

## LECTURE 15: SUMMARY

In today's lecture, we proved the following result (which is half of David's conjecture from last lecture):

**Theorem 1.** *If  $(m, n) = 1$ , then  $\varphi(mn) = \varphi(m)\varphi(n)$ .*

Right off the bat, note that the hypothesis that  $m$  and  $n$  are relatively prime is necessary. For example,  $\varphi(12) \neq \varphi(2)\varphi(6)$ . We also practiced using this theorem to calculate  $\varphi(n)$ . As we saw, whenever we could factor  $n$ , the theorem made it easy to figure out  $\varphi(n)$ . Unfortunately, if  $n$  is not easy to factor, then it's less clear how to determine  $\varphi(n)$ . We will discuss this in more depth later, when talking about the RSA encryption algorithm.

Before writing down the proof of theorem, we discuss the strategy. By definition, we have

$$\varphi(mn) = |\mathbb{Z}_{mn}^\times|.$$

What about  $\varphi(m)\varphi(n)$ ? A bit of thought showed that this, too, measures the size of a set:

$$\varphi(m)\varphi(n) = |\mathbb{Z}_m^\times \times \mathbb{Z}_n^\times|$$

where  $A \times B := \{(a, b) : a \in A, b \in B\}$ . Thus, if we can show that the two sets  $\mathbb{Z}_{mn}^\times$  and  $\mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$  have the same size, we win. How will we do this? We look for a bijection between the two sets, i.e. a way of pairing off elements of the two sets. Shichu suggested the following function:

$$\begin{aligned}\sigma : \mathbb{Z}_{mn}^\times &\longrightarrow \mathbb{Z}_m^\times \times \mathbb{Z}_n^\times \\ a &\longmapsto (a \pmod{m}, a \pmod{n})\end{aligned}$$

where  $x \pmod{d}$  denotes the unique element of  $\mathbb{Z}_d$  which is congruent to  $x$  modulo  $d$ . If we can show that this is a bijection – i.e. that for every  $(x, y) \in \mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$ , there exists a unique  $a \in \mathbb{Z}_{mn}^\times$  such that  $\sigma(a) = (x, y)$  – then it would immediately follow that  $\mathbb{Z}_{mn}^\times$  and  $\mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$  have the same number of elements.

Before going into the proof of the theorem, we state a useful tool:

**Lemma 2.** *Suppose  $(a, N) = 1$ . Then the integer  $a \pmod{N}$  is also relatively prime to  $N$ , i.e.  $a \pmod{N} \in \mathbb{Z}_N^\times$ .*

I leave the proof of this lemma as an exercise.

*Proof.* Consider the function  $\sigma$  defined above. We prove that it's a bijection in three steps:

- (1)  $\sigma$  is **well-defined**, i.e. for all  $x \in \mathbb{Z}_{mn}^\times$  there exists a unique  $(a, b) \in \mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$  such that  $\sigma(x) = (a, b)$ ;

- (2)  $\sigma$  is **surjective**, i.e. for all  $(a, b) \in \mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$  there exists at least one  $x \in \mathbb{Z}_{mn}^\times$  such that  $\sigma(x) = (a, b)$ ; and
- (3)  $\sigma$  is **injective**, i.e. for all  $(a, b) \in \mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$  there exists at most one  $x \in \mathbb{Z}_{mn}^\times$  such that  $\sigma(x) = (a, b)$ .

First, why is  $\sigma$  well-defined? Well, certainly  $\sigma(x) \in \mathbb{Z}_m \times \mathbb{Z}_n$ ; what's not immediate is that  $\sigma(x) \in \mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$ . However, armed with the Lemma above this isn't so difficult. Since  $x \in \mathbb{Z}_{mn}^\times$ , we know that  $(x, mn) = 1$ . It follows that  $(x, m) = 1$ , whence (by the lemma) the integer  $x \pmod{m} \in \mathbb{Z}_m^\times$ . The same goes for  $x \pmod{n}$ , of course.

Next, why is  $\sigma$  surjective? Given  $(a, b) \in \mathbb{Z}_m^\times \times \mathbb{Z}_n^\times$ , can we find an  $x \in \mathbb{Z}_{mn}^\times$  such that  $\sigma(x) = (a, b)$ ? It's easy to see that this is equivalent to finding an  $x \in \mathbb{Z}_{mn}^\times$  such that

$$x \equiv a \pmod{m} \quad \text{and} \quad x \equiv b \pmod{n}.$$

The trick is to write  $x = (\cdots)m + (\cdots)n$ , and find appropriate ways to fill in the blanks. The advantage of writing  $x$  this way is that when we reduce  $x \pmod{m}$  we can focus on just the second term, while when we reduce  $\pmod{n}$  we can focus on just the first term. A bit of thought showed that we should choose the first blank to be  $bm^{-1}$ , where  $m^{-1}$  denotes the inverse of  $m$  in  $\mathbb{Z}_n^\times$ , and the second blank to be  $an^{-1}$ , where  $n^{-1}$  denotes the inverse of  $n$  in  $\mathbb{Z}_m^\times$ .<sup>1</sup> In any event, let

$$x = (bm^{-1})m + (an^{-1})n.$$

It's easy to check that  $x \pmod{m} = a$  and  $x \pmod{n} = b$ . The only remaining difficulty is that  $x$  is just some integer; it might not be an element of  $\mathbb{Z}_{mn}^\times$ ! Fortunately, this can be fixed. I leave this as an exercise.

Finally, why is  $\sigma$  injective? Well, suppose  $\sigma(x) = \sigma(y)$  for some  $x, y \in \mathbb{Z}_{mn}^\times$ . Then

$$x \equiv y \pmod{m} \quad \text{and} \quad x \equiv y \pmod{n}.$$

It follows that  $m \mid x - y$  and also  $n \mid x - y$ . By problem 1.9 from your homework, it follows that  $mn \mid x - y$ , i.e. that  $x \equiv y \pmod{mn}$ . Thus,  $x = y$ , so  $\sigma$  is injective.  $\square$

Make sure that you go through and understand the theorem properly; there were some gaps in the sketch above. Among other questions, you should ask yourself: where did we use that  $(m, n) = 1$ ?

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<sup>1</sup>Actually, if we were being super careful, we should be referring to  $n \pmod{m}$  and  $m \pmod{n}$  in the previous sentence, rather than to  $n$  and  $m$  themselves.