

On Summand Minimality of Generalized Zeckendorf Decompositions

Chi Huynh (nhuynh30@gatech.edu)

Carsten Peterson (carsten.peterson@yale.edu)

Yen Nhi Truong Vu (ytruongvu17@amherst.edu)

Joint work with Katherine Cordwell, Max Hlavacek and
Steven J. Miller

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Introduction

Base d Representation and Zeckendorf's theorem

Theorem (Base d Representation)

Every natural number can be represented uniquely as a sum of powers of d with coefficients in $\{0, 1, \dots, d - 1\}$.

Theorem (Zeckendorf)

Every natural number can be represented uniquely as a sum of non-consecutive Fibonacci numbers (with $F_1 = 1$ and $F_2 = 2$).

Definition: PLRS

Definition

A **positive linear recurrence sequence (PLRS)** is the sequence given by a recurrence of the following form:

$$H_n = c_1 H_{n-1} + \cdots + c_t H_{n-t}$$

where all $c_i \geq 0$ and $c_1, c_t \geq 1$. We use **ideal initial conditions** $H_{-(n-1)} = 0, \dots, H_{-1} = 0, H_0 = 1$.

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Definition

We call (c_1, \dots, c_t) the **signature of the sequence**.

Representation Notation

Definition

For a sequence $\{H_i\}$ with signature (c_1, \dots, c_t) , suppose we have:

$$n = b_k H_k + b_{k-1} H_{k-1} + \dots + b_0 H_0$$

Then we call $[b_k, b_{k-1}, \dots, b_0, \infty_{(t-1)}]$ a **representation** of n over $\{H_i\}$.

Example: Representation Notation

Representations of 18

F_7	F_6	F_5	F_4	F_3	F_2	F_1	F_0	F_{-1}
21	13	8	5	3	2	1	1	0
<hr/>								
	1	0	1	0	0	0	0	∞
	1	0	0	1	1	0	0	∞

Allowable Blocks

Definition

Given a PLRS with signature (c_1, \dots, c_t) , we say that $[b_1, \dots, b_k]$ is an **allowable block** if $k \leq t$ and $b_i = c_i$ for $i < k$ and $0 \leq b_k < c_k$. In other words, allowable blocks are *lexicographically less than* the signature.

Examples:

$$\sigma = (1, 1)$$

$$[0]$$

$$[1, 0]$$

$$\sigma = (2, 2, 1)$$

$$[0]$$

$$[1]$$

$$[2, 0]$$

$$[2, 1]$$

$$[2, 2, 0]$$

$$\sigma = (3, 0, 2)$$

$$[0]$$

$$[1]$$

$$[2]$$

$$[3, 0, 0]$$

$$[3, 0, 1]$$

Generalized Zeckendorf Decompositions

Theorem (Miller et al., Hamlin)

*Given a positive linear recurrence sequence, every natural number has a unique representation composed of allowable blocks. This representation is called the **Generalized Zeckendorf Decomposition** (GZD).*

Examples

Signature: $\sigma = (d)$

Allowable blocks: $\{[0], [1], \dots, [d-1]\}$

$$\begin{array}{r} 1000 \quad 100 \quad 10 \quad 1 \\ \hline \quad 7 \quad \quad 3 \quad \quad 0 \quad 9 \end{array}$$

This is the base d representation.

Signature: $\sigma = (1, 1)$

Allowable blocks: $\{[0], [1, 0]\}$

$$\begin{array}{r} 5 \quad 3 \quad 2 \quad 1 \quad 1 \quad 0 \\ \hline 1 \quad 0 \quad 0 \quad 1 \quad 0 \quad \infty \end{array}$$

This is the Zeckendorf representation.

Example

Signature: $\sigma = (2, 2, 1)$

Allowable blocks: $\{[0], [1], [2, 0], [2, 1], [2, 2, 0]\}$

385	136	48	17	6	2	1	0	0
2	0	1	2	2	0	0	∞	∞

Summand minimality

Number of Summands

Signature: $\sigma = (1, 1)$

Representations of 6:

# of summands	5	3	2	1	1	0
2	1	0	0	1	0	∞
2		2	0	0	0	∞
3		1	1	1	0	∞
3			3	0	0	∞
4			2	2	0	∞
5			1	0	4	∞
6					6	∞

The GZD has the fewest number of summands.

Number of Summands

Signature: $\sigma = (3, 0, 2)$

Representations of 116:

# of summands	93	29	9	3	1	0	0
4		4	0	0	0	∞	∞
6	1	0	2	1	2	∞	∞

There exists a representation with fewer summands than the GZD.

Summand minimality

Definition

A positive linear recurrence sequence, H , is called **summand minimal** if for all $n \in \mathbb{N}$, the GZD of n has the fewest number of summands among all representations for n using H .

Question

When is a positive linear recurrence summand minimal?

Summand minimality

Definition

A positive linear recurrence sequence, H , is called **summand minimal** if for all $n \in \mathbb{N}$, the GZD of n has the fewest number of summands among all representations for n using H .

Question

When is a positive linear recurrence summand minimal?

Theorem (CHHMPT '16)

A positive linear recurrence sequence is summand minimal if and only if its signature is weakly decreasing, i.e. $\sigma = (c_1, c_2, \dots, c_t)$ with $c_1 \geq c_2 \geq \dots \geq c_t$.

Sketch of techniques

From one representation to another: borrowing

- Signature $\sigma = (3, 2)$
- Strings $[-1, 3, 2]$ and $[1, -3, -2]$ represent 0.
- Representations of 39

H_4	H_3	H_2	H_1	H_0	H_{-1}
139	39	11	3	1	0
<hr/>					
	1	0	0	0	∞
	-1	3	2		
<hr/>					
	0	3	2	0	∞
			-1	3	2
<hr/>					
		3	1	3	∞

From one representation to another: carrying

- Signature $\sigma = (2, 1, 3)$
- Representations of 42

H_4	H_3	H_2	H_1	H_0	H_{-1}	H_{-2}
41	15	5	2	1	0	0
<hr/>						
	1	3	4	4	∞	∞
	1	-2	-1	-3		
<hr/>						
	2	1	3	1	∞	∞
1	-2	-1	-3			
<hr/>						
1	0	0	0	1	∞	∞

From any representation to the GZD

Theorem (CHHMPT '16)

There exists an algorithm which, given any $n \in \mathbb{N}$ and any representation for n , terminates in the GZD for n through a finite sequence of borrows and carries.

Example of the algorithm

- Signature $\sigma = (3, 2, 2)$
- Representations of 44

H_5	H_4	H_3	H_2	H_1	H_0	H_{-1}	H_{-2}
537	148	41	11	3	1	0	0
			4	0	0	∞	∞
			-1	3	2	2	
			3	3	2	∞	∞
		1	-3	-2	-2		
		1	0	1	0	∞	∞

Weakly decreasing signature implies summand minimality

Corollary

The definition of the algorithm immediately implies that positive linear recurrences sequences with weakly decreasing signature are summand minimal.

Idea: after every borrow we are immediately able to carry which causes the number of summands to weakly decrease.

Summand minimality implies weakly decreasing signature

- What about non-weakly decreasing signature?
- When $c_1 > 1$, we can construct an explicit non-GZD representation with fewer summands than the GZD. Idea is to construct a representation such that during the algorithm we need at least two borrows before we can carry (so that the number of summands increases as we move towards the GZD).
- When $c_1 = 1$, using the theory of irreducible polynomials, we are able to non-constructively prove the existence of a non-GZD representation with fewer summands than the GZD.

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Acknowledgments

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<https://arxiv.org/abs/1608.08764>
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