## Introduction to Linear Programming

### Steven Miller -- Fall 2014

#### ? LinearProgramming

LinearProgramming[c, m, b] finds a vector x that minimizes the quantity c.x subject to the constraints  $m.x \ge b$  and  $x \ge 0$ . LinearProgramming[ $c, m, \{\{b_1, s_1\}, \{b_2, s_2\}, ...\}$ ] finds a vector x that minimizes c.x subject

to  $x \ge 0$  and linear constraints specified by the matrix *m* and the pairs  $\{b_i, s_i\}$ . For each row  $m_i$  of  $m_i$ 

the corresponding constraint is  $m_i \cdot x \ge b_i$  if  $s_i == 1$ , or  $m_i \cdot x == b_i$  if  $s_i == 0$ , or  $m_i \cdot x \le b_i$  if  $s_i == -1$ .

LinearProgramming[c, m, b, l] minimizes c.x subject to the constraints specified by m and b and  $x \ge l$ .

LinearProgramming[ $c, m, b, \{l_1, l_2, ...\}$ ] minimizes c.x subject to the constraints specified by m and b and  $x_i \ge l_i$ .

LinearProgramming[ $c, m, b, \{\{l_1, u_1\}, \{l_2, u_2\}, ...\}$ ] minimizes c.x subject to the constraints specified by m and b and  $l_i \le x_i \le u_i$ .

LinearProgramming[c, m, b, lu, dom] takes the elements of x to be in the domain dom, either Reals or Integers.

LinearProgramming[ $c, m, b, lu, \{dom_1, dom_2, ...\}$ ] takes  $x_i$  to be in the domain  $dom_i$ .  $\gg$ 

## Diet Problem : 10 x + 4 y >= 9, 5 x + 8 y >= 11, 3 x + 2 y >= 5, minimize 2 x + 3 y

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In[26]:= (* Here is how we do a simple linear
     programming problem if real variables and sizes small *)
     (* we first enter the matrix A, where we enter row by row *)
     (* we then enter the b-vector,
    which must have the same number of entries as rows of A *)
     (* we then enter the cost vector c,
    which must have the same number of entries as the columns of A *)
     (* the way the code works is we solve Ax \ge b, x \ge 0, min(c.x) *)
    A = \{\{10, 4\}, \{5, 8\}, \{3, 2\}\};\
    b = \{9, 11, 5\};
    c = \{2, 3\};
    soln = LinearProgramming[c, A, b];
    Print["The solution is ", soln]
    Print["The cost is ", c.soln]
    Print["Now let's do it where we require integer values."];
    intsoln = LinearProgramming[c, A, b, Automatic, Integers];
    Print["The solution in integers is ", intsoln]
    Print["The integer solution cost is ", c.intsoln]
```

The solution is  $\left\{\frac{9}{7}, \frac{4}{7}\right\}$ The cost is  $\frac{30}{7}$ Now let's do it where we require integer values. LinearProgramming::lpip: Warning: integer linear programming will use a machine-precision approximation of the inputs. >> The solution in integers is {1, 1}

The integer solution cost is 5

# Chess Problem: n queens on n x n board ....

```
In[3]:= (* making this a program so can call it with different arguments easily *)
    (* have a board that is boardsize by boardsize and placing numberofqueens *)
    (* use underscore to indicate a variable *)
    (* Module indicates module / procedure / function,
    the {} means no local variables *)
    (* Thus all output can be felt / seen later *)
    (* if printlinearprogrammingconditions is 1 print the matrix A and rest,
    else don't *)
   pawnqueenproblem[boardsize_, numberofqueens_,
      printlinearprogrammingconditions_] := Module[{},
      n = boardsize; (*just change for different board *)
      numqueens = numberofqueens;
      (* allows the number of queens to differ from board size *)
      (* when programming you want freedom -- see
       what matters and what depends on what *)
      (* in some of the constraints what matters is the board size,
      and in others its the number of queens *)
      (* variables x1, ..., xn, x_{n+1} = x_{2,1}, ... *)
      (* followed by p1, ..., pn, p_{n+1}, ... *)
      (* the above requires some explanation. We
       can't have our variables having indices unfortunately *)
      (* we need a linear list of variables, not an array *)
      (* fortunately it's very simple to pass from indices to a linear list *)
      (* if we have x_{ij} with 1 \le i, j \le n then we let num(i,j) = (i-1)*n + j *)
      (* note that counts from 1 to n^2 *)
      (* we have two variables, the x_ij,
      which say whether or not a queen is at (i,j), and *)
```

```
(* p ij, which say whether or not we can place a pawn at (i,j). *)
(* so x_{ij} \rightarrow x_{num}(i,j) is 1 if a queen is at (i,j) and 0 otherwise *)
(* and p_ij -->
 x_{n^2 + num(i,j)} is 1 if a pawn is at (i,j) and 0 otherwise *)
(* notice that we are always using x subscript
 an integer for our variables. *)
(* we write the queen variables first and then the pawn variables,
and hence adding n^2 in the subscript *)
A = {}; (* initializes our constraint matrix A to be empty,
we'll add constraints *)
(* notice we don't want to type all the constraints by hand,
but write code to add *)
bvec = {}; (* initializes bvector to empty *)
(* the next lines take into account
 which squares on the chessboard can attack (i,j) *)
(* we first record which squares can have a queen attacking (i,j),
remembering we count linearly and must convert *)
(* numbers to pairs; thus if we want to investigate what happens
  for a queen at (a,b) that corresponds to index (a-1)*n + b *)
(* while similarly we could go from inex (a-1)*n + b to (a,b);
there is some small issue with how Mathematica looks at remainders *)
(* and so we change b if it is 0 mod n to n. *)
(* the constraint is the following:
  -Sum_{(a,b)} \text{ attacks } (i,j) \times_{ab} - numqueens p_{ij} \ge - numsqueens *)
(* this is NOT my first choice for how to write the constraint,
but remember Mathematica does Ax \ge b *)
(* eventually we will maximize the sum of p_{ij},
so if it is available for a pawn the program will place one there *)
(* (of course we maximize the sum of p_{ij} by minimizing the sum of -p_{ij},
 as we work with minimums!) *)
(* returning to the constraint: if all the x_{ab} that can
   attack (i,j) are zero then we may take p_{ij} to be 1 *)
(* if even one x_{ab} is 1 then the constraint is too negative on
 the LHS if p_{ij} is 1, and thus we have p_{ij} = 0 as desired *)
For[i = 1, i \le n, i++,
 For[j = 1, j \le n, j++,
  (* the i and j for statements go over all board locations *)
  {
   temp = {}; (* initialize temp to be empty,
   we'll start putting the constraint info here and then append to A *)
   For [num = 1, num \leq n^2, num++, (* remember we index variables linearly,
    this goes over the n<sup>2</sup> squares that can attack (i,j) *)
     b = Mod[num, n];
     If[b = 0, b = n];
     a = ((num - b) / n) + 1; (* this converts from the linear index
```

```
to a pair (a,b) for the board space under consideration *)
      (*numat = (a-1)*n + b;*) (* moves from (a,b) to num from 1 to n^2 *)
      (* we now see if a queen at (a,b) could attack
      square (i,j). there are four ways this could happen. *)
      (* they could have the same x-coord, so i = a;
      they could have the same y-coord, so j = b. *)
      (* they could also be on the same upward sloping diagonal,
      so i-j = a-b, or downward, so i-j = a-b *)
      (* the || is how we do an or condition; if either of the four
        conditions hold we append a -1 to our list, else a 0 *)
      (* this will give us n^2 elements, the first half of a row for A;
     we then do the p_ij entries *)
     temp = AppendTo[temp, If[a == i || b == j ||
          i-j = a-b || i+j = a+b, -1, 0]];
    }]; (* end of nuum loop *)
   numij = (i - 1) * n + j;
   (* this gives the linear index corresponding to (i,j) *)
   (* we now finish the row constraint; we put a -
    numqueens at the place corresponding to (i,j) and a 0 elsewhere *)
   For [num = 1, num \leq n^2, num++, temp =
     AppendTo[temp, If[num == numij, -numqueens, 0]]];
   A = AppendTo[A, temp]; (* add constraint to A *)
   bvec = AppendTo[bvec, -numqueens]; (* add the entry to bvector *)
  }]; (* end of j loop *)
]; (* end of i loop *)
(* now we want to add constraints saying there are EXACTLY numqueen queens,
ie, numqueen of the xij are 1 and the rest are 0 *)
(* as Mathematica does Ax \ge b we need two constraints to get an
  equality: one blah \geq numqueens and one -blah \geq -numqueens. *)
(* we initialize temp to be zero, and put a 1 in the first n<sup>2</sup>
 elements (those corresponding to xij) and a 0 elsewhere *)
(* we then put a numqueens for the entry of b-vec;
this gives the constraint sum xij ≥ numqueens *)
(* the next lines are similar and give -sum x ij \geq -numqueens *)
temp = {}; (* always remember to reinitialize
 the temp list to empty before adding things to it! *)
For [num = 1, num \leq 2n^2, num++, temp = AppendTo[temp, If[num \leq n^2, 1, 0]];
A = AppendTo[A, temp]; (* now we add our constraint to the A matrix *)
bvec = AppendTo[bvec, numqueens];
temp = \{\};
For [num = 1, num \leq 2n^2, num++, temp = AppendTo[temp, If[num \leq n^2, -1, 0]];
A = AppendTo[A, temp];
bvec = AppendTo[bvec, -numqueens];
(* this finishes making sure we have numqueen queens *)
(* annoyingly I could only find commands for
 Mathematica to do integer programming, NOT binary programming *)
```

```
(* we can declare the variables to be integers but not 0,
1 integers. fortunately this is easily fixed *)
(* we just need to add a constraint that each variable is at most 1,
as they are assumed to be at least 0 (advantages canonical form!) *)
(* As the constraints are always Ax \ge b,
if we want x \leq 1 we have to program that as -x \geq -1 *)
(* for each of the 2n^2 variables, the n^2 choices of xij and the n^2 of pij,
we make sure it is at most 1 *)
temp = {}; (* always initialize to empty *)
For [num = 1, num \leq 2n<sup>2</sup>, num++, (* go through the 2n<sup>2</sup> variables *)
 {
  temp = {}; (* for each variable choice make our list empty,
  put a -1 in the right spot and 0's elsewhere *)
  For[counter = 1, counter ≤ 2 n^2, counter++, AppendTo[temp,
    If [counter == num, -1, 0]]; (* if in right spot 1, else 0 *)
  A = AppendTo[A, temp]; (* appends new constraint to A,
  and then next line appends -1 as needed to b-vector *)
  bvec = AppendTo[bvec, -1];
 }];
(* now we deine the vector needed for the optimization. we
 initiallize it to empty, and then make it the right size *)
(* since we can only do minima we use -1 as the entries for c
  corresponding to the pawn variable locations, and 0 for the queen *)
c = {};
For [num = 1, num \le 2n^2, num + +, c = AppendTo[c, If[num > n^2, -1, 0]]];
(* appending information to c *)
(* below is the key line -- it calls the linear program solver,
and the last bits inform it that the variables are integers *)
(* we save the output to a quantity we named soln,
for solution; we will then format the answer nicely *)
soln = LinearProgramming[c, A, bvec, Automatic, Integers];
Print["Solution is ", soln]; (* prints the soln vector,
but hard to parse so we work on it a bit *)
queenlist = {}; (* initializes the list of queens to empty,
this is the first n<sup>2</sup> variables *)
(* we then go through the soln list and save that info to the queen list *)
(* the next lines after this redo
 this and make a list of pawn solution information *)
(* we really don't need to do this -- we can work directly with
 the solution list, but thought this might be easier to parse *)
For[num = 1, num ≤ Length[soln], num++,
 If[num ≤ n^2, queenlist = AppendTo[queenlist, soln[[num]]]];
pawnlist = {};
For[num = 1, num ≤ Length[soln], num++,
 If[num > n^2, pawnlist = AppendTo[pawnlist, soln[[num]]]]];
```

```
(* now we will draw a board and place
  Q for queen at P for pawn at the correct locations *)
board = {}; (* as always initialize to empty, this will be a matrix *)
For [i = 1, i \le n, i++, (* will construct the board row by row *)
  {
   temp = {}; (* initialize new row of matrix to empty and will add *)
   For [j = 1, j \le n, j++, (* go through the n elements of the row *)
    {
     numat = (i-1) * n + j;
     (* convert from (i,j) board location to linear index number *)
     If[queenlist[[numat]] == 1, temp = AppendTo[temp, "Q"],
      (* if queen there write Q *)
      If[pawnlist[[numat]] == 1, temp = AppendTo[temp, "P"],
       (* if pawn there write P *)
       temp = AppendTo[temp, "-"]]]; (* else write -
      to show space empty but show the space is there *)
    }];
   board = AppendTo[board, temp]; (* save the new row to the board matrix *)
  }1;
 Print[MatrixForm[board]]; (* print the board matrix NICELY *)
 Print["Number of pawns = ", Sum[pawnlist[[i]], {i, 1, Length[pawnlist]}];
 (* prints number pawns *)
 Print["Number of queens = ", Sum[queenlist[[i]], {i, 1, Length[queenlist]}];
 (* prints number queens *)
 If[printlinearprogrammingconditions == 1,
  {
   Print["Constraint Matrix is"];
   Print[MatrixForm[A]]; (* prints the constraint matrix *)
   Print["b-vector is"];
   Print[bvec]; (* prints the b-vector *)
   Print["c vector for optimization is"];
   Print[c] ; (* prints the c vector *)
  }]; (* end of print condition --
  only print if printlinearprogrammingconditions is 1 *)
]; (* end of module *)
```

#### In[4]:= pawnqueenproblem[1, 1, 0]

LinearProgramming::lpip : Warning: integer linear programming will use a machine-precision approximation of the inputs. >>
Solution is {1, 0}
(Q)
Number of pawns = 0
Number of queens = 1

#### In[5]:= pawnqueenproblem[2, 2, 0]

LinearProgramming:: lpip : Warning: integer linear programming will use a machine-precision approximation of the inputs. >>

```
Solution is {1, 1, 0, 0, 0, 0, 0, 0}

\begin{pmatrix} Q & Q \\ - & - \end{pmatrix}

Number of pawns = 0

Number of queens = 2
```

#### In[6]:= pawnqueenproblem[3, 3, 0]

LinearProgramming:: lpip: Warning: integer linear programming will use a machine-precision approximation of the inputs.  $\gg$  Solution is {0, 0, 1, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0}

```
\begin{pmatrix} - & - & Q \\ - & - & - \\ Q & Q & - \end{pmatrix}
Number of pawns = 0
Number of queens = 3
```

#### In[18]:= Timing[pawnqueenproblem[4, 4, 0]]

LinearProgramming:: lpip : Warning: integer linear programming will use a machine-precision approximation of the inputs. >>

#### In[20]:= Timing[pawnqueenproblem[5, 5, 1]]

- P - - P

LinearProgramming:: lpip: Warning: integer linear programming will use a machine-precision approximation of the inputs. >>

- - - P Q - - - -Q – – Q – - - 0 0 -Number of pawns = 3Number of queens = 5Constraint Matrix is 0 0 0 -1 0 0 0 0 -1 0 0 0 -1 -1 -1 -1 -1 0 0 -1 -1 0 -1 0 -1 0 -1 00 -1 0 0 0 -1 0 0 0 -1 -1 -1 -1 -1 0 0 0 -1 -1 0 0 -1 0 -1 0-1 0 0 -1 -1 0 0 0 -1 0 -1 -1 0 0 0 -1 -1 -1 -1 -1 -1 -1 0 0 0 - 1 0 -1 0 0 - 1 0 0 - 1 0 0 -1 -1 -1 0 0 -1 -1 -1 -1 -1 -1 -1 -1 0 0 0 -1 0 -1 0 0 -1 0 0 -1 0 0  $\sim$  $\cap$  $\circ$ 1 1 1  $\cap$  $\cap$ 

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0 b-vector is 

c vector for optimization is

Out[20]= {0.436803, Null}

#### In[21]:= Timing[pawnqueenproblem[6, 6, 0]]

LinearProgramming:: lpip : Warning: integer linear programming will use a machine-precision approximation of the inputs. >>

#### In[22]:= Timing[pawnqueenproblem[7, 7, 0]]

LinearProgramming:: lpip : Warning: integer linear programming will use a machine-precision approximation of the inputs. >>

#### In[23]:= Timing[pawnqueenproblem[8, 8, 0]]

LinearProgramming:: lpip : Warning: integer linear programming will use a machine-precision approximation of the inputs. >>

In[24]:= 2759 / 60.

Out[24]= 45.9833