

# THE CIRCLE METHOD

PROBLEMS: Waring: Is there an  $s$  st all pos ints (or all suff large pos ints) are a sum of at most  $s$   $k$ -th powers?

$$\hookrightarrow x_1^k + \dots + x_s^k = n$$

• Goldbach:  $P_1 + P_2 + P_3 = N$

$$P_1 + P_2 = N$$

• Goldbach:  $P_1 + P_2 = N$

• Goldbach:  $P$  st  $P, \frac{P-1}{2}$  prime or  $P, P+2$  prime

How to solve? Especially Goldbach

$\hookrightarrow$  primes are multiplicative, these sunset questions

## NEEDED INPUT: Number Theory

See section 13.2.6

PRIME NUMBER THEM (PNT):  $\sum_{p \leq x} \log p = x + O(x \exp(-c\sqrt{\log x}))$

for some  $c > 1$ . By partial sum,  $\sum_{p \leq x} 1 = \text{Li}(x) + O(x \exp(-\frac{c\sqrt{\log x}}{2}))$

$\hookrightarrow$  See Section 2.2.2

$$\hookrightarrow \text{Li}(x) \approx \frac{x}{\log x}$$

SIEGEL-WALFISZ:  $\sum_{\substack{p \leq x \\ p \equiv a \pmod{q}}} \log p = \frac{x}{\phi(q)} + O\left(\frac{x}{\log^c x}\right) \forall q \leq \log^B x$

$\forall C, B > 0$  and  $(a, q) = 1$   
 $\hookrightarrow$  wish could do  $q \leq x^C$

$\hookrightarrow$  Main term  $\gg$  error term if  $C$  suff large

$\hookrightarrow \phi(q)$  Euler's  $\phi$ -function; See Section 2.1

Need basic relation from Fourier Analysis:

$$\int_0^1 \mathbb{E}(nx) \mathbb{E}(-mx) dx = \begin{cases} 1 & n=m \\ 0 & \text{otherwise} \end{cases}$$

# ABIDE: PNT and SIEGEL-WALFISZ

## Counting Primes

① Euclid:  $P_1 \cdots P_n + 1$

↳ investigate on OEIS where  $q_l$  is  $l^{\text{th}}$  prime generated by this

↳ can generalize to some but not all arithmetic prog  
 $p \equiv 4n+1$  no     $p \equiv 4n+3$  yes (see Exe 2.3.5)

② See "Proofs from THE BOOK"

↳ Two favorites:  $\zeta(s) = \sum_{n=1}^{\infty} \frac{1}{n^s} = \prod_{p \text{ prime}} (1 - \frac{1}{p^s})^{-1}$   $\text{Re}(s) > 1$

$\lim_{s \rightarrow 1^+} \zeta(s) = +\infty \Rightarrow \infty$  many primes

$\zeta(2) = \pi^2/6$  and  $\pi^2 \notin \mathbb{Q} \Rightarrow \infty$  many primes (see also § 11.3.4)

③ Chebyshev's Thm (Thm 2.3.9)  $\exists A, B$  with  $0 < A < B < \infty$  st

$$\forall x \geq 30 \quad \frac{Ax}{\log x} \leq \pi(x) \leq \frac{Bx}{\log x}$$

↳ many proofs, one in book uses nice approx technique and telescoping sums

↳ corr: Bertrand's Postulate:  $\forall n \geq 1 \exists p \in [n, 2n]$ !

Questions on primes in short gaps

Good Exercise: Exe 2.3.15 (Prime Deserts)

↳ also Exe 2.3.16, 2.3.17

④ Dirichlet's Thm (Section 3.3):  $(a, m) = 1$ . Then

$$\pi_{m,a}(x) = \frac{1}{\phi(m)} \frac{x}{\log x} + O_a\left(\frac{1}{\phi(m)} \frac{x}{\log x}\right)$$

↳ Siegel-Walfisz is controlling error

# THE CIRCLE METHOD

## ORIGINS / BASIC PROBLEMS

"Cookie Problem":  $N$  cookies,  $p$  people:  $\binom{N+p-1}{p-1}$

↳ equiv to  $x_1 + \dots + x_p = N$

(note: easy add conds such as  $x_i \geq a_i$ )

↳ solving by combinatorics; doesn't generalize to high powers

such as  $x_1^k + \dots + x_s^k = N$

## Generating Functions

↳ look at  $x_1 + \dots + x_s = N$  again

$$f(x) = \sum_{m=0}^{\infty} x^m = \frac{1}{1-x}$$

$$f(x)^s = \left( \sum_{m_1} x^{m_1} \right) \dots \left( \sum_{m_s} x^{m_s} \right) = \sum_{n=0}^{\infty} r_{1,s}(n) x^n$$

↳ Thus  $r_{1,s}(n)$  is answer, but how to recover?

BUT  $f(x)^s = \left( \frac{1}{1-x} \right)^s = \frac{1}{(s-1)!} \frac{d^{s-1}}{dx^{s-1}} \frac{1}{1-x}$

$$\Rightarrow f(x)^s = \frac{1}{(s-1)!} \frac{d^{s-1}}{dx^{s-1}} \left( \sum_{n=0}^{\infty} x^n \right) = \sum_{n=0}^{\infty} \binom{n+s-1}{s-1} x^n$$

$$\Rightarrow r_{1,s}(n) = \binom{n+s-1}{s-1}$$

WARNING:  $f_{k,N}(x) = \sum_{n=0}^N x^{nk}$ ,  $f_{k,N}^s(x) = \dots$

Goldbach:  $f_N(x) = \sum_{p \leq N} \mathcal{O}(px)$ ,  $f_N^s(x) = \dots$   $\mathcal{O}(u) = e^{2\pi i u}$

↳ technically easier to do finite sums and exponentials

↳ get say  $r_{N,s}(m) = r_s(m)$  if  $m \leq N$

ASIDE: Differentiating identities **POWERFUL** technique, especially for calculating moments in prob / generating identities. See handout



# THE CIRCLE METHOD

Write  $n = p_1 + \dots + p_s$

$$F_N(x) = \sum_{p \leq N} \log p \cdot e(px)$$

↳  $\log p$  is to simplify some sums  
technically easier to count primes @  $\log$  weight

ASIDE: Section 2.3.4 and 3.2.2 for  
why we count primes with  $\log p$  weight

$$F_N^s(x) = \sum_{m=0}^{\infty} R_{N,s}(m) e(mx)$$

where  $R_{N,s}(m)$  is # ways write  $m$  as a sum  
of  $s$  primes at most  $N$ . Note  $R_{N,s}(m) = R_s(m)$   
for  $m \leq N$

## Basic Fourier analysis

$$R_{N,s}(n) = \langle F_N^s, e(-nx) \rangle = \int_0^1 F_N^s(x) e(-nx) dx$$

↳ Note: we have "solved" our problem!

Unfortunately, not very illuminating

Goal: Show  $R_{N,s}(N) > 1 \rightarrow$  at least one way

↳ analyzing that can often only prove one item exists  
by showing as many do

↳ ex: general  $a, q$  rel prime: to show  $\exists p \equiv a(q)$   
with  $p$  prime requires proving Dirichlet's  
Thm ( $\pi_{a,q}(x) \sim \frac{1}{\phi(q)} \pi(x)$ )

# THE CIRCLE METHOD: SIZE OF $M F_N(x)$

$$F_N(x) = \sum_{p \leq N} \log p \cdot e(p x)$$

Trivial:  $|F_N(x)| \leq \sum_{p \leq N} \log p = N + o(N)$

Average: On average  $|F_N(x)|^2$  is like  $N \log N + o(N \log N)$

↳ Proof: VERY IMPORTANT TECHNIQUE  
works for  $L^2$  norm, not others...

$$\int_0^1 |F_N(x)|^2 dx = \int_0^1 F_N(x) F_N(-x) dx \dots$$

Gives  $\sum_{p \leq N} \log^2 p = N \log N + o(N \log N)$

↳ could see by  $p \leq \frac{N}{\log^2 N}$  and  $\frac{N}{\log^2 N} \leq p \leq N$

ASIDE:

Philosophy  
of Square  
root  
cancellation

Special Values:  $F_N(0) = F_N(1) = N + o(N)$

$$F_N(1/2) = -N + o(N)$$

↳ much larger than ave value

Idea:  $F_N$  "large" at  $x$  near  $1/2$  with  $\epsilon$  "small"  
and  $F_N$  "small" otherwise

- Split  $\int_0^1 F_N(x) e(-nx) dx$  into two integrals,  
one where  $F_N$  "large", other "small"

↳ technicalities galore!

# THE CIRCLE METHOD: SIZE OF $F_N(x)$ (CONT)

THM: Fix  $B$ , set  $Q = \log^B N$ , take  $a, q$  rel prime with  $q \leq Q$

$$\text{then } \forall c \text{ have } F_N\left(\frac{a}{q}\right) = \frac{N}{\phi(q)} \sum_{\substack{r=1 \\ (r,q)=1}}^q \mathcal{E}\left(\frac{ar}{q}\right) + O\left(\frac{N}{\log^{c-B} N}\right)$$

$$\text{Proof: } F_N\left(\frac{a}{q}\right) = \sum_{p \leq N} \log p \cdot \mathcal{E}\left(p \frac{a}{q}\right)$$

only depends on  $p \bmod q$

$$= \sum_{r=1}^q \log \mathcal{E}\left(\frac{ar}{q}\right) \sum_{\substack{p \equiv r(q) \\ p \leq N}} \log p$$

trivial estimation + Siegel-Walfisz

$$\text{MAJOR ARCS: } \mathcal{M} = \bigcup_{q=1}^Q \bigcup_{\substack{a=1 \\ (a,q)=1}}^q \mathcal{M}_{a,q}$$

$$\text{where } \mathcal{M}_{a,q} = \left\{ x \in [0,1] : \left| x - \frac{a}{q} \right| < \frac{Q}{N} \right\}$$

$$\text{and } \mathcal{M}_{1,1} = \left[ 0, \frac{Q}{N} \right) \cup \left( 1 - \frac{Q}{N}, 1 \right]$$

$$\text{MINOR ARCS: } \mathcal{m} = [0,1] - \mathcal{M}$$

$$\text{Trivially have } |\mathcal{M}| \leq \sum_{q=1}^Q q \cdot \frac{2Q}{N} \leq \frac{2Q^3}{N}$$

$$\text{as } Q = \log^B N, \text{ see } \lim_{N \rightarrow \infty} |\mathcal{M}| = 0, \lim_{N \rightarrow \infty} |\mathcal{m}| = 1$$

↳ Discuss terminology for major/minor arcs



# 13.3.5. THE MAJOR ARCS AND THE SINGULAR SERIES

## MAJOR ARC CONTRIBUTION (HEURISTIC)

lets do  $s=3$  and  $m=N$  (ie,  $P_1+P_2+P_3=N$ )

$$\int_0^1 F_N(x)^3 dx \approx \sum_{\substack{P_1, P_2, P_3 \leq N \\ P_1+P_2+P_3=N}} \log P_1 \log P_2 \log P_3$$

Major arcs

$$\int_M F_N(x)^3 e(-Nx) dx = \sum_{q=1}^Q \sum_{\substack{a=1 \\ (a,q)=1}}^q \int_{M_{a,q}} F_N(x)^3 e(-Nx) dx$$

↳ estimate  $F_N(x)$  by  $F_N(\frac{a}{q})$

ignore all lower order terms, note  $|M_{a,q}| = 2Q/N$

$$\text{yields } \frac{2Q}{N} \sum_{q=1}^Q \sum_{\substack{r=1 \\ (r,q)=1}}^q \left( \sum_{\substack{r=1 \\ (r,q)=1}}^q e\left(\frac{ar}{q}\right) \right)^3 e\left(-\frac{Na}{q}\right)$$

$$= \mathcal{G}(N) \cdot (N^2 \cdot 2Q) \text{ with } Q = \log^B N$$

$$\text{where } \mathcal{G}(N) = \sum_{q=1}^{\infty} \frac{1}{\phi(q)^3} \sum_{\substack{a=1 \\ (a,q)=1}}^q \left( \sum_{\substack{r=1 \\ (r,q)=1}}^q e\left(\frac{ar}{q}\right) \right)^3 e\left(-N\frac{a}{q}\right)$$

↳ More careful analysis (see Chap 14) gives  $\mathcal{G}(N) \cdot \frac{N^2}{2}$

Comments: ① Close: correct power of  $N$ , missed by a few logarithms, which could matter

② not surprising, as doing a zeroth order Taylor Expansion.

↳ surprisingly, do much better (ie, easier) NOT to Taylor Expand  $F_N(x)$  but to find a  $f_n(u_N(x))$  st  $u_N(a/q) = F_N(a/q)$  and  $u_N$  easily integrated

③ Main term is from this (we hope!), but tough to look at  $\mathcal{G}(N)$  and "see" its behavior: ie:  $\mathcal{G}(1776) = 0!$

# THE SINGULAR SERIES

$$G(N) = \sum_{e=1}^{\infty} \frac{1}{\phi(e)^3} \sum_{\substack{a=1 \\ (a,e)=1}}^e \left( \sum_{\substack{r=1 \\ (r,e)=1}}^e \phi\left(\frac{ar}{e}\right) \right)^3 \phi\left(-N \frac{a}{e}\right)$$

Using methods of Chapter 14:

$$G(N) = \prod_p \left( 1 - \frac{c_p(N)}{\phi(p)^3} \right) \text{ with } c_p(N) = \begin{cases} p-1 & \text{if } p|N \\ 0 & \text{otherwise} \end{cases}$$

↳ very important

says  $G(N)$  is a multiplicative function (see Defn 2.1.2)

view as  $G(N) = \prod_p \delta_p(N)$ , each  $\delta_p(N)$  a

"local density", measuring obstructions

↳ if  $2|N$  then  $\delta_2(N) = 0$  and main term  $G(N) \frac{N^2}{2} = 0!$

↳ What does this mean?

Circle Method knows Goldbach ( $P_1 + P_2 = 2m$ ) hard

Say  $P_1 + P_2 + P_3 = 2N = 2m + 2$

Then wlog  $P_3 = 2$  and writing  $2m$  as sum of two primes!

ASIDE ON OBSTRUCTIONS: Hasse Principle, Section 4.4

Exe 13.3.16: Important

N odd  $\rightarrow \exists c_1, c_2 > 0$  (and indep of  $N$ ) s.t.  $c_1 < G(N) < c_2$



## MINOR ARCS CONTRIBUTION

Know  $|F_N(x)|$ , on average, is about  $\sqrt{N \log N}$  and at most  $N$   
can be as large as order  $N$  near  $0/e$  with  $q \ll N$ .

Vinogradov:  $\max_{x \in m} |F_N(x)| \ll \frac{N}{\log^D N}$   $D = D(B, C) > 1$

$$\begin{aligned} \left| \int_m F_N^3(x) e(-Nx) dx \right| &\leq \int_m |F_N(x)|^2 \cdot |F_N(x)| dx \\ &\leq \max_{x \in m} |F_N(x)| \cdot \int_m |F_N(x)|^2 dx \\ &\ll \frac{N}{\log^D N} \cdot \int_0^1 |F_N(x)|^2 dx \\ &= \frac{N}{\log^D N} \cdot N \log N = \frac{N^2}{\log^{D-1} N} \end{aligned}$$

Thus a "small" cancellation is enough to win when  $s=3$

Aside: Very common/important technique  
easy to evaluate  $\int_0^1 |F_N(x)|^{2k} dx$ , and  
need slight savings in an  $|F_N(x)|$  for  $x \in m$

# WHY IS GOLDBACH HARD?

$$\text{Cauchy-Schwarz (verb)}: \left| \int_0^1 fg \right| \leq \left( \int_0^1 f^2 \right)^{\frac{1}{2}} \left( \int_0^1 g^2 \right)^{\frac{1}{2}}$$

↳ Proof: Lemma A.6.9

When  $s=2$  ( $p_1+p_2=N$ ) Major arcs give  $\sim \mathcal{O}_2(N) \cdot N$

Estimate minor arcs: insert abs values (lose all oscillation)

$$\text{↳ } \int_m |F_N(x)|^2 dx \sim N \log N \quad \text{too big}$$

$$\text{↳ Pull out } |F_N(x)|: \max_{x \in m} |F_N(x)| \cdot \int_m |F_N(x)| dx$$

↳ real bad: even if  $|F_N(x)|$  is always its average value,  $\sqrt{N \log N}$   
we get at least  $N \log N$  again

↳ What if "most" of the time  $|F_N(x)|$  small?

Using Cauchy-Schwarz:

$$\begin{aligned} \int_m |F_N(x)| dx &\leq \left( \int_m |F_N(x)|^2 dx \right)^{1/2} \left( \int_m 1^2 dx \right)^{1/2} \\ &\leq (N \log N)^{1/2} \cdot 1 \end{aligned}$$

(note: applying Cauchy-Schwarz good on finite intervals)

↳ still need more than average value (cancellation!)

ASIDE: Littlewood Problem (includes results of Paul Cohen)

Remark: Is Goldbach hopeless?

↳ This is too crude of a method: moving absolute value inside the integration loses all oscillation.

ASIDE: Research projects / detailed calculations in Chapter 14 for German Primes