

Large Subsets are Sumsets

Benjamin Baily
bmb2@williams.edu

This work was done jointly with the SMALL REU in 2015 and 2021.

With Justine Dell, Sophia Dever, Adam Dionne, Henry Fleischmann, Leo Goldmakher, Gal Gross, Faye Jackson, Steven J. Miller, Ethan Pesikoff, Huy Tuan Pham, Luke Reifenberg, and Vidya Venkatesh.

Williams College

May 24, 2022

Set Addition

Definition

Given two sets $A, B \subseteq \mathbb{Z}$, we say that $A + B = \{a + b \mid a \in A, b \in B\}$.

Set Addition

Definition

Given two sets $A, B \subseteq \mathbb{Z}$, we say that $A + B = \{a + b \mid a \in A, b \in B\}$.

Example

$$\{0, 1, 2\} + \{0, 1, 2, 4\} = \{0, 1, 2, 3, 4, 5, 6\}$$

Set Addition

Definition

Given two sets $A, B \subseteq \mathbb{Z}$, we say that $A + B = \{a + b \mid a \in A, b \in B\}$.

Example

$$\{0, 1, 2\} + \{0, 1, 2, 4\} = \{0, 1, 2, 3, 4, 5, 6\}$$

| | 0 | 1 | 2 |
|---|---|---|---|
| 0 | 0 | 1 | 2 |
| 1 | 1 | 2 | 3 |
| 2 | 2 | 3 | 4 |
| 4 | 4 | 5 | 6 |

Irreducibility

Definition

A set $S \subseteq \mathbb{Z}$ is **reducible** if $S = A + B$ for two sets A, B such that $|A|, |B| \geq 2$. Otherwise, S is **irreducible**.

Irreducibility

Definition

A set $S \subseteq \mathbb{Z}$ is **reducible** if $S = A + B$ for two sets A, B such that $|A|, |B| \geq 2$. Otherwise, S is **irreducible**.

Example

For any set $S \subset \mathbb{Z}$ and $n \in \mathbb{Z}$, we have $S = (S + \{-n\}) + \{n\}$, so it's important to require $|A|, |B| \geq 2$.

Irreducibility

Definition

A set $S \subseteq \mathbb{Z}$ is **reducible** if $S = A + B$ for two sets A, B such that $|A|, |B| \geq 2$. Otherwise, S is **irreducible**.

Example

For any set $S \subset \mathbb{Z}$ and $n \in \mathbb{Z}$, we have $S = (S + \{-n\}) + \{n\}$, so it's important to require $|A|, |B| \geq 2$.

Example

The set $\{0, 1, 2\} = \{0, 1\} + \{0, 1\}$ is reducible.

Irreducibility

Definition

A set $S \subseteq \mathbb{Z}$ is **reducible** if $S = A + B$ for two sets A, B such that $|A|, |B| \geq 2$. Otherwise, S is **irreducible**.

Example

For any set $S \subset \mathbb{Z}$ and $n \in \mathbb{Z}$, we have $S = (S + \{-n\}) + \{n\}$, so it's important to require $|A|, |B| \geq 2$.

Example

The set $\{0, 1, 2\} = \{0, 1\} + \{0, 1\}$ is reducible. In contrast, $\{0, 1, 3\}$ is irreducible.

Higher-Dimensional Irreducibility

Definition

Let $S \subset \mathbb{Z}^d$. S is **reducible** if $S = A + B$ for $A, B \subset \mathbb{Z}^d$ with $|A|, |B| \geq 2$. Otherwise, S is **irreducible**.

Higher-Dimensional Irreducibility

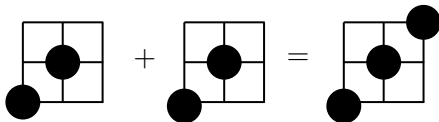
Definition

Let $S \subset \mathbb{Z}^d$. S is **reducible** if $S = A + B$ for $A, B \subset \mathbb{Z}^d$ with $|A|, |B| \geq 2$. Otherwise, S is **irreducible**.

Let $S = \{(0, 0), (1, 1), (2, 2)\}$. Then,

$$S = \{(0, 0), (1, 1)\} + \{(0, 0), (1, 1)\}$$

so S is reducible.



Results

Definition

Let $[n]^d = \underbrace{\{0, 1, \dots, n\} \times \{0, 1, \dots, n\} \cdots \{0, 1, \dots, n\}}_{d \text{ copies}}.$

Results

Definition

Let $[n]^d = \underbrace{\{0, 1, \dots, n\} \times \{0, 1, \dots, n\} \cdots \{0, 1, \dots, n\}}_{d \text{ copies}}.$

For example, $[5] = \{0, 1, 2, 3, 4, 5\}.$

Results

Definition

Let $[n]^d = \underbrace{\{0, 1, \dots, n\} \times \{0, 1, \dots, n\} \cdots \{0, 1, \dots, n\}}_{d \text{ copies}}.$

For example, $[5] = \{0, 1, 2, 3, 4, 5\}.$

Definition

For $d, n \geq 1$ we define $f_d(n)$ as follows.

$$f_d(n) = \min_{S \subset [n]^d, S \text{ irreducible}} |[n]^d \setminus S|$$

Results

Definition

Let $[n]^d = \underbrace{\{0, 1, \dots, n\} \times \{0, 1, \dots, n\} \cdots \{0, 1, \dots, n\}}_{d \text{ copies}}.$

For example, $[5] = \{0, 1, 2, 3, 4, 5\}.$

Definition

For $d, n \geq 1$ we define $f_d(n)$ as follows.

$$f_d(n) = \min_{S \subset [n]^d, S \text{ irreducible}} |[n]^d \setminus S|$$

Theorem

For all d , we have $f_d(n) = \Theta(\log n).$

What is the largest size of an irreducible subset of $[n]^d$?

Lemma

Fix $S \subset \mathbb{Z}^d$ such that $|S| \geq 3$ and $0 \in S$. Let $A = \{0, r\}$ for $r \in \mathbb{Z}^d \setminus \{0\}$. $S = A + B$ for some $B \subset S$ iff for all $s \in S$, $s - r \in S$ or $s + r \in S$.

What is the largest size of an irreducible subset of $[n]^d$?

Lemma

Fix $S \subset \mathbb{Z}^d$ such that $|S| \geq 3$ and $0 \in S$. Let $A = \{0, r\}$ for $r \in \mathbb{Z}^d \setminus \{0\}$. $S = A + B$ for some $B \subset S$ iff for all $s \in S$, $s - r \in S$ or $s + r \in S$.

- We consider the minimum size of the complement of an irreducible subset of $[n]^d$.

What is the largest size of an irreducible subset of $[n]^d$?

Lemma

Fix $S \subset \mathbb{Z}^d$ such that $|S| \geq 3$ and $0 \in S$. Let $A = \{0, r\}$ for $r \in \mathbb{Z}^d \setminus \{0\}$. $S = A + B$ for some $B \subset S$ iff for all $s \in S$, $s - r \in S$ or $s + r \in S$.

- We consider the minimum size of the complement of an irreducible subset of $[n]^d$.
- For each $r \in [n]^d$, in order for the set $\{0, r\}$ to not be a summand of S , there must exist $s \in S$ such that both $r - s, r + s \notin S$.

What is the largest size of an irreducible subset of $[n]^d$?

Lemma

Fix $S \subset \mathbb{Z}^d$ such that $|S| \geq 3$ and $0 \in S$. Let $A = \{0, r\}$ for $r \in \mathbb{Z}^d \setminus \{0\}$. $S = A + B$ for some $B \subset S$ iff for all $s \in S$, $s - r \in S$ or $s + r \in S$.

- We consider the minimum size of the complement of an irreducible subset of $[n]^d$.
- For each $r \in [n]^d$, in order for the set $\{0, r\}$ to not be a summand of S , there must exist $s \in S$ such that both $r - s, r + s \notin S$.
- If $|[n]^d \setminus S| \ll \log n$, the complement of S is too small for this to hold for all $r \in [n]^d$.

The Largest Irreducible Subset of $[n]^d$

Theorem

Let $S \subseteq [n]^d$. Let $k = |[n]^d \setminus S|$. Then, S is reducible if

$$\frac{k}{d} \log 2 + H_{\lceil \frac{k}{d} \rceil} + H_{\lceil \binom{k}{2}/d \rceil} < H_{n-1}$$

where H_n is the n th Harmonic number ($H_n \approx \log(n)$).

The Largest Irreducible Subset of $[n]^d$

Theorem

Let $S \subseteq [n]^d$. Let $k = |[n]^d \setminus S|$. Then, S is reducible if

$$\frac{k}{d} \log 2 + H_{\lceil \frac{k}{d} \rceil} + H_{\lceil \binom{k}{2}/d \rceil} < H_{n-1}$$

where H_n is the n th Harmonic number ($H_n \approx \log(n)$).

Corollary

We have $f_d(n) = \Omega(\log n)$.

How do you show a set is irreducible?

K. H. Kim and F. W. Roush showed in 2007 that the problem of determining if a set is irreducible is NP-complete.

How do you show a set is irreducible?

K. H. Kim and F. W. Roush showed in 2007 that the problem of determining if a set is irreducible is NP-complete.

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① If m is the smallest nonzero element of S , then $2m \notin S$.*
- ② For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.*

Then, S is irreducible.

How do you show a set is irreducible?

K. H. Kim and F. W. Roush showed in 2007 that the problem of determining if a set is irreducible is NP-complete.

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① If m is the smallest nonzero element of S , then $2m \notin S$.*
- ② For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.*

Then, S is irreducible.

- Local irreducibility is defined to be easily computer verifiable.

How do you show a set is irreducible?

K. H. Kim and F. W. Roush showed in 2007 that the problem of determining if a set is irreducible is NP-complete.

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① *If m is the smallest nonzero element of S , then $2m \notin S$.*
- ② *For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.*

Then, S is irreducible.

- Local irreducibility is defined to be easily computer verifiable.
- If $S \subset [n]$, then verification takes $O(\max(|S|^2, ([n] \setminus S)^3))$ time.

How do you show a set is irreducible?

K. H. Kim and F. W. Roush showed in 2007 that the problem of determining if a set is irreducible is NP-complete.

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① *If m is the smallest nonzero element of S , then $2m \notin S$.*
- ② *For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.*

Then, S is irreducible.

- Local irreducibility is defined to be easily computer verifiable.
- If $S \subset [n]$, then verification takes $O(\max(|S|^2, ([n] \setminus S)^3))$ time.
- Irreducibility follows by an iterative argument.

Local Irreducibility Example

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① If m is the smallest nonzero element of S , then $2m \notin S$.
- ② For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.

Then, S is irreducible.

Suppose $S = \{0, 1, 3\} = A + B$.

- WLOG, $0 \in A, B$ and $A, B \subset S$ by shifting the sets.
- 1 must be in A or B . Suppose $1 \in B$.
- Then, since $1 + 1 \notin S$, $1 \notin A$.
- Since $1 + 3 \notin S$, we know $3 \notin A$.
- So, $A = \{0\}$ and S is irreducible.

$$A = \{\} \quad B = \{\}$$

Local Irreducibility Example

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① If m is the smallest nonzero element of S , then $2m \notin S$.
- ② For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.

Then, S is irreducible.

Suppose $S = \{0, 1, 3\} = A + B$.

- WLOG, $0 \in A, B$ and $A, B \subset S$ by shifting the sets.
- 1 must be in A or B . Suppose $1 \in B$.
- Then, since $1 + 1 \notin S$, $1 \notin A$.
- Since $1 + 3 \notin S$, we know $3 \notin A$.
- So, $A = \{0\}$ and S is irreducible.

$$A = \{0\} \quad B = \{0\}$$

Local Irreducibility Example

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① *If m is the smallest nonzero element of S , then $2m \notin S$.*
- ② *For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.*

Then, S is irreducible.

Suppose $S = \{0, 1, 3\} = A + B$.

- WLOG, $0 \in A, B$ and $A, B \subset S$ by shifting the sets.
- 1 must be in A or B . Suppose $1 \in B$.
- **Then, since $1 + 1 \notin S$, $1 \notin A$.**
- Since $1 + 3 \notin S$, we know $3 \notin A$.
- So, $A = \{0\}$ and S is irreducible.

$$A = \{0\} \quad B = \{0, 1\}$$

Local Irreducibility Example

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① If m is the smallest nonzero element of S , then $2m \notin S$.
- ② For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.

Then, S is irreducible.

Suppose $S = \{0, 1, 3\} = A + B$.

- WLOG, $0 \in A, B$ and $A, B \subset S$ by shifting the sets.
- 1 must be in A or B . Suppose $1 \in B$.
- Then, since $1 + 1 \notin S$, $1 \notin A$.
- Since $1 + 3 \notin S$, we know $3 \notin A$.
- So, $A = \{0\}$ and S is irreducible.

$$A = \{0\} \quad B = \{0, 1, 3\}$$

Local Irreducibility Example

Proposition (Local Irreducibility)

Suppose $S \subset \mathbb{Z}^{\geq 0}$ with $0 \in S$ satisfies the following.

- ① If m is the smallest nonzero element of S , then $2m \notin S$.
- ② For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t < s$ and $s + t \notin S$.

Then, S is irreducible.

Suppose $S = \{0, 1, 3\} = A + B$.

- WLOG, $0 \in A, B$ and $A, B \subset S$ by shifting the sets.
- 1 must be in A or B . Suppose $1 \in B$.
- Then, since $1 + 1 \notin S$, $1 \notin A$.
- Since $1 + 3 \notin S$, we know $3 \notin A$.
- So, $A = \{0\}$ and S is irreducible.

$$A = \{0\} \quad B = \{0, 1, 3\}$$

Local Irreducibility in Higher Dimensions

Proposition (Local Irreducibility)

Suppose a set $S \subset \mathbb{Z}_{\geq 0}^d$ with $0 \in S$ satisfies the following.

- ① If m is the lexicographically first nonzero element of S , then $2m \notin S$.*
- ② For each $s \in S \setminus \{0, m\}$ there is some $t \in S$ with $t \prec s$ and $s + t \notin S$.*

Then, S is irreducible.

Constructive Bounds

- We construct explicit families of large locally irreducible subsets of $\mathbb{Z}^{\geq 0}$

Constructive Bounds

- We construct explicit families of large locally irreducible subsets of $\mathbb{Z}^{\geq 0}$
- Use pseudo-greedy construction algorithm to construct the 1-dimensional set

Constructive Bounds

- We construct explicit families of large locally irreducible subsets of $\mathbb{Z}^{\geq 0}$
- Use pseudo-greedy construction algorithm to construct the 1-dimensional set
- In higher dimensions, place a copy of the 1-dimensional set along each axis of the cube

Constructive Bounds

- We construct explicit families of large locally irreducible subsets of $\mathbb{Z}^{\geq 0}$
- Use pseudo-greedy construction algorithm to construct the 1-dimensional set
- In higher dimensions, place a copy of the 1-dimensional set along each axis of the cube

Proposition

For all n, d there exists an irreducible subset $S \subset [n]^d$ such that $|[n]^d \setminus S| = O(\log n)$.

Constructive Bounds

- We construct explicit families of large locally irreducible subsets of $\mathbb{Z}^{\geq 0}$
- Use pseudo-greedy construction algorithm to construct the 1-dimensional set
- In higher dimensions, place a copy of the 1-dimensional set along each axis of the cube

Proposition

For all n, d there exists an irreducible subset $S \subset [n]^d$ such that $|[n]^d \setminus S| = O(\log n)$.

Corollary

We have $f_d(n) = \Theta(\log n)$.

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i
- Find the largest number N such that $\mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$ is locally irreducible

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i
- Find the largest number N such that $\mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$ is locally irreducible
- Set $\mathcal{Q}_{i+1} = \mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i
- Find the largest number N such that $\mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$ is locally irreducible
- Set $\mathcal{Q}_{i+1} = \mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$

$$\mathcal{Q}_0 = \{0, 1\}$$

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i
- Find the largest number N such that $\mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$ is locally irreducible
- Set $\mathcal{Q}_{i+1} = \mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$

$$\mathcal{Q}_1 = \{0, 1, 3\}$$

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i
- Find the largest number N such that $\mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$ is locally irreducible
- Set $\mathcal{Q}_{i+1} = \mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$

$$\mathcal{Q}_2 = \{0, 1, 3, 5, 6, 7\}$$

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i
- Find the largest number N such that $\mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$ is locally irreducible
- Set $\mathcal{Q}_{i+1} = \mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$

$$\mathcal{Q}_3 = \{0, 1, 3, 5, 6, 7, 9, 10\}$$

1-Dimensional Construction

- Start with the set $\mathcal{Q}_0 = \{0, 1\}$
- After stage i of the process, \mathcal{Q}_i has maximum element M_i
- Find the largest number N such that $\mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$ is locally irreducible
- Set $\mathcal{Q}_{i+1} = \mathcal{Q}_i \cup \{M_i + 1, \dots, N - 1\}$

$$\mathcal{Q}_4 = \{0, 1, 3, 5, 6, 7, 9, 10, 12, 13, 14, 15\}$$

1-Dimensional Construction (Continued)

- The set $\mathcal{Q} := \bigcup_{i=1}^{\infty} \mathcal{Q}_i$ is locally irreducible by construction.

1-Dimensional Construction (Continued)

- The set $\mathcal{Q} := \bigcup_{i=1}^{\infty} \mathcal{Q}_i$ is locally irreducible by construction.
- For each n , the set $[n] \cap \mathcal{Q}$ is also locally irreducible.

1-Dimensional Construction (Continued)

- The set $\mathcal{Q} := \bigcup_{i=1}^{\infty} \mathcal{Q}_i$ is locally irreducible by construction.
- For each n , the set $[n] \cap \mathcal{Q}$ is also locally irreducible.
- **Problem:** Hard to estimate the size of $[n] \cap \mathcal{Q}$.

1-Dimensional Construction (Continued)

- The set $\mathcal{Q} := \bigcup_{i=1}^{\infty} \mathcal{Q}_i$ is locally irreducible by construction.
- For each n , the set $[n] \cap \mathcal{Q}$ is also locally irreducible.
- **Problem:** Hard to estimate the size of $[n] \cap \mathcal{Q}$.
- **Solution:** After constructing \mathcal{Q}_9 , change the construction.

1-Dimensional Construction (Continued)

- The set $\mathcal{Q} := \bigcup_{i=1}^{\infty} \mathcal{Q}_i$ is locally irreducible by construction.
- For each n , the set $[n] \cap \mathcal{Q}$ is also locally irreducible.
- **Problem:** Hard to estimate the size of $[n] \cap \mathcal{Q}$.
- **Solution:** After constructing \mathcal{Q}_9 , change the construction.
- For $0 \leq i \leq 9$, set $\mathcal{P}_i := \mathcal{Q}_i$, and for each $0 \leq i \leq 9$, set $N_i = \max(\mathcal{P}_i) + 1$.

1-Dimensional Construction (Continued)

- The set $\mathcal{Q} := \bigcup_{i=1}^{\infty} \mathcal{Q}_i$ is locally irreducible by construction.
- For each n , the set $[n] \cap \mathcal{Q}$ is also locally irreducible.
- **Problem:** Hard to estimate the size of $[n] \cap \mathcal{Q}$.
- **Solution:** After constructing \mathcal{Q}_9 , change the construction.
- For $0 \leq i \leq 9$, set $\mathcal{P}_i := \mathcal{Q}_i$, and for each $0 \leq i \leq 9$, set $N_i = \max(\mathcal{P}_i) + 1$.
- For $i \geq 10$, set $N_i = N_{i-2} + N_{i-3}$ and let

$$\mathcal{P}_i = \mathcal{P}_{i-1} \cup \{N_{i-1} + 1, \dots, N_i - 1\}$$

1-Dimensional Results

- Let $\mathcal{P} := \bigcup_{i=1}^{\infty} \mathcal{P}_i$.

1-Dimensional Results

- Let $\mathcal{P} := \bigcup_{i=1}^{\infty} \mathcal{P}_i$.
- For each n , the set $[n] \cap \mathcal{P}$ is locally irreducible and its complement in $[n]$ is (up to finitely many edits) a linear recurrence sequence.

1-Dimensional Results

- Let $\mathcal{P} := \bigcup_{i=1}^{\infty} \mathcal{P}_i$.
- For each n , the set $[n] \cap \mathcal{P}$ is locally irreducible and its complement in $[n]$ is (up to finitely many edits) a linear recurrence sequence.
- We are able to estimate the number of terms in a linear recurrence sequence up to n .

1-Dimensional Results

- Let $\mathcal{P} := \bigcup_{i=1}^{\infty} \mathcal{P}_i$.
- For each n , the set $[n] \cap \mathcal{P}$ is locally irreducible and its complement in $[n]$ is (up to finitely many edits) a linear recurrence sequence.
- We are able to estimate the number of terms in a linear recurrence sequence up to n .
- Let λ denote the largest complex root of the polynomial $x^3 - x - 1 = 0$.

Proposition

Let $\lambda = 1.325\dots$ denote the largest complex root of the polynomial $x^3 - x - 1 = 0$. Then $|[n] \setminus \mathcal{P}| \sim \log_{\lambda} n$.

Recap

- 1 We let $f_d(n)$ denote the minimum size complement of an irreducible subset of $[n]^d$.
- 2 By showing that every sufficiently large subset $S \subset [n]^d$ admits a factorization of the form $S = A + B$ with $A = \{0, r\}$, we proved that $f_d(n) = \Omega(\log n)$.
- 3 Define **local irreducibility**, a strong condition on a subset of \mathbb{Z}^d such that every locally irreducible set is also irreducible.
- 4 Construct large locally irreducible subsets of $[n]^d$ to prove that $f_d(n) = O(\log n)$.
- 5 Conclude that $f_d(n) = \Theta(\log n)$.

Acknowledgments

This research was done as part of the SMALL REU program and was funded by NSF grant number 1947438.

Special thanks to Professors Leo Goldmakher and Steven J. Miller for their mentorship.



References



B. Baily, J. Dell, S. Dever, A. Dionne, H. Fleischmann, L. Goldmakher, G. Gross, F. Jackson, S. J. Miller, E. Pesikoff, H. Pham, L. Reifenberg, and V. Venkatesh, Large Subsets are Sumsets, in preparation.



<https://oeis.org/A349775>