

Distribution of Missing Sums in Sumsets

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Background

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- Goldbach's conjecture: $\{4, 6, 8, \dots\} \subseteq P + P$.
- Fermat's last theorem: let A_n be the n th powers and then ask if $(A_n + A_n) \cap A_n = \emptyset$ for all $n > 2$.

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Key Question: What is the structure of $A + A$?

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Note: Both theorems can be more naturally stated in terms of missing sums (independent of n).

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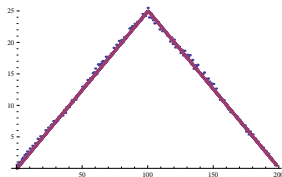


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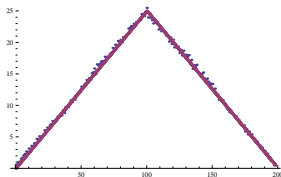


Figure: Comparison of predicted and observed number of representations of possible elements of the sumset.

- Key fact:** if $k < n$, then $P(k \notin A + A) \sim \left(\frac{3}{4}\right)^{k/2}$.

Results

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Theorem: Bounds on the distribution (Lazarev-Miller, 2011)

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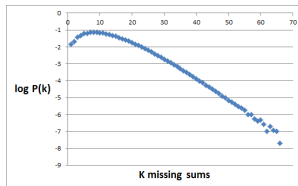


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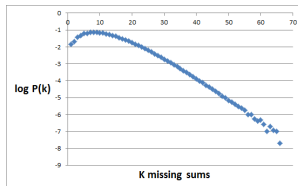


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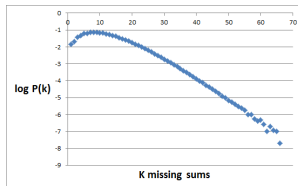


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Main idea: Use graph theory.

More Results

Theorem: Variance (Lazarev-Miller)

$$\text{Var}|A + A| = 4 \sum_{i < j \leq n-1} P(i \text{ and } j \notin A + A) - 40 \sim 35.98.$$

Theorem: Distribution of configurations (Lazarev-Miller)

For any fixed a_1, \dots, a_m , exists $\lambda_{a_1, \dots, a_m}$ such that

$$P(k + a_1, k + a_2, \dots, \text{ and } k + a_m \notin A + A) = \Theta(\lambda_{a_1, \dots, a_m}^k).$$

Theorem: Consecutive missing sums (Lazarev-Miller)

$$P(k, k + 1, \dots, \text{ and } k + m \notin A + A) = \left(\frac{1}{2} + o(1) \right)^{(k+m)/2}.$$

Bounds on the Distribution

Bound on Distribution: Lower Bound

Lower bound: $P(A + A \text{ has } k \text{ missing sums}) > 0.01 \cdot 0.70^k$

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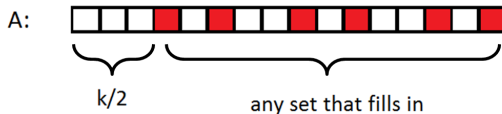
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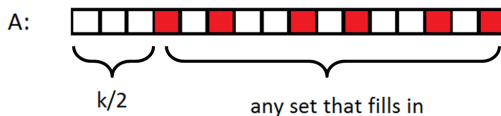


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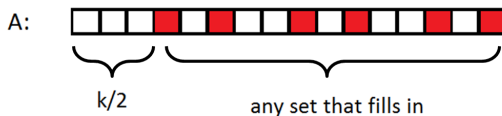


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$$P(A + A \text{ has } k \text{ missing sums}) > 0.01 \left(\frac{1}{2}\right)^{k/2} \sim 0.01 \cdot 0.70^k.$$

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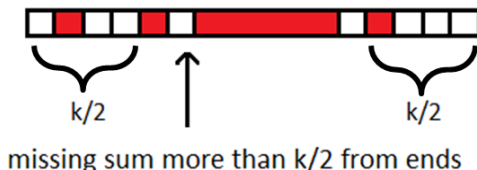
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- $P(A + A \text{ has } k \text{ missing sums}) < P(k/2 \notin A + A) < \left(\frac{3}{4}\right)^{k/4} \sim 0.93^k$.

Note: Bounds on $P(k + a_1, k + a_2, \dots, \text{ and } k + a_m \notin A + A)$ yield upper bounds on $P(A + A \text{ has } k \text{ missing sums})$.

Variance

Problem: Dependent Random Variables

Variances reduces to $\sum_{0 \leq i, j \leq 2n-2} \mathbf{P}(A : i \text{ and } j \notin A + A)$.

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Example: $P(A : 3 \text{ and } 7 \notin A + A)$

● Conditions:

$i = 3 :$ 0 or 3 $\notin A$
and 1 or 2 $\notin A$

$j = 7 :$ 0 or 7 $\notin A$
and 1 or 6 $\notin A$
and 2 or 5 $\notin A$
and 3 or 4 $\notin A$.

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$j = 7$: $0 \text{ or } 7 \notin A$
and $1 \text{ or } 6 \notin A$
and $2 \text{ or } 5 \notin A$
and $3 \text{ or } 4 \notin A$.

- Since there are common integers in both lists, the events $3 \notin A + A$ and $7 \notin A + A$ are dependent.

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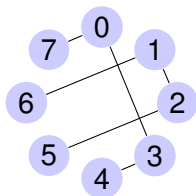
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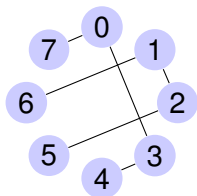
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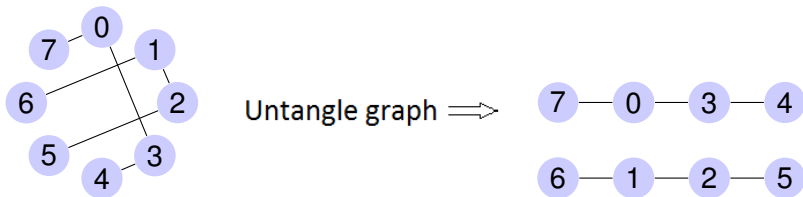


Untangle graph \Rightarrow

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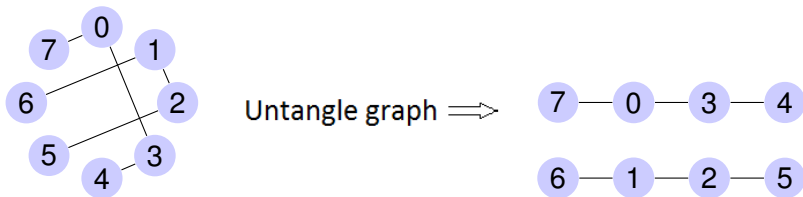
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Example $i = 3, j = 7$:



- One-to-one correspondence between conditions/edges (and integers/vertices).

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- So need to pick a **vertex cover!**

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Lemma (Lazarev-Miller)

$$P(i, j \notin A + A) = P(\text{pick a vertex cover for graph}).$$

Number of Vertex Covers

Condition graphs are always ‘segment’ graphs. So we just need $g(n)$, the number of vertex covers for a ‘segment’ graph with n vertices.

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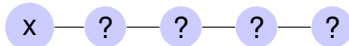
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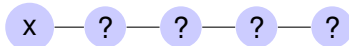
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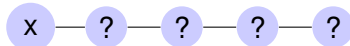


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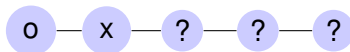
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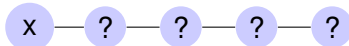


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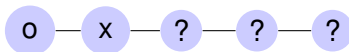
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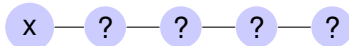
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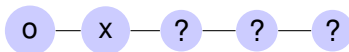
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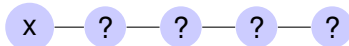
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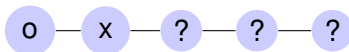
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$$\implies g(n) = F_{n+2}$$

General i, j

- In particular

$$P(3 \text{ and } 7 \notin A + A) = \frac{1}{2^8} F_{4+2} F_{4+2} = \frac{1}{4}$$

since there were two graphs each of length 4.

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$$\begin{aligned} &P(A : i \text{ and } j \notin A + A) \\ &= \frac{1}{2^{j+1}} F_{\frac{1}{2} \left((j-i) \left\lceil \frac{i+1}{j-i} \right\rceil - (i+1) \right)} \times F_{\frac{1}{2} \left(j+1 - (j-i) \left\lceil \frac{i+1}{j-i} \right\rceil \right)} \\ &\quad \times F_{2 \left\lceil \frac{i+1}{j-i} \right\rceil + 2}. \end{aligned}$$

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$$\begin{aligned} &P(A : i \text{ and } j \notin A + A) \\ &= \frac{1}{2^{j+1}} F_{\frac{1}{2} \left((j-i) \left\lceil \frac{i+1}{j-i} \right\rceil - (i+1) \right)} \times F_{\frac{1}{2} \left(j+1 - (j-i) \left\lceil \frac{i+1}{j-i} \right\rceil \right)} \\ &\quad \times F_{2 \left\lceil \frac{i+1}{j-i} \right\rceil + 2}. \end{aligned}$$

- In general $P(k \text{ and } k+1 \notin A + A) < C(\phi/2)^k \sim 0.81^k$, giving upper bound.

Variance Formula

$$\begin{aligned}
 \text{Var}|A + A| &= -40 + 4 \sum_{i < j < n} P(i, j \notin A + A) \\
 &= -40 + O(c^n) \\
 &+ 4 \sum_{i, j \text{ odd}} \frac{1}{2^{j+1}} F_{2^{\lceil \frac{j+1}{j-i} \rceil + 2}}^{\frac{1}{2} \left((j-i) \lceil \frac{j+1}{j-i} \rceil - (i+1) \right)} F_{2^{\lceil \frac{j+1}{j-i} \rceil + 4}}^{\frac{1}{2} \left((j+1 - (j-i)) \lceil \frac{j+1}{j-i} \rceil \right)} \\
 &+ 4 \sum_{i \text{ even}, j \text{ odd}} \frac{1}{2^{j+1}} F_{2^{\lceil \frac{j+1}{j-i} \rceil + 1}}^{\frac{1}{2} \left((j-i-1) \lceil \frac{j+1}{j-i} \rceil - (i+1) + 2 \lceil \frac{j+1}{j-i} \rceil - 1 \right)} F_{2^{\lceil \frac{j+1}{j-i} \rceil + 4}}^{\frac{1}{2} \left((j+1 - (j-i-1)) \lceil \frac{j+1}{j-i} \rceil - 2 \lceil \frac{j+1}{j-i} \rceil \right)} \\
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 &+ 4 \sum_{i, j \text{ even}} \frac{1}{2^{j+1}} F_{2^{\lceil \frac{j+1}{j-i} \rceil + 1}}^{\frac{1}{2} \left((j-i-2) \lceil \frac{j+1}{j-i} \rceil - (i+1) + 2 \lceil \frac{j+1}{j-i} \rceil + 2 \lceil \frac{j+1}{j-i} \rceil - 3 \right)} F_{2^{\lceil \frac{j+1}{j-i} \rceil + 4}}^{\frac{1}{2} \left((j+2 - (j-i-2)) \lceil \frac{j+1}{j-i} \rceil + 2 \lceil \frac{j+1}{j-i} \rceil + 2 \lceil \frac{j+1}{j-i} \rceil \right)}
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 \end{aligned}$$

So clearly

$$\text{Var}|A + A| \sim 35.98.$$

Consecutive Missing Sums

Consecutive Missing Sums in $A+A$

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Consecutive Missing Sums in $A+A$

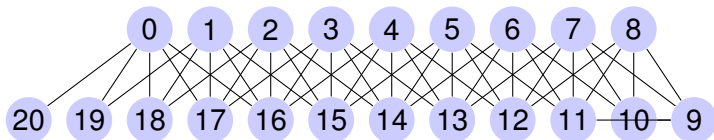
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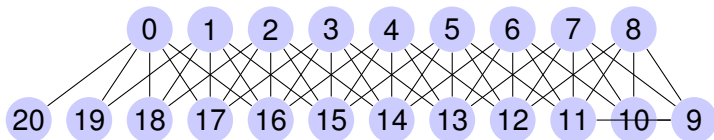
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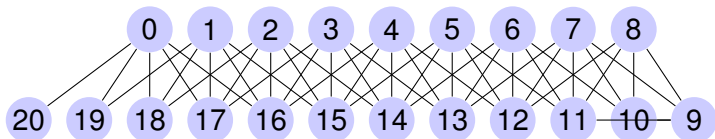


\implies Transforms to:

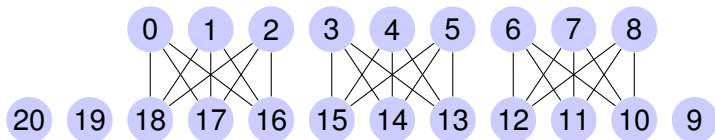
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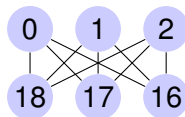


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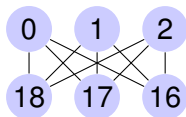
Consecutive Missing Sums in A+A

- So have 3 **complete bipartite** graphs like:



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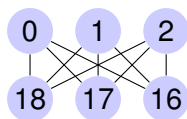
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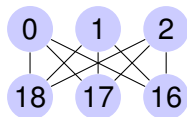
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- By independence**, $P(16, 17, 18, 19, 20) \leq \left(\frac{1}{4}\right)^3 \sim \left(\frac{1}{4}\right)^{20/6}$.
- In general,
$$P(k, k+1, k+2, k+3, k+4) \leq \left(\frac{1}{4}\right)^{(k+4)/6} \sim 0.79^{k+4}.$$

Consecutive Missing Sums

- Most general case is:

$$P(k, k+1, \dots, k+i \notin A+A) \leq \left(\frac{1}{2}\right)^{(k+i)/2} (1 + \epsilon_i)^k.$$

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- **Why interesting?** Bounds almost match!
- Essentially the only way to miss a block of i consecutive sums is to miss all elements before the block as well.

Summary

Use graph theory to study $P(a_1, \dots, \text{ and } a_m \notin A + A)$.

Currently investigating:

- Is distribution of missing sums approximately exponential?
- Higher moments: third moment involves $P(i, j, k \notin A + A)$, with more complicated graphs.
- Distribution of $A - A$.

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Use graph theory to study $P(a_1, \dots, \text{ and } a_m \notin A + A)$.

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Thank you!

CANT 2012



Williams

