

# Cookie Monster Meets the Fibonacci Numbers. Mmmmmm – Theorems!

Steven J Miller, Williams College

Steven.J.Miller@williams.edu

<http://www.williams.edu/go/math/sjmilller>

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## Summary / Acknowledgements



- **Previous results:** Zeckendorf and Lekkerkerker.
- **New approach:** Joint with Carlos Dominguez, Gene Kopp, Murat Kolğlu and Yinghui Wang.
- **Thanks:** Ed Burger and his SMALL REU students (David Clyde, Cory Colbert, Gea Shin and Nancy

## Previous Results

**Fibonacci Numbers:**  $F_{n+1} = F_n + F_{n-1}$ ;

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### Lekkerkerker's Theorem

The average number of non-consecutive Fibonacci summands in the Zeckendorf decomposition for integers in  $[F_n, F_{n+1})$  tends to  $\frac{n}{\varphi^2 + 1} \approx .276n$ , where  $\varphi = \frac{1+\sqrt{5}}{2}$  is the golden mean.

## Main Results

### Lemma: Application of Cookie Counting

The 'probability' (ie, percentage of the time) an integer in  $[F_n, F_{n+1})$  has exactly  $k + 1$  non-consecutive Fibonacci summands is  $\binom{n-1-k}{k} / F_{n-1}$ .

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The above lemma yields Zeckendorf's Theorem, Lekkerkerker's Theorem, and

### An Erdos-Kac Type Theorem: SMALL 2010

As  $n \rightarrow \infty$ , the distribution of the number of non-consecutive Fibonacci summands in the Zeckendorf decomposition for integers in  $[F_n, F_{n+1})$  is Gaussian.

## Properties of Fibonacci Numbers and needed Combinatorial Results

## Binet's Formula

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$$F_n = \frac{1}{\sqrt{5}} \left( \frac{1 + \sqrt{5}}{2} \right)^{n+1} - \frac{1}{\sqrt{5}} \left( \frac{1 - \sqrt{5}}{2} \right)^{n+1}.$$

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Alternate proof via generating functions useful for

## Combinatorial Review

### The Cookie Problem

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**Example:** 10 cookies and 5 people:



## Cookie Problem: Reinterpretation

### Reinterpreting the Cookie Problem

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This follows by setting  $x_i = y_i + c_i$  with  $y_i$  a non-negative integer.

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Must have  $F_n$  one of the summands, must not have  $F_{n-1}$ .

For each Fibonacci number from  $F_1$  to  $F_{n-1}$  we either include or not, cannot have two consecutive, must end with a non-taken number.

## Proof of Zeckendorf's Theorem (continued)

Consider all subsets of  $k + 1$  non-consecutive Fibonacci from  $\{F_1, \dots, F_n\}$  where  $F_n$  is taken. Let  $y_0$  be number of Fibonacci not taken until first one taken, and then  $y_i$  ( $1 \leq i \leq k$ ) be the number not taken between two taken.

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**Example:**  $2010 = 1597 + 377 + 34 + 2 = F_{16} + F_{13} + F_8 + F_2$ ,  
so  $n = 16$ ,  $k + 1 = 4$ ,  $y_0 = 1$ ,  $y_1 = 5$ ,  $y_2 = 4$ ,  $y_3 = 2$ .

Equivalently:  $y_0 + y_1 + \dots + y_k + k = n - 1$ ,  $y_i \geq 1$  if  $i \geq 1$ .

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Obtain  $\sum_{k=0}^{\lfloor \frac{n-1}{2} \rfloor} \binom{n-1-k}{k} = F_{n-1}$  integers in  $[F_n, F_{n+1})$ ; as all distinct and this many integers in interval, done.  $\square$

## Lekkerkerker's Theorem

## Preliminaries

$$\mathcal{E}(n) := \sum_{k=0}^{\lfloor \frac{n-1}{2} \rfloor} k \binom{n-1-k}{k}.$$

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Proof by algebra (details in appendix):

$$\begin{aligned}\mathcal{E}(n) &= \sum_{k=0}^{\lfloor \frac{n-1}{2} \rfloor} k \binom{n-1-k}{k} \\ &= (n-2) \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} \binom{n-3-\ell}{\ell} - \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} \ell \binom{n-3-\ell}{\ell} \\ &= (n-2)F_{n-3} - \mathcal{E}(n-2).\end{aligned}$$

## Solving Recurrence Relation

### Formula for $\mathcal{E}(n)$ (i.e., Lekkerkerker's Theorem)

$$\mathcal{E}(n) = \frac{nF_{n-1}}{\varphi^2 + 1} + O(F_{n-2}).$$

$$\begin{aligned} & \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-1)^\ell (\mathcal{E}(n-2\ell) + \mathcal{E}(n-2(\ell+1))) \\ &= \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-1)^\ell (n-2-2\ell)F_{n-3-2\ell}. \end{aligned}$$

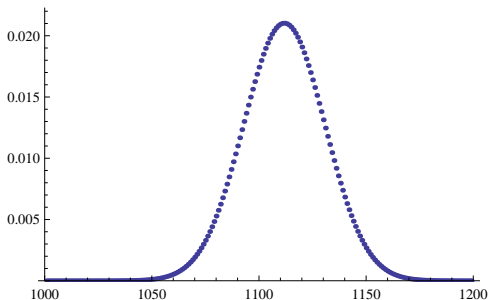
Result follows from Binet's formula, the geometric series formula, and differentiating identities:  $\sum_{j=0}^m jx^j = x \frac{(m+1)x^m(x-1) - (x^{m+1}-1)}{(x-1)^2}$ . Details in appendix.

## An Erdos-Kac Type Theorem

## Generalizing Lekkerkerker

### Theorem (SMALL 2010)

As  $n \rightarrow \infty$ , the distribution of the number of summands in Zeckendorf's Theorem is a Gaussian.



**Figure:** Number of summands in  $[F_{2010}, F_{2011})$

## Generalizing Lekkerkerker: Erdos-Kac type result

### Theorem (SMALL 2010)

As  $n \rightarrow \infty$ , the distribution of the number of summands in Zeckendorf's Theorem is a Gaussian.

**Numerics:** At  $F_{100,000}$ : Ratio of  $2m^{\text{th}}$  moment  $\sigma_{2m}$  to  $(2m-1)!!\sigma_2^m$  is between .999955 and 1 for  $2m \leq 10$ .

**Sketch of proof:** Use Stirling's formula,

$$n! \approx n^n e^{-n} \sqrt{2\pi n}$$

to approximate binomial coefficients, after a few pages of algebra find the probabilities are approximately Gaussian.

## Additional Generalizations

## Further Generalization

### Generalized Fibonacci Numbers

Let  $H_n = c_1 H_{n-1} + \cdots + c_L H_{n-L}$  with  $c_1 \geq \cdots \geq c_L \geq 1$ . Then every positive integer can be written as a unique sum of the  $H_i$ 's such that cannot use the recurrence relation to remove any summands.

*Key ingredients in proof: generating functions, matching coefficients of polynomials.*

## Further Generalization (cont)

In 2009 Hannah Alpert proved every positive integer can be written uniquely as a sum and difference of Fibonacci numbers, such that all terms of the same sign are at least 4 apart and those of different sign at least 3. We can show

### Signed Representations

The number of positive and negative summands are Gaussianly distributed as  $n \rightarrow \infty$ . They are not independent, and have a negative correlation coefficient.

## Conclusion

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- Re-derive Zeckendorf and Lekkerkerker's results through combinatorics.
- Method yields an Erdos-Kac type result on Gaussian behavior of the number of summands.
- Method applicable to other, related questions.

NOTE: These and similar questions are being studied by the students at the 2010 SMALL REU at Williams College; we expect to be able to provide papers and proofs by the end of the summer.

## Appendix: Details of Computations

## Needed Binomial Identity

### Binomial identity involving Fibonacci Numbers

Let  $F_m$  denote the  $m^{\text{th}}$  Fibonacci number, with  $F_1 = 1$ ,  $F_2 = 2$ ,  $F_3 = 3$ ,  $F_4 = 5$  and so on. Then

$$\sum_{k=0}^{\lfloor \frac{n-1}{2} \rfloor} \binom{n-1-k}{k} = F_{n-1}.$$

**Proof by induction:** The base case is trivially verified. Assume our claim holds for  $n$  and show that it holds for  $n+1$ . We may extend the sum to  $n-1$ , as  $\binom{n-1-k}{k} = 0$  whenever  $k > \lfloor \frac{n-1}{2} \rfloor$ . Using the standard identity that

$$\binom{m}{\ell} + \binom{m}{\ell+1} = \binom{m+1}{\ell+1},$$

and the convention that  $\binom{m}{\ell} = 0$  if  $\ell$  is a negative integer, we find

$$\begin{aligned} \sum_{k=0}^n \binom{n-k}{k} &= \sum_{k=0}^n \left[ \binom{n-1-k}{k-1} + \binom{n-1-k}{k} \right] \\ &= \sum_{k=1}^n \binom{n-1-k}{k-1} + \sum_{k=0}^n \binom{n-1-k}{k} \\ &= \sum_{k=1}^n \binom{n-2-(k-1)}{k-1} + \sum_{k=0}^n \binom{n-1-k}{k} = F_{n-2} + F_{n-1} \end{aligned}$$

by the inductive assumption; noting  $F_{n-2} + F_{n-1} = F_n$  completes the proof.  $\square$

# Derivation of Recurrence Relation for $\mathcal{E}(n)$

$$\begin{aligned}
 \mathcal{E}(n) &= \sum_{k=0}^{\lfloor \frac{n-1}{2} \rfloor} k \binom{n-1-k}{k} \\
 &= \sum_{k=1}^{\lfloor \frac{n-1}{2} \rfloor} k \frac{(n-1-k)!}{k!(n-1-2k)!} \\
 &= \sum_{k=1}^{\lfloor \frac{n-1}{2} \rfloor} (n-1-k) \frac{(n-2-k)!}{(k-1)!(n-1-2k)!} \\
 &= \sum_{k=1}^{\lfloor \frac{n-1}{2} \rfloor} (n-2-(k-1)) \frac{(n-3-(k-1))!}{(k-1)!(n-3-2(k-1))!} \\
 &= \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (n-2-\ell) \binom{n-3-\ell}{\ell} \\
 &= (n-2) \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} \binom{n-3-\ell}{\ell} - \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} \ell \binom{n-3-\ell}{\ell} \\
 &= (n-2)F_{n-3} - \mathcal{E}(n-2),
 \end{aligned}$$

which proves the claim (note we used the binomial identity to replace the sum of binomial coefficients with a Fibonacci number).

# Formula for $\mathcal{E}(n)$

## Formula for $\mathcal{E}(n)$

$$\mathcal{E}(n) = \frac{nF_{n-1}}{\varphi^2 + 1} + O(F_{n-2}).$$

**Proof:** The proof follows from using telescoping sums to get an expression for  $\mathcal{E}(n)$ , which is then evaluated by inputting Binet's formula and differentiating identities. Explicitly, consider

$$\begin{aligned} \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-1)^\ell (\mathcal{E}(n-2\ell) + \mathcal{E}(n-2(\ell+1))) &= \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-1)^\ell (n-2-2\ell)F_{n-3-2\ell} \\ &= \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-1)^\ell (n-3-2\ell)F_{n-3-2\ell} + \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-1)^\ell (2\ell)F_{n-3-2\ell} \\ &= \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-1)^\ell (n-3-2\ell)F_{n-3-2\ell} + O(F_{n-2}); \end{aligned}$$

while we could evaluate the last sum exactly, trivially estimating it suffices to obtain the main term (as we have a sum of every other Fibonacci number, the sum is at most the next Fibonacci number after the largest one in our sum).

## Formula for $\mathcal{E}(n)$ (continued)

We now use Binet's formula to convert the sum into a geometric series. Letting  $\varphi = \frac{1+\sqrt{5}}{2}$  be the golden mean, we have

$$F_n = \frac{\varphi}{\sqrt{5}} \cdot \varphi^n - \frac{1-\varphi}{\sqrt{5}} \cdot (1-\varphi)^n$$

(our constants are because our counting has  $F_1 = 1$ ,  $F_2 = 2$  and so on). As  $|1-\varphi| < 1$ , the error from dropping the  $(1-\varphi)^n$  term is  $O(\sum_{\ell \leq n} n) = O(n^2) = o(F_{n-2})$ , and may thus safely be absorbed in our error term. We thus find

$$\begin{aligned} \mathcal{E}(n) &= \frac{\varphi}{\sqrt{5}} \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (n-3-2\ell)(-1)^\ell \varphi^{n-3-2\ell} + O(F_{n-2}) \\ &= \frac{\varphi^{n-2}}{\sqrt{5}} \left[ (n-3) \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} (-\varphi^{-2})^\ell - 2 \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} \ell (-\varphi^{-2})^\ell \right] + O(F_{n-2}). \end{aligned}$$

## Formula for $\mathcal{E}(n)$ (continued)

We use the geometric series formula to evaluate the first term. We drop the upper boundary term of

$(-\varphi^{-1})^{\lfloor \frac{n-3}{2} \rfloor}$ , as this term is negligible since  $\varphi > 1$ . We may also move the 3 from the  $n-3$  into the error term, and are left with

$$\begin{aligned}\mathcal{E}(n) &= \frac{\varphi^{n-2}}{\sqrt{5}} \left[ \frac{n}{1+\varphi^{-2}} - 2 \sum_{\ell=0}^{\lfloor \frac{n-3}{2} \rfloor} \ell (-\varphi^{-2})^{\ell} \right] + O(F_{n-2}) \\ &= \frac{\varphi^{n-2}}{\sqrt{5}} \left[ \frac{n}{1+\varphi^{-2}} - 2S\left(\left\lfloor \frac{n-3}{2} \right\rfloor, -\varphi^{-2}\right) \right] + O(F_{n-2}),\end{aligned}$$

where

$$S(m, x) = \sum_{j=0}^m jx^j.$$

There is a simple formula for  $S(m, x)$ . As

$$\sum_{j=0}^m x^j = \frac{x^{m+1} - 1}{x - 1},$$

applying the operator  $x \frac{d}{dx}$  gives

$$S(m, x) = \sum_{j=0}^m jx^j = x \frac{(m+1)x^m(x-1) - (x^{m+1} - 1)}{(x-1)^2} = \frac{mx^{m+2} - (m+1)x^{m+1} + x}{(x-1)^2}.$$

## Formula for $\mathcal{E}(n)$ (continued)

Taking  $x = -\varphi^{-2}$ , we see that the contribution from this piece may safely be absorbed into the error term  $O(F_{n-2})$ , leaving us with

$$\mathcal{E}(n) = \frac{n\varphi^{n-2}}{\sqrt{5}(1+\varphi^{-2})} + O(F_{n-2}) = \frac{n\varphi^n}{\sqrt{5}(\varphi^2+1)} + O(F_{n-2}).$$

Noting that for large  $n$  we have  $F_{n-1} = \frac{\varphi^n}{\sqrt{5}} + O(1)$ , we finally obtain

$$\mathcal{E}(n) = \frac{nF_{n-1}}{\varphi^2+1} + O(F_{n-2}). \square$$

## (Sketch of the) Proof of Gaussianness

The probability density for the number of Fibonacci numbers that add up to an integer in  $[F_n, F_{n+1})$  is  $f_n(k) = \binom{n-1-k}{k} / F_{n-1}$ . Consider the density for the  $n+1$  case. Then we have, by Stirling

$$\begin{aligned} f_{n+1}(k) &= \binom{n-k}{k} \frac{1}{F_n} \\ &= \frac{(n-k)!}{(n-2k)!k!} \frac{1}{F_n} = \frac{1}{\sqrt{2\pi}} \frac{(n-k)^{(n-k+\frac{1}{2})}}{k^{(k+\frac{1}{2})}(n-2k)^{(n-2k+\frac{1}{2})}} \frac{1}{F_n} \end{aligned}$$

plus a lower order correction term.

Also we can write  $F_n = \frac{1}{\sqrt{5}} \phi^{n+1} = \frac{\phi}{\sqrt{5}} \phi^n$  for large  $n$ , where  $\phi$  is the golden ratio (we are using relabeled Fibonacci numbers where  $1 = F_1$  occurs once to help dealing with uniqueness and  $F_2 = 2$ ). We can now split the terms that exponentially depend on  $n$ .

$$f_{n+1}(k) = \left( \frac{1}{\sqrt{2\pi}} \sqrt{\frac{(n-k)}{k(n-2k)}} \frac{\sqrt{5}}{\phi} \right) \left( \phi^{-n} \frac{(n-k)^{(n-k)}}{k^k (n-2k)^{(n-2k)}} \right)$$

Define

$$N_n = \frac{1}{\sqrt{2\pi}} \sqrt{\frac{(n-k)}{k(n-2k)}} \frac{\sqrt{5}}{\phi}, \quad S_n = \phi^{-n} \frac{(n-k)^{(n-k)}}{k^k (n-2k)^{(n-2k)}}.$$

Thus, write the density function as

$$f_{n+1}(k) = N_n S_n$$

where  $N_n$  is the first term that is of order  $n^{-1/2}$  and  $S_n$  is the second term with exponential dependence on  $n$ .

## (Sketch of the) Proof of Gaussianity (cont)

Model the distribution as centered around the mean by the change of variable  $k = \mu + x\sigma$  where  $\mu$  and  $\sigma$  are the mean and the standard deviation, and depend on  $n$ . The discrete weights of  $f_n(k)$  will become continuous. This requires us to use the change of variable formula to compensate for the change of scales:

$$f_n(k)dk = f_n(\mu + \sigma x)\sigma dx.$$

Using the change of variable, we can write  $N_n$  as

$$\begin{aligned} N_n &= \frac{1}{\sqrt{2\pi}} \sqrt{\frac{n-k}{k(n-2k)}} \frac{\phi}{\sqrt{5}} \\ &= \frac{1}{\sqrt{2\pi n}} \sqrt{\frac{1-k/n}{(k/n)(1-2k/n)}} \frac{\sqrt{5}}{\phi} \\ &= \frac{1}{\sqrt{2\pi n}} \sqrt{\frac{1-(\mu+\sigma x)/n}{((\mu+\sigma x)/n)(1-2(\mu+\sigma x)/n)}} \frac{\sqrt{5}}{\phi} \\ &= \frac{1}{\sqrt{2\pi n}} \sqrt{\frac{1-C-y}{(C+y)(1-2C-2y)}} \frac{\sqrt{5}}{\phi} \end{aligned}$$

where  $C = \mu/n \approx 1/(\phi+2)$  (note that  $\phi^2 = \phi+1$ ) and  $y = \sigma x/n$ . But for large  $n$ , the  $y$  term vanishes since  $\sigma \sim \sqrt{n}$  and thus  $y \sim n^{-1/2}$ . Thus

$$N_n \approx \frac{1}{\sqrt{2\pi n}} \sqrt{\frac{1-C}{C(1-2C)}} \frac{\sqrt{5}}{\phi} = \frac{1}{\sqrt{2\pi n}} \sqrt{\frac{(\phi+1)(\phi+2)}{\phi}} \frac{\sqrt{5}}{\phi} = \frac{1}{\sqrt{2\pi n}} \sqrt{\frac{5(\phi+2)}{\phi}} = \frac{1}{\sqrt{2\pi\sigma^2}}$$

since  $\sigma^2 = n \frac{\phi}{5(\phi+2)}$ .

## (Sketch of the) Proof of Gaussianity (cont)

For the second term  $S_n$ , take the logarithm and once again change variable  $k = \mu + x\sigma$ ,

$$\begin{aligned}
 \log(S_n) &= \log \left( \phi^{-n} \frac{(n-k)^{(n-k)}}{k^k (n-2k)^{(n-2k)}} \right) \\
 &= -n \log(\phi) + (n-k) \log(n-k) - (k) \log(k) \\
 &\quad - (n-2k) \log(n-2k) \\
 &= -n \log(\phi) + (n - (\mu + x\sigma)) \log(n - (\mu + x\sigma)) \\
 &\quad - (\mu + x\sigma) \log(\mu + x\sigma) \\
 &\quad - (n - 2(\mu + x\sigma)) \log(n - 2(\mu + x\sigma)) \\
 &= -n \log(\phi) \\
 &\quad + (n - (\mu + x\sigma)) \left( \log(n - \mu) + \log \left( 1 - \frac{x\sigma}{n - \mu} \right) \right) \\
 &\quad - (\mu + x\sigma) \left( \log(\mu) + \log \left( 1 + \frac{x\sigma}{\mu} \right) \right) \\
 &\quad - (n - 2(\mu + x\sigma)) \left( \log(n - 2\mu) + \log \left( 1 - \frac{x\sigma}{n - 2\mu} \right) \right) \\
 &= -n \log(\phi) \\
 &\quad + (n - (\mu + x\sigma)) \left( \log \left( \frac{n}{\mu} - 1 \right) + \log \left( 1 - \frac{x\sigma}{n - \mu} \right) \right) \\
 &\quad - (\mu + x\sigma) \log \left( 1 + \frac{x\sigma}{\mu} \right) \\
 &\quad - (n - 2(\mu + x\sigma)) \left( \log \left( \frac{n}{\mu} - 2 \right) + \log \left( 1 - \frac{x\sigma}{n - 2\mu} \right) \right).
 \end{aligned}$$

## (Sketch of the) Proof of Gaussianity (cont)

Note that, since  $n/\mu = \phi + 2$  for large  $n$ , the constant terms vanish. We have  $\log(S_n)$

$$\begin{aligned}
 &= -n \log(\phi) + (n-k) \log\left(\frac{n}{\mu} - 1\right) - (n-2k) \log\left(\frac{n}{\mu} - 2\right) + (n-(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-\mu}\right) \\
 &\quad - (\mu+x\sigma) \log\left(1 + \frac{x\sigma}{\mu}\right) - (n-2(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-2\mu}\right) \\
 &= -n \log(\phi) + (n-k) \log(\phi+1) - (n-2k) \log(\phi) + (n-(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-\mu}\right) \\
 &\quad - (\mu+x\sigma) \log\left(1 + \frac{x\sigma}{\mu}\right) - (n-2(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-2\mu}\right) \\
 &= n(-\log(\phi) + \log(\phi^2) - \log(\phi)) + k(\log(\phi^2) + 2\log(\phi)) + (n-(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-\mu}\right) \\
 &\quad - (\mu+x\sigma) \log\left(1 + \frac{x\sigma}{\mu}\right) - (n-2(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-2\mu}\right) \\
 &= (n-(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-\mu}\right) - (\mu+x\sigma) \log\left(1 + \frac{x\sigma}{\mu}\right) \\
 &\quad - (n-2(\mu+x\sigma)) \log\left(1 - \frac{x\sigma}{n-2\mu}\right).
 \end{aligned}$$

## (Sketch of the) Proof of Gaussianity (cont)

Finally, we expand the logarithms and collect powers of  $x\sigma/n$ .  $\log(S_n)$

$$\begin{aligned}
 &= (n - (\mu + x\sigma)) \left( -\frac{x\sigma}{n - \mu} - \frac{1}{2} \left( \frac{x\sigma}{n - \mu} \right)^2 + \dots \right) \\
 &\quad - (\mu + x\sigma) \left( \frac{x\sigma}{\mu} - \frac{1}{2} \left( \frac{x\sigma}{\mu} \right)^2 + \dots \right) \\
 &\quad - (n - 2(\mu + x\sigma)) \left( -2\frac{x\sigma}{n - 2\mu} - \frac{1}{2} \left( 2\frac{x\sigma}{n - 2\mu} \right)^2 + \dots \right) \\
 &= (n - (\mu + x\sigma)) \left( -\frac{x\sigma}{n \frac{(\phi+1)}{(\phi+2)}} - \frac{1}{2} \left( \frac{x\sigma}{n \frac{(\phi+1)}{(\phi+2)}} \right)^2 + \dots \right) \\
 &\quad - (\mu + x\sigma) \left( \frac{x\sigma}{\frac{n}{\phi+2}} - \frac{1}{2} \left( \frac{x\sigma}{\frac{n}{\phi+2}} \right)^2 + \dots \right) \\
 &\quad - (n - 2(\mu + x\sigma)) \left( -\frac{2x\sigma}{n \frac{\phi}{\phi+2}} - \frac{1}{2} \left( \frac{2x\sigma}{n \frac{\phi}{\phi+2}} \right)^2 + \dots \right) \\
 &= \frac{x\sigma}{n} n \left( -\left(1 - \frac{1}{\phi+2}\right) \frac{(\phi+2)}{(\phi+1)} - 1 + 2\left(1 - \frac{2}{\phi+2}\right) \frac{\phi+2}{\phi} \right) \\
 &\quad - \frac{1}{2} \left( \frac{x\sigma}{n} \right)^2 n \left( -2\frac{\phi+2}{\phi+1} + \frac{\phi+2}{\phi+1} + 2(\phi+2) - (\phi+2) + 4\frac{\phi+2}{\phi} \right) \\
 &\quad + O(n(x\sigma/n)^3)
 \end{aligned}$$

## (Sketch of the) Proof of Gaussianity (cont)

$$\begin{aligned}
 &= \frac{x\sigma}{n} n \left( -\frac{\phi+1}{\phi+2} \frac{\phi+2}{\phi+1} - 1 + 2 \frac{\phi}{\phi+2} \frac{\phi+2}{\phi} \right) \\
 &\quad - \frac{1}{2} \left( \frac{x\sigma}{n} \right)^2 n(\phi+2) \left( -\frac{1}{\phi+1} + 1 + \frac{4}{\phi} \right) \\
 &\quad + O \left( n \left( \frac{x\sigma}{n} \right)^3 \right) \\
 &= -\frac{1}{2} \frac{(x\sigma)^2}{n} (\phi+2) \left( \frac{3\phi+4}{\phi(\phi+1)} + 1 \right) + O \left( n \left( \frac{x\sigma}{n} \right)^3 \right) \\
 &= -\frac{1}{2} \frac{(x\sigma)^2}{n} (\phi+2) \left( \frac{3\phi+4+2\phi+1}{\phi(\phi+1)} \right) + O \left( n \left( \frac{x\sigma}{n} \right)^3 \right) \\
 &= -\frac{1}{2} x^2 \sigma^2 \left( \frac{5(\phi+2)}{\phi n} \right) + O \left( n (x\sigma/n)^3 \right)
 \end{aligned}$$

## (Sketch of the) Proof of Gaussianity (cont)

But recall that

$$\sigma^2 = \frac{\phi n}{5(\phi + 2)}$$

Also, since  $\sigma \sim n^{-1/2}$ ,  $n \left( \frac{x\sigma}{n} \right)^3 \sim n^{-1/2}$ . So for large  $n$ , the  $O \left( n \left( \frac{x\sigma}{n} \right)^3 \right)$  term vanishes. Thus we are left with

$$\begin{aligned} \log S_n &= -\frac{1}{2}x^2 \\ S_n &= e^{-\frac{1}{2}x^2} \end{aligned}$$

Hence, as  $n$  gets large, the density converges to the normal distribution.

$$\begin{aligned} f_n(k)dk &= N_n S_n dk \\ &= \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{1}{2}x^2} \sigma dx \\ &= \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2}x^2} dx \end{aligned}$$

□

## References



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