarrow \infty} \|h_M - 1\|_1 = 0$ .

Triangle inequality:

$$\|h_M - 1\|_1 \leq \|h_M - T_N h_M\|_1 + \|T_N h_M - 1\|_1.$$

Choices of  $N$  and  $\epsilon$ :

$$\|h_M - T_N h_M\|_1 < \epsilon/2.$$

Show  $\|T_N h_M - 1\|_1 < \epsilon/2$ .

## Proof of Modulo 1 CLT

$$\begin{aligned} \|T_N h_M - 1\|_1 &= \int_0^1 \left| \sum_{\substack{n=-N \\ n \neq 0}}^N \left(1 - \frac{|n|}{N}\right) \widehat{h}_M(n) e^{2\pi i n x} \right| dx \\ &\leq \sum_{\substack{n=-N \\ n \neq 0}}^N \left(1 - \frac{|n|}{N}\right) |\widehat{h}_M(n)| \end{aligned}$$

$$\widehat{h}_M(n) = \widehat{g}_1(n) \cdots \widehat{g}_M(n) \xrightarrow{M \rightarrow \infty} 0.$$

For fixed  $N$  and  $\epsilon$ , choose  $M$  large so that

$$|\widehat{h}_M(n)| < \epsilon/4N \text{ whenever } n \neq 0 \text{ and } |n| \leq N.$$

## Conditions for Chains of Random Variables

### Conditions

- $\{\mathcal{D}_i(\theta)\}_{i \in I}$ : one-parameter distributions, densities  $f_{\mathcal{D}_i(\theta)}$  on  $[0, \infty)$ .
- $p : \mathbb{N} \rightarrow I$ ,  $X_1 \sim \mathcal{D}_{p(1)}(1)$ ,  $X_m \sim \mathcal{D}_{p(m)}(X_{m-1})$ .
- $m \geq 2$ ,  $f_m(x_m) =$

$$\int_0^\infty f_{\mathcal{D}_{p(m)}(1)} \left( \frac{x_m}{x_{m-1}} \right) f_{m-1}(x_{m-1}) \frac{dx_{m-1}}{x_{m-1}}$$

•

$$\lim_{n \rightarrow \infty} \sum_{\substack{\ell = -\infty \\ \ell \neq 0}}^{\infty} \prod_{m=1}^n (\mathcal{M} f_{\mathcal{D}_{p(m)}(1)}) \left( 1 - \frac{2\pi i \ell}{\log B} \right) = 0$$

## Behavior of Chains of Random Variables

### Theorem (JKKKM)

- *If conditions hold, as  $n \rightarrow \infty$  the distribution of leading digits of  $X_n$  tends to Benford's law.*
- *The error is a nice function of the Mellin transforms: if  $Y_n = \log_B X_n$ , then*

$$\begin{aligned}
 & |\text{Prob}(Y_n \bmod 1 \in [a, b]) - (b - a)| \leq \\
 & \left| (b - a) \cdot \sum_{\substack{\ell=-\infty \\ \ell \neq 0}}^{\infty} \prod_{m=1}^n (\mathcal{M} f_{\mathcal{D}_{p(m)}(1)}) \left( 1 - \frac{2\pi i \ell}{\log B} \right) \right|
 \end{aligned}$$

## Example: All $X_i \sim \text{Unif}(0, k)$

- $X_i \sim \text{Unif}(0, k)$ : without loss of generality  $k \in [1, 10)$ .
- $P_n(s) = \text{Prob}(M_{10}(\Xi_n) \leq s)$ .
- $|P_n(s) - \log_{10}(s)| \leq$

$$\frac{k (\log k)^{n-1}}{s \Gamma(n)} + \left( \frac{1}{2.9^n} + \frac{\zeta(n) - 1}{2.7^n} \right) 2 \log_{10} s.$$

## Example: All $X_i \sim \text{Exp}(1)$

- $X_i \sim \text{Exp}(1)$ ,  $Y_n = \log_B \Xi_n$ .
- Needed ingredients:
  - ◇  $\int_0^\infty \exp(-x) x^{s-1} dx = \Gamma(s)$ .
  - ◇  $|\Gamma(1 + ix)| = \sqrt{\pi x / \sinh(\pi x)}$ ,  $x \in \mathbb{R}$ .
- $|P_n(s) - \log_{10}(s)| \leq$

$$\log_B s \sum_{\ell=1}^{\infty} \left( \frac{2\pi^2 \ell / \log B}{\sinh(2\pi^2 \ell / \log B)} \right)^{n/2}.$$

## Example: All $X_i \sim \text{Exp}(1)$ (continued)

### Bounds on the error

- $|P_n(s) - \log_{10} s| \leq$ 
  - ◇  $3.3 \cdot 10^{-3} \log_B s$  if  $n = 2$ ,
  - ◇  $1.9 \cdot 10^{-4} \log_B s$  if  $n = 3$ ,
  - ◇  $1.1 \cdot 10^{-5} \log_B s$  if  $n = 5$ , and
  - ◇  $3.6 \cdot 10^{-13} \log_B s$  if  $n = 10$ .
- Error at most

$$\log_{10} s \sum_{\ell=1}^{\infty} \left( \frac{17.148\ell}{\exp(8.5726\ell)} \right)^{n/2} \leq .057^n \log_{10} s$$

## More Sums Than Difference Sets

## Statement

A finite set of integers,  $|A|$  its size. Form

- Sumset:  $A + A = \{a_i + a_j : a_i, a_j \in A\}$ .
- Difference set:  $A - A = \{a_i - a_j : a_i, a_j \in A\}$ .

### Definition

We say  $A$  is **difference dominated** if

$|A - A| > |A + A|$ , **balanced** if  $|A - A| = |A + A|$

and **sum dominated (or an MSTD set)** if

$|A + A| > |A - A|$ .

## Questions

Expect **generic** set to be difference dominated:

- addition is commutative, subtraction isn't:
- Generic pair  $(x, y)$  gives 1 sum, 2 differences.

## Questions

- Do there exist sum-dominated sets?
- If yes, how many?

## Examples

- Conway:  $\{0, 2, 3, 4, 7, 11, 12, 14\}$ .
- Marica (1969):  $\{0, 1, 2, 4, 7, 8, 12, 14, 15\}$ .
- Freiman and Pigarev (1973):  $\{0, 1, 2, 4, 5, 9, 12, 13, 14, 16, 17, 21, 24, 25, 26, 28, 29\}$ .
- Computer search: random subsets of  $\{1, \dots, 100\}$ :  
 $\{2, 6, 7, 9, 13, 14, 16, 18, 19, 22, 23, 25, 30, 31, 33, 37, 39, 41, 42, 45, 46, 47, 48, 49, 51, 52, 54, 57, 58, 59, 61, 64, 65, 66, 67, 68, 72, 73, 74, 75, 81, 83, 84, 87, 88, 91, 93, 94, 95, 98, 100\}$ .

## Binomial model

### Binomial model, parameter $p(n)$

Each  $k \in \{0, \dots, n\}$  is in  $A$  with probability  $p(n)$ .

Consider uniform model ( $p(n) = 1/2$ ):

- Let  $A \in \{0, \dots, n\}$ . Most elements in  $\{0, \dots, 2n\}$  in  $A + A$  and in  $\{-n, \dots, n\}$  in  $A - A$ .
- $\mathbb{E}[|A + A|] = 2n - 11$ ,  $\mathbb{E}[|A - A|] = 2n - 7$ .

## Martin and O'Bryant '06

### Theorem

*A from  $\{0, \dots, N\}$  by binomial model with constant parameter  $p$  (so  $k \in A$  with probability  $p$ ). At least  $k_{\text{SD};p} 2^{N+1}$  subsets are sum dominated.*

- $k_{\text{SD};1/2} \geq 10^{-7}$ , expect about  $10^{-3}$ .
- Proof ( $p = 1/2$ ): Generically  
 $|A| = \frac{N}{2} + O(\sqrt{N})$ .
  - ◇ about  $\frac{N}{4} - \frac{|N-k|}{4}$  ways write  $k \in A + A$ .
  - ◇ about  $\frac{N}{4} - \frac{|k|}{4}$  ways write  $k \in A - A$ .

## Notation

- $X \sim f(N)$  means  $\forall \epsilon_1, \epsilon_2 > 0, \exists N_{\epsilon_1, \epsilon_2}$  st  $\forall N \geq N_{\epsilon_1, \epsilon_2}$

$$\text{Prob}(X \notin [(1 - \epsilon_1)f(N), (1 + \epsilon_1)f(N)]) < \epsilon_2.$$

- $\mathcal{S} = |A + A|, \mathcal{D} = |A - A|,$   
 $\mathcal{S}^c = 2N + 1 - \mathcal{S}, \mathcal{D}^c = 2N + 1 - \mathcal{D}.$

New model: Binomial with parameter  $p(N)$ :

- $1/N = o(p(N))$  and  $p(N) = o(1)$ ;
- $\text{Prob}(k \in A) = p(N).$

### Conjecture (Martin-O'Bryant)

As  $N \rightarrow \infty$ ,  $A$  is a.s. difference dominated.

## Main Result

### Theorem (Hegarty-Miller)

$p(N)$  as above,  $g(x) = 2 \frac{e^{-x} - (1-x)}{x}$ .

- $p(N) = o(N^{-1/2})$ :  $\mathcal{D} \sim 2\mathcal{S} \sim (Np(N))^2$ ;
- $p(N) = cN^{-1/2}$ :  $\mathcal{D} \sim g(c^2)N$ ,  $\mathcal{S} \sim g\left(\frac{c^2}{2}\right)N$   
( $c \rightarrow 0$ ,  $\mathcal{D}/\mathcal{S} \rightarrow 2$ ;  $c \rightarrow \infty$ ,  $\mathcal{D}/\mathcal{S} \rightarrow 1$ );
- $N^{-1/2} = o(p(N))$ :  $\mathcal{S}^c \sim 2\mathcal{D}^c \sim 4/p(N)^2$ .

Can generalize to binary linear forms, still have  
critical threshold.

# Inputs

Key input: recent strong concentration results of Kim and Vu  
(Applications: combinatorial number theory, random graphs, ...).

**Example (Chernoff):**  $t_i$  iid binary random variables,  $Y = \sum_{i=1}^n t_i$ , then

$$\forall \lambda > 0 : \text{Prob} \left( |Y - \mathbb{E}[Y]| \geq \sqrt{\lambda n} \right) \leq 2e^{-\lambda/2}.$$

Need to allow dependent random variables.

Sketch of proofs:  $\mathcal{X} \in \{\mathcal{S}, \mathcal{D}, \mathcal{S}^c, \mathcal{D}^c\}$ .

- 1 Prove  $\mathbb{E}[\mathcal{X}]$  behaves asymptotically as claimed;
- 2 Prove  $\mathcal{X}$  is strongly concentrated about mean.

## Setup for Proofs

Note: only need strong concentration for  $N^{-1/2} = o(p(N))$ .

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Will assume  $p(N) = o(N^{-1/2})$  as proofs are elementary (i.e., Chebyshev:  $\text{Prob}(|Y - \mathbb{E}[Y]| \geq k\sigma_Y) \leq 1/k^2$ ).

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For convenience let  $p(N) = N^{-\delta}$ ,  $\delta \in (1/2, 1)$ .

i.i.d binary indicator variables:

$$X_{n;N} = \begin{cases} 1 & \text{with probability } N^{-\delta} \\ 0 & \text{with probability } 1 - N^{-\delta}. \end{cases}$$

$$X = \sum_{i=1}^N X_{n;N}, \quad \mathbb{E}[X] = N^{1-\delta}.$$

# Proof

## Lemma

$P_1(N) = 4N^{-(1-\delta)}$ ,  $\mathcal{O} = \#\{(m, n) : m < n \in \{1, \dots, N\} \cap A\}$ .

With probability at least  $1 - P_1(N)$  have

- 1  $X \in [\frac{1}{2}N^{1-\delta}, \frac{3}{2}N^{1-\delta}]$ .

- 2  $\frac{\frac{1}{2}N^{1-\delta}(\frac{1}{2}N^{1-\delta}-1)}{2} \leq \mathcal{O} \leq \frac{\frac{3}{2}N^{1-\delta}(\frac{3}{2}N^{1-\delta}-1)}{2}$ .

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*With probability at least  $1 - P_1(N)$  have*

①  $X \in [\frac{1}{2}N^{1-\delta}, \frac{3}{2}N^{1-\delta}]$ .

②  $\frac{\frac{1}{2}N^{1-\delta}(\frac{1}{2}N^{1-\delta}-1)}{2} \leq \mathcal{O} \leq \frac{\frac{3}{2}N^{1-\delta}(\frac{3}{2}N^{1-\delta}-1)}{2}$ .

Proof:

- (1) is Chebyshev:  $\text{Var}(X) = N\text{Var}(X_{n,N}) \leq N^{1-\delta}$ .
- (2) follows from (1) and  $\binom{r}{2}$  ways to choose 2 from  $r$ .

## Concentration

### Lemma

- $f(\delta) = \min\left(\frac{1}{2}, \frac{3\delta-1}{2}\right)$ ,  $g(\delta)$  any function st  $0 < g(\delta) < f(\delta)$ .
- $p(N) = N^{-\delta}$ ,  $\delta \in (1/2, 1)$ ,  $P_1(N) = 4N^{-(1-\delta)}$ ,  
 $P_2(N) = CN^{-(f(\delta)-g(\delta))}$ .

With probability at least  $1 - P_1(N) - P_2(N)$  have  
 $\mathcal{D}/S = 2 + O(N^{-g(\delta)})$ .

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With probability at least  $1 - P_1(N) - P_2(N)$  have  
 $\mathcal{D}/\mathcal{S} = 2 + O(N^{-g(\delta)})$ .

Proof: Show  $\mathcal{D} \sim 2\mathcal{O} + O(N^{3-4\delta})$ ,  $\mathcal{S} \sim \mathcal{O} + O(N^{3-4\delta})$ .

As  $\mathcal{O}$  is of size  $N^{2-2\delta}$  with high probability, need  $2 - 2\delta > 3 - 4\delta$  or  
 $\delta > 1/2$ .

## Analysis of $\mathcal{D}$

Contribution from 'diagonal' terms lower order, ignore.

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$$Y_{m,n,m',n'} = \begin{cases} 1 & \text{if } n - m = n' - m' \\ 0 & \text{otherwise.} \end{cases}$$

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$$Y_{m,n,m',n'} = \begin{cases} 1 & \text{if } n - m = n' - m' \\ 0 & \text{otherwise.} \end{cases}$$

$\mathbb{E}[Y] \leq N^3 \cdot N^{-4\delta} + N^2 \cdot N^{-3\delta} \leq 2N^{3-4\delta}$ . As  $\delta > 1/2$ ,  
Expected number bad pairs  $\lll |\mathcal{O}|$ .

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Expected number bad pairs  $\lll |\mathcal{O}|$ .

**Claim:**  $\sigma_Y \leq N^{r(\delta)}$  with  $r(\delta) = \frac{1}{2} \max(3 - 4\delta, 5 - 7\delta)$ . This and Chebyshev conclude proof of theorem.

## Proof of claim

Cannot use CLT as  $Y_{m,n,m',n'}$  are not independent.

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Use  $\text{Var}(U + V) \leq 2\text{Var}(U) + 2\text{Var}(V)$ .

Write

$$\sum Y_{m,n,m',n'} = \sum U_{m,n,m',n'} + \sum V_{m,n,n'}$$

with all indices distinct (at most one in common, if so must be  $n = m'$ ).

$$\text{Var}(U) = \sum \text{Var}(U_{m,n,m',n'}) + 2 \sum_{\substack{(m,n,m',n') \neq \\ (\tilde{m}, \tilde{n}, \tilde{m}', \tilde{n}')}} \text{CoVar}(U_{m,n,m',n'}, U_{\tilde{m}, \tilde{n}, \tilde{m}', \tilde{n}'}).$$

## Analyzing $\text{Var}(U_{m,n,m',n'})$

At most  $N^3$  tuples.

Each has variance  $N^{-4\delta} - N^{-8\delta} \leq N^{-4\delta}$ .

Thus  $\sum \text{Var}(U_{m,n,m',n'}) \leq N^{3-4\delta}$ .

## Analyzing $\text{CoVar}(U_{m,n,m',n'}, U_{\tilde{m},\tilde{n},\tilde{m}',\tilde{n}'})$

- All 8 indices distinct: independent, covariance of 0.
- 7 indices distinct: At most  $N^3$  choices for first tuple, at most  $N^2$  for second, get

$$\mathbb{E}[U_{(1)}U_{(2)}] - \mathbb{E}[U_{(1)}]\mathbb{E}[U_{(2)}] = N^{-7\delta} - N^{-4\delta}N^{-4\delta} \leq N^{-7\delta}.$$

- Argue similarly for rest, get  $\ll N^{5-7\delta} + N^{3-4\delta}$ .