

# Benford's Law, Values of $L$ -Functions and the $3x + 1$ Problem, or: Why the IRS should care about Number Theory

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[http://web.williams.edu/Mathematics/sjmilller/public\\_html/index.htm](http://web.williams.edu/Mathematics/sjmilller/public_html/index.htm)

Five Colleges Number Theory Seminar  
January 31, 2012

# Introduction

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**Plausible answers:** 10%, 11%, about 30%.

## Summary

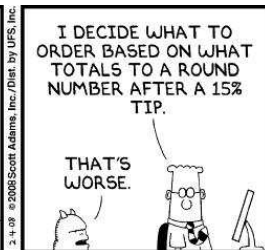
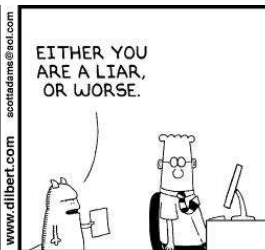
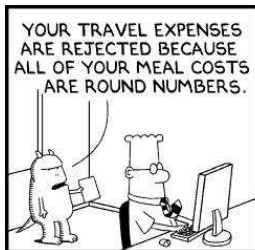
- State Benford's Law.
- Discuss examples and applications.
- Sketch proofs.
- Describe open problems.

## Caveats!

- A math test indicating fraud is *not* proof of fraud:  
unlikely events, alternate reasons.

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  - ◇ **Many streets of different sizes: close to Benford.**

## Examples

- recurrence relations
- special functions (such as  $n!$ )
- iterates of power, exponential, rational maps
- products of random variables
- $L$ -functions, characteristic polynomials
- iterates of the  $3x + 1$  map
- differences of order statistics
- hydrology and financial data
- many hierarchical Bayesian models

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## Applications

- analyzing round-off errors
- determining the optimal way to store numbers
- detecting tax and image fraud, and data integrity

## General Theory

## Mantissas (or Significands)

Mantissa:  $x = M_{10}(x) \cdot 10^k$ ,  $k$  integer.

$M_{10}(x) = M_{10}(\tilde{x})$  if and only if  $x$  and  $\tilde{x}$  have the same leading digits.

**Key observation:**  $\log_{10}(x) = \log_{10}(\tilde{x}) \bmod 1$  if and only if  $x$  and  $\tilde{x}$  have the same leading digits. Thus often study  $y = \log_{10} x$ .

## Equidistribution and Benford's Law

### Equidistribution

$\{y_n\}_{n=1}^{\infty}$  is equidistributed modulo 1 if probability  $y_n \bmod 1 \in [a, b]$  tends to  $b - a$ :

$$\frac{\#\{n \leq N : y_n \bmod 1 \in [a, b]\}}{N} \rightarrow b - a.$$

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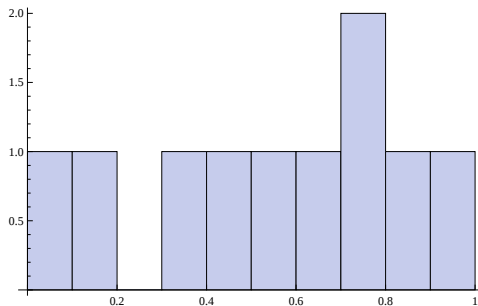
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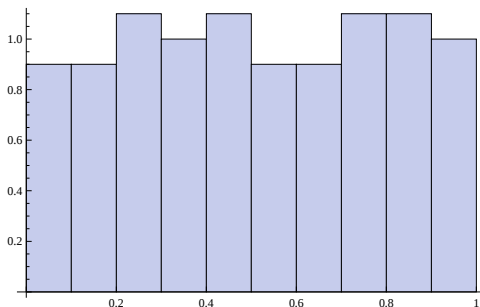
- Thm:  $\beta \notin \mathbb{Q}$ ,  $n\beta$  is equidistributed mod 1.
- Examples:  $\log_{10} 2, \log_{10} \left( \frac{1+\sqrt{5}}{2} \right) \notin \mathbb{Q}$ .

## Example of Equidistribution: $n\sqrt{\pi} \bmod 1$



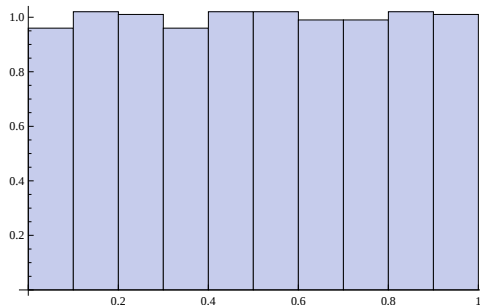
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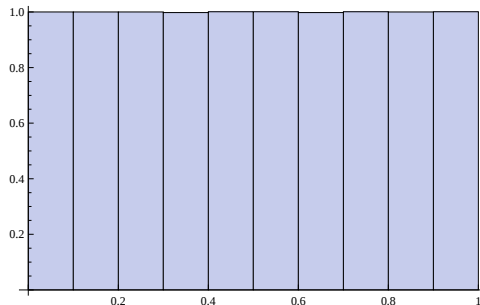
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$n\sqrt{\pi} \bmod 1$  for  $n \leq 1000$

## Example of Equidistribution: $n\sqrt{\pi} \bmod 1$



$n\sqrt{\pi} \bmod 1$  for  $n \leq 10,000$

## Logarithms and Benford's Law

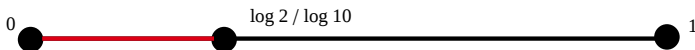
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Data set  $\{x_i\}$  is Benford base  $B$  if  $\{y_i\}$  is equidistributed mod 1, where  $y_i = \log_B x_i$ .

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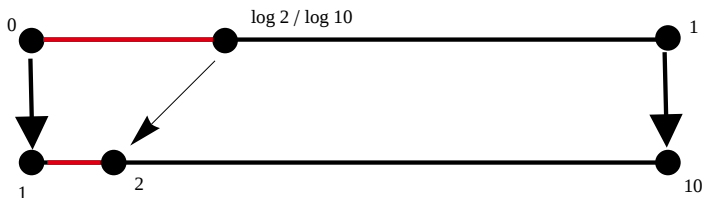
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$$a_{n+1} = a_n + a_{n-1}.$$

$$\text{Binet: } a_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^n.$$

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- Most linear recurrence relations Benford.

$$a_{n+1} = 2a_n$$

$$a_{n+1} = 2a_n - a_{n-1}$$

Take  $a_0 = a_1 = 1$  or  $a_0 = 0, a_1 = 1$ .

# Digits of $2^n$

First 60 values of  $2^n$  (only displaying 30)

|     |        |           | digit | #  | Obs Prob | Benf Prob |
|-----|--------|-----------|-------|----|----------|-----------|
| 1   | 1024   | 1048576   |       |    |          |           |
| 2   | 2048   | 2097152   | 1     | 18 | .300     | .301      |
| 4   | 4096   | 4194304   | 2     | 12 | .200     | .176      |
| 8   | 8192   | 8388608   | 3     | 6  | .100     | .125      |
| 16  | 16384  | 16777216  | 4     | 6  | .100     | .097      |
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| 128 | 131072 | 134217728 | 7     | 2  | .033     | .058      |
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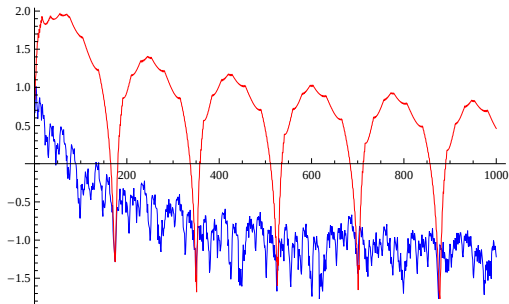
## Logarithms and Benford's Law

$\chi^2$  values for  $\alpha^n$ ,  $1 \leq n \leq N$  (5% 15.5).

| $N$  | $\chi^2(\gamma)$ | $\chi^2(e)$ | $\chi^2(\pi)$ |
|------|------------------|-------------|---------------|
| 100  | 0.72             | 0.30        | 46.65         |
| 200  | 0.24             | 0.30        | 8.58          |
| 400  | 0.14             | 0.10        | 10.55         |
| 500  | 0.08             | 0.07        | 2.69          |
| 700  | 0.19             | 0.04        | 0.05          |
| 800  | 0.04             | 0.03        | 6.19          |
| 900  | 0.09             | 0.09        | 1.71          |
| 1000 | 0.02             | 0.06        | 2.90          |

## Logarithms and Benford's Law: Base 10

$\log_{10}(\chi^2)$  vs  $N$  for  $\pi^n$  (red) and  $e^n$  (blue),  
 $n \in \{1, \dots, N\}$ . Note  $\pi^{175} \approx 1.0028 \cdot 10^{87}$ , (5%  
and 8 d.f.,  $\log_{10}(\chi^2) \approx .44$ ).



## Applications

## Applications for the IRS: Detecting Fraud

[illegible]

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1040 U.S. Individual Income Tax Return 1992

93-4670

Label: WILLIAM J. CLINTON  
HILLARY RODHAM CLINTON  
THE WHITE HOUSE  
1600 PENNSYLVANIA AVENUE N.W.  
WASHINGTON, DC 20500

Use the IRS label. Otherwise, place name in type.

President: Do you want \$1 to go to this fund? ☒ Yes ☐ No  
Election Campaign: If you return, does your spouse want \$1 to go to the fund? ☒ Yes ☐ No

Filing Status: ☒ Married filing joint return (even if only one had income)  
☐ Married filing separate return. Enter spouse's SSN above and full name here. ☐  
☐ Single  
☐ Head of household  
☐ Qualifying widow(er) with dependent child

Check only one box.

Exemptions: ☒ Yourself. If your spouse or someone else can claim you as a dependent on his or her return, do not check this box. ☐ Other. If you are claiming an exemption for someone else, check this box on line 2b or page 2.

Dependent: ☐ If you are claiming an exemption for someone else, check this box on line 2b or page 2.

Child: ☒ DAUGHTER

Other: ☐ If you are claiming an exemption for someone else, check this box on line 2b or page 2.

Income:

7 Wages, salaries, tips, etc. (Attach Form W-2) **237,659.7**

8 Taxable interest income. (Do not include on line 8a) **6,624.00**

9 Tax-exempt interest income. (Do not include on line 8a) **743.00**

10 Dividend income. Attach Schedule D if over \$400 **1,401.00**

11 Taxable refunds, credits, or offsets of state and local income taxes **16,336.00**

12 Alimony received **16,336.00**

13 Business income or (loss). Attach Schedule C or C-EZ **16,336.00**

14 Capital gain or (loss). Attach Schedule D **16,336.00**

15 Other gains or (losses). Attach Form 4797 **16,336.00**

16 Total IRA distributions **16,336.00**

17 Total pensions and annuities **16,336.00**

18 Rents, royalties, partnerships, estates, trusts, etc. Attach Schedule E **16,336.00**

19 Farm income or (loss). Attach Schedule F **16,336.00**

20 Unemployment compensation **16,336.00**

21 Social Security benefits **16,336.00**

22 Other income. **16,336.00**

23 Add the amounts in the far right column for lines 7 through 22. This is your total income **297,172.00**

Adjustments to income:

24 Your IRA deduction **250.00**

25 Spouse's IRA deduction **250.00**

26 Overhead of self-employment tax **250.00**

27 Self-employed health insurance deduction **250.00**

28 Keogh retirement plan and self-employed SEP deduction **250.00**

29 Penalty on early withdrawal of savings **250.00**

30 Alimony paid. Recipient's SSN: **250.00**

31 Add lines 24 through 30. This is your total adjustments **250.00**

32 Subtract line 31 from line 23. This is your adjusted gross income **296,922.00**

AGI **296,922.00**

Form 1040 (1992)

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- Write-off limit of \$5,000. Officer had friends applying for credit cards, ran up balances just under \$5,000 then he would write the debts off.

## Benford Good Processes

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- Poisson Summation Formula:  $f$  nice:

$$\sum_{\ell=-\infty}^{\infty} f(\ell) = \sum_{\ell=-\infty}^{\infty} \hat{f}(\ell),$$

$$\text{Fourier transform } \hat{f}(\xi) = \int_{-\infty}^{\infty} f(x) e^{-2\pi i x \xi} dx.$$

## Benford Good Process

$X_T$  is **Benford Good** if there is a nice  $f$  st

$$\text{CDF}_{\vec{Y}_{T,B}}(y) = \int_{-\infty}^y \frac{1}{T} f\left(\frac{t}{T}\right) dt + E_T(y) := G_T(y)$$

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- **Small translated error:**  $\mathcal{E}(a, b, T) =$   
 $\sum_{|\ell| \leq Th(T)} [E_T(b + \ell) - E_T(a + \ell)] = o(1).$

## Main Theorem

### Theorem (Kontorovich and M–, 2005)

*$X_T$  converging to  $X$  as  $T \rightarrow \infty$  (think spreading Gaussian). If  $X_T$  is Benford good, then  $X$  is Benford.*

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- **Examples**

- ◇  $L$ -functions
- ◇ characteristic polynomials (RMT)
- ◇  $3x + 1$  problem
- ◇ geometric Brownian motion.

## Sketch of the proof

- **Structure Theorem:**
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  - ◇ main term is something nice spreading out
  - ◇ apply Poisson summation
- **Control translated errors:**
  - ◇ hardest step
  - ◇ techniques problem specific

## Sketch of the proof (continued)

$$\sum_{\ell=-\infty}^{\infty} \mathbb{P} \left( a + \ell \leq \overrightarrow{Y}_{T,B} \leq b + \ell \right)$$

## Sketch of the proof (continued)

$$\begin{aligned}
 & \sum_{\ell=-\infty}^{\infty} \mathbb{P} \left( a + \ell \leq \overrightarrow{Y}_{T,B} \leq b + \ell \right) \\
 &= \sum_{|\ell| \leq Th(T)} [G_T(b + \ell) - G_T(a + \ell)] + o(1)
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 &= \int_a^b \sum_{|\ell| \leq Th(T)} \frac{1}{T} f \left( \frac{t}{T} \right) dt + \mathcal{E}(a, b, T) + o(1)
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 &= \int_a^b \sum_{|\ell| \leq Th(T)} \frac{1}{T} f \left( \frac{t}{T} \right) dt + \mathcal{E}(a, b, T) + o(1) \\
 &= \hat{f}(0) \cdot (b - a) + \sum_{\ell \neq 0} \hat{f}(T\ell) \frac{e^{2\pi i b \ell} - e^{2\pi i a \ell}}{2\pi i \ell} + o(1).
 \end{aligned}$$

## *L*-functions and Random Matrix Theory

## Good $L$ -functions ( $\zeta(s)$ , full level cusp form)

$L$ -function is **good** if:

- Euler product:

$$L(s, f) = \sum_{n=1}^{\infty} \frac{a_f(n)}{n^s} = \prod_{p \text{ prime}} \prod_{j=1}^d (1 - \alpha_{f,j}(p)p^{-s})^{-1}.$$

- meromorphic continuation to  $\mathbb{C}$ , of finite order, at most finitely many poles (all on the line  $\operatorname{Re}(s) = 1$ ).
- Functional equation:  $\omega \in \mathbb{R}$ ,  $G(s)$  prod  $\Gamma$ -fns:

$$e^{i\omega} G(s) L(s, f) = e^{-i\omega} \overline{G(1 - \bar{s}) L(1 - \bar{s})}.$$

## Good L-functions

- For some  $\aleph > 0$ ,  $c \in \mathbb{C}$ ,  $x \geq 2$ :

$$\sum_{p \leq x} \frac{|a_f(p)|^2}{p} = \aleph \log \log x + c + O\left(\frac{1}{\log x}\right).$$

- The  $\alpha_{f,j}(p)$  are (Ramanujan-Petersson) tempered:  $|\alpha_{f,j}(p)| \leq 1$ .
- $N(\sigma, T) = \#\{\rho : L(\rho, f) = 0, \operatorname{Re}(\rho) \geq \sigma, \operatorname{Im}(\rho) \in [0, T]\}$ .  $\exists \beta > 0$

$$N(\sigma, T) = O\left(T^{1-\beta(\sigma-\frac{1}{2})} \log T\right).$$

# Log-Normal Law (Hejhal, Laurinćikas, Selberg)

## Log-Normal Law

$$\frac{\mu(\{t \in [T, 2T] : \log |L(\sigma + it, f)| \in [a, b]\})}{T} =$$

$$\frac{1}{\sqrt{\psi(\sigma, T)}} \int_a^b e^{-\pi u^2 / \psi(\sigma, T)} du + \text{Error}$$

$$\psi(\sigma, T) = \aleph \log \left[ \min \left( \log T, \frac{1}{\sigma - \frac{1}{2}} \right) \right] + O(1)$$

$$\frac{1}{2} \leq \sigma \leq \frac{1}{2} + \frac{1}{\log^\delta T}, \quad \delta \in (0, 1).$$

## Values of $L$ -functions and Benford's Law

### Theorem (Kontorovich and M–, 2005)

$L(s, f)$  a good  $L$ -function, as  $T \rightarrow \infty$ ,  
 $L(\sigma_T + it, f)$  is Benford.

### Ingredients

- Approximate  $\log L(\sigma_T + it, f)$  with  

$$\sum_{n \leq x} \frac{c(n)\Lambda(n)}{\log n} \frac{1}{n^{\sigma_T + it}}.$$

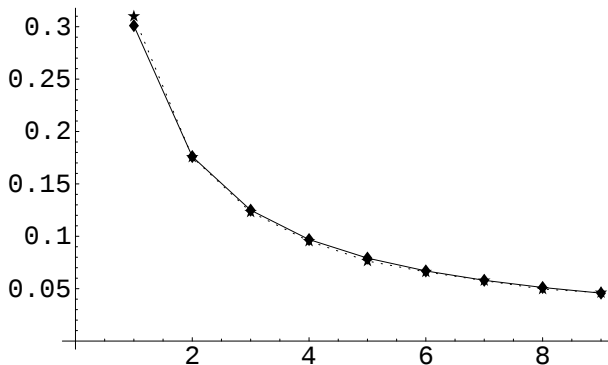
- study moments  $\int_T^{2T} |\cdot|^k, k \leq \log^{1-\delta} T.$

- Montgomery-Vaughan:

$$\int_T^{2T} \sum a_n n^{-it} \overline{\sum b_m m^{-it}} dt = H \sum a_n \bar{b}_n + O(1) \sqrt{\sum n |a_n|^2 \sum n |b_n|^2}.$$

# Riemann Zeta Function

$$\left| \zeta \left( \frac{1}{2} + i \frac{k}{4} \right) \right|, k \in \{0, 1, \dots, 65535\}.$$



## Random Matrices: Preliminaries

- $N \times N$  unitary matrices  $U$  (Haar measure):

$$p_N(U) = \frac{1}{(2\pi)^N N!} \prod_{1 \leq j < m \leq N} |e^{i\theta_j} - e^{i\theta_m}|.$$

- characteristic polynomial:

$$Z(U, \theta) = \det(I - Ue^{-i\theta}) = \prod (1 - e^{i(\theta_n - \theta)}).$$

- $\rho_N(x)$  the probability density for  $\log |Z(U, \theta)|$ :

$$\tilde{\rho}_N(x) = \sqrt{Q_2(N)} \rho_N(\sqrt{Q_2(N)} x),$$

variance  $Q_2(N) \sim (\log N)/2$ .

## Random Matrices and Benford's Law

### Theorem (Kontorovich and M–, 2005)

*As  $N \rightarrow \infty$ , the distribution of digits of the absolute values of the characteristic polynomials of  $N \times N$  unitary matrices (with respect to Haar measure) converges to the Benford probabilities.*

- Key Ingredient: Keating-Snaith:

$$\tilde{\rho}_N(x)dx = \frac{1}{\sqrt{2\pi}} e^{-x^2/2} dx + O\left((\log N)^{-3/2} dx\right).$$

# The $3x + 1$ Problem and Benford's Law

## 3x + 1 Problem

- Kakutani (conspiracy), Erdős (not ready).
- $x$  odd,  $T(x) = \frac{3x+1}{2^k}$ ,  $2^k || 3x + 1$ .

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2-path (1, 1), 5-path (1, 1, 2, 3, 4).  
 $m$ -path:  $(k_1, \dots, k_m)$ .

## Heuristic Proof of $3x + 1$ Conjecture

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Geometric Brownian Motion, drift  $\log(3/4) < 1$ .

## Structure Theorem: Sinai, Kontorovich-Sinai

$$\mathbb{P}(A) = \lim_{N \rightarrow \infty} \frac{\#\{n \leq N : n \equiv 1, 5 \pmod{6}, n \in A\}}{\#\{n \leq N : n \equiv 1, 5 \pmod{6}\}}.$$

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## 3x + 1 and Benford

### Theorem (Kontorovich and M–, 2005)

*As  $m \rightarrow \infty$ ,  $x_m / (3/4)^m x_0$  is Benford.*

### Theorem (Lagarias-Soundararajan 2006)

*$X \geq 2^N$ , for all but at most  $c(B)N^{-1/36}X$  initial seeds the distribution of the first  $N$  iterates of the  $3x + 1$  map are within  $2N^{-1/36}$  of the Benford probabilities.*

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$$C = \log_B 2 \text{ of irrationality type } \kappa < \infty:$$

$$\#\{k \in I_\ell : \overline{kC} \in [a, b]\} = M(b-a) + O(M^{1+\epsilon-1/\kappa}).$$

## Irrationality Type

### Irrationality type

$\alpha$  has irrationality type  $\kappa$  if  $\kappa$  is the supremum of all  $\gamma$  with

$$\underline{\lim}_{q \rightarrow \infty} q^{\gamma+1} \min_p \left| \alpha - \frac{p}{q} \right| = 0.$$

- Algebraic irrationals: type 1 (Roth's Thm).
- Theory of Linear Forms:  $\log_B 2$  of finite type.

## Linear Forms

### Theorem (Baker)

$\alpha_1, \dots, \alpha_n$  algebraic numbers height  $A_j \geq 4$ ,  
 $\beta_1, \dots, \beta_n \in \mathbb{Q}$  with height at most  $B \geq 4$ ,

$$\Lambda = \beta_1 \log \alpha_1 + \dots + \beta_n \log \alpha_n.$$

If  $\Lambda \neq 0$  then  $|\Lambda| > B^{-C \Omega \log \Omega'}$ , with  
 $d = [\mathbb{Q}(\alpha_i, \beta_j) : \mathbb{Q}]$ ,  $C = (16nd)^{200n}$ ,  
 $\Omega = \prod_j \log A_j$ ,  $\Omega' = \Omega / \log A_n$ .

Gives  $\log_{10} 2$  of finite type, with  $\kappa < 1.2 \cdot 10^{602}$ :

$$|\log_{10} 2 - p/q| = |q \log 2 - p \log 10| / q \log 10.$$

## Quantified Equidistribution

### Theorem (Erdős-Turan)

$$D_N = \frac{\sup_{[a,b]} |N(b-a) - \#\{n \leq N : x_n \in [a,b]\}|}{N}$$

*There is a  $C$  such that for all  $m$ :*

$$D_N \leq C \cdot \left( \frac{1}{m} + \sum_{h=1}^m \frac{1}{h} \left| \frac{1}{N} \sum_{n=1}^N e^{2\pi i h x_n} \right| \right)$$

## Proof of Erdős-Turan

Consider special case  $x_n = n\alpha$ ,  $\alpha \notin \mathbb{Q}$ .

- Exponential sum  $\leq \frac{1}{|\sin(\pi h\alpha)|} \leq \frac{1}{2||h\alpha||}$ .
- Must control  $\sum_{h=1}^m \frac{1}{h||h\alpha||}$ , see irrationality type enter.
- type  $\kappa$ ,  $\sum_{h=1}^m \frac{1}{h||h\alpha||} = O(m^{\kappa-1+\epsilon})$ , take  $m = \lfloor N^{1/\kappa} \rfloor$ .

## 3x + 1 Data: random 10,000 digit number, $2^k \parallel 3x + 1$

80,514 iterations ( $(4/3)^n = a_0$  predicts 80,319);  
 $\chi^2 = 13.5$  (5% 15.5).

| Digit | Number | Observed | Benford |
|-------|--------|----------|---------|
| 1     | 24251  | 0.301    | 0.301   |
| 2     | 14156  | 0.176    | 0.176   |
| 3     | 10227  | 0.127    | 0.125   |
| 4     | 7931   | 0.099    | 0.097   |
| 5     | 6359   | 0.079    | 0.079   |
| 6     | 5372   | 0.067    | 0.067   |
| 7     | 4476   | 0.056    | 0.058   |
| 8     | 4092   | 0.051    | 0.051   |
| 9     | 3650   | 0.045    | 0.046   |

## Conclusions

## Conclusions and Future Investigations

- See many different systems exhibit Benford behavior.
- Ingredients of proofs (logarithms, equidistribution).
- Applications to fraud detection / data integrity.