

# Why I love Monovariants: From Zombies to Conway's Soldiers to Fibonacci Games

**Steven J. Miller: [sjm1@williams.edu](mailto:sjm1@williams.edu)  
President Fibonacci Association, Williams College**

[http://www.williams.edu/Mathematics/sjmiller/public\\_html](http://www.williams.edu/Mathematics/sjmiller/public_html)

Dartmouth, February 10, 2025



## Invariants / Monovariants

**Invariant:** a quantity that is unchanged throughout the process / operations. (Big application: Noether's theorem).

**Monovariant:** a quantity that only changes in one direction throughout the process / operations. See

<https://howardhalim.com/math/Invariants%20and%20Monovariants.pdf>

for a nice collection of problems.

Often a challenge to find a useful monovariant.

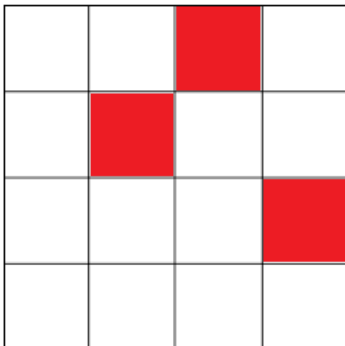
# Zombies

## Zombie Infection: Rules

- If share walls with 2 or more infected, become infected.
- Once infected, always infected.

## Zombie Infection: Rules

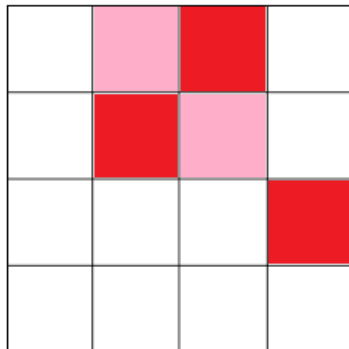
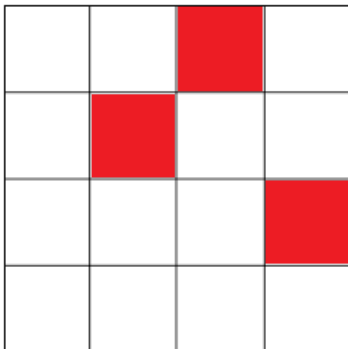
- If share walls with 2 or more infected, become infected.
- Once infected, always infected.



*Initial Configuration*

## Zombie Infection: Rules

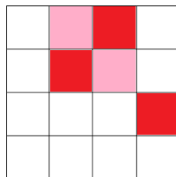
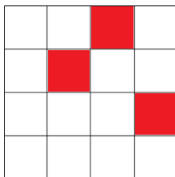
- If share walls with 2 or more infected, become infected.
- Once infected, always infected.



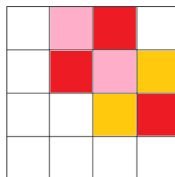
*Initial Configuration    One moment later*

## Zombie Infection: Rules

- If share walls with 2 or more infected, become infected.
- Once infected, always infected.



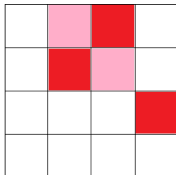
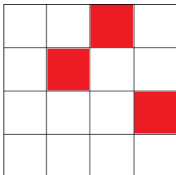
*Initial Configuration    One moment later*



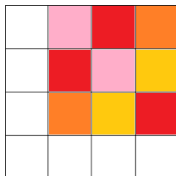
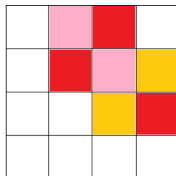
*Two moments later*

# Zombie Infection: Rules

- If share walls with 2 or more infected, become infected.
- Once infected, always infected.



*Initial Configuration One moment later*

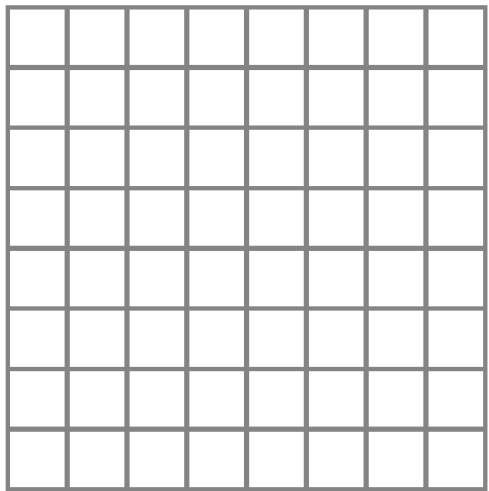


*Two moments later Three moments later*



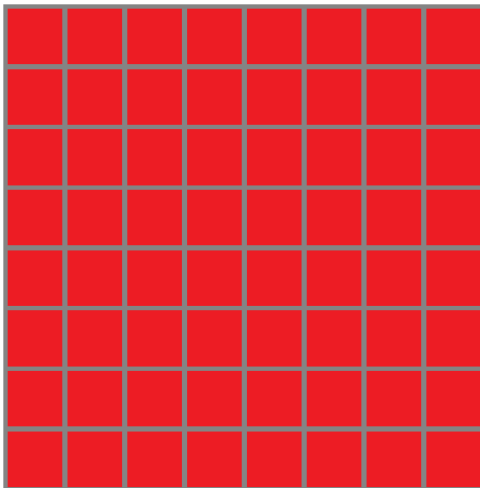
# Zombie Infection: Conquering The World

Easiest initial state that ensures all eventually infected is...?



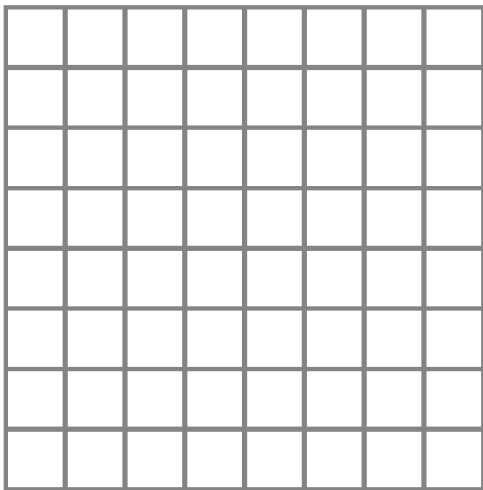
# Zombie Infection: Conquering The World

Easiest initial state that ensures all eventually infected is...?



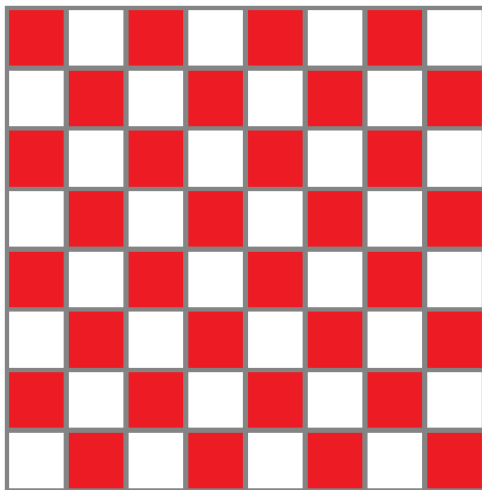
## Zombie Infection: Conquering The World

Next simplest initial state ensuring all eventually infected...?



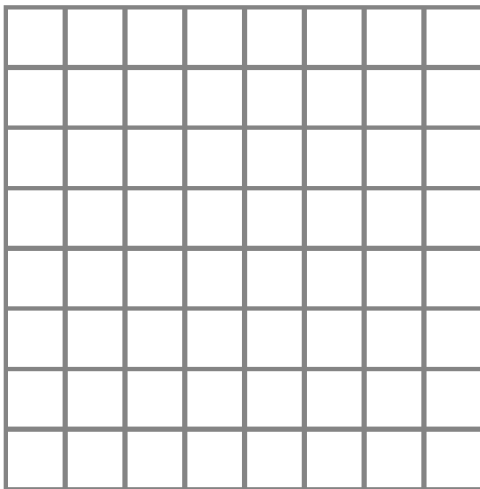
## Zombie Infection: Conquering The World

Next simplest initial state ensuring all eventually infected...?



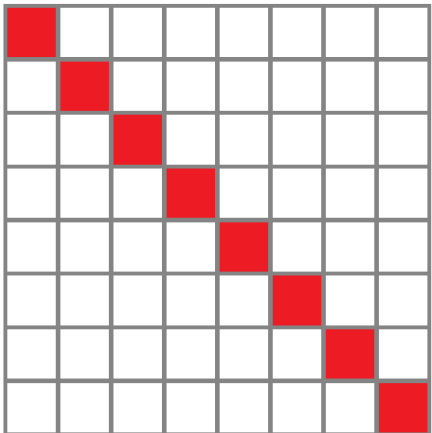
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



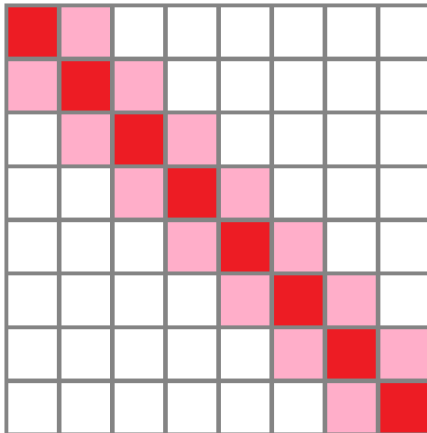
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



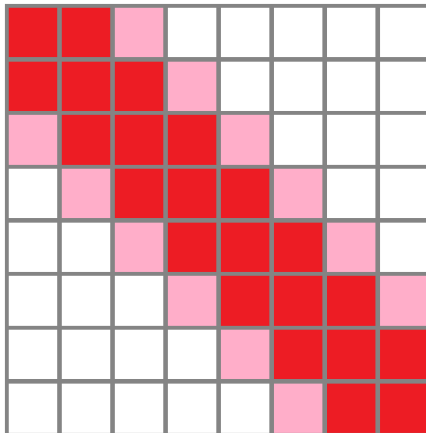
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



# Zombie Infection: Conquering The World

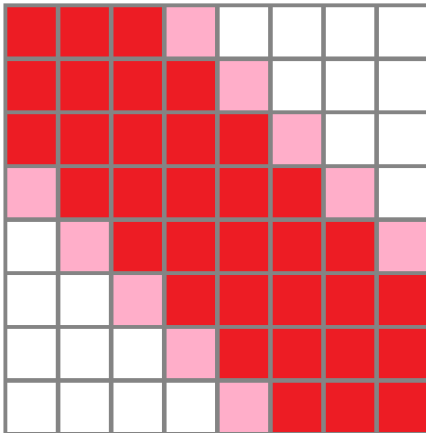
Fewest number of initial infections needed to get all...?





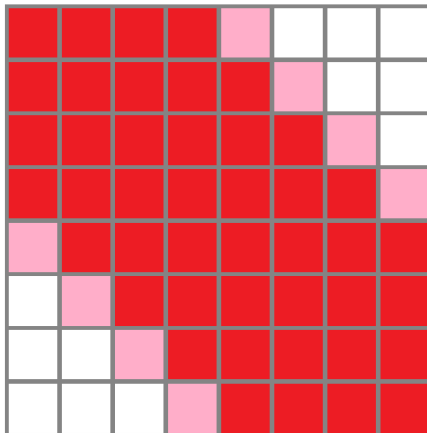
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



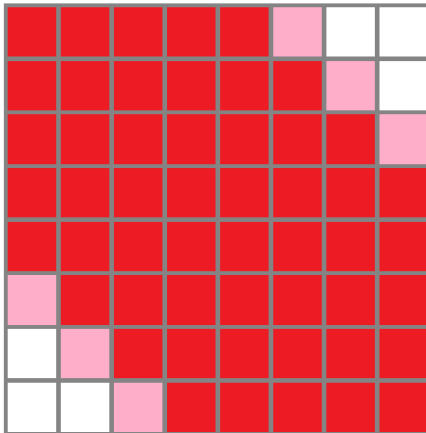
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



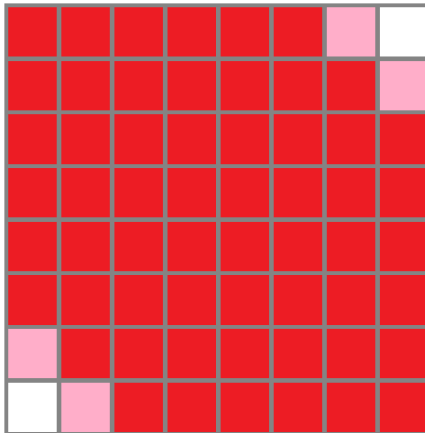
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



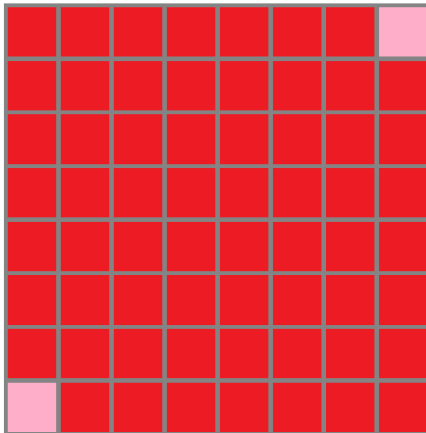
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



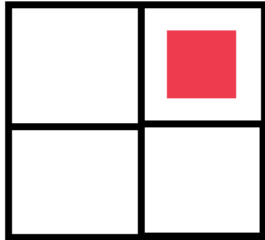
# Zombie Infection: Conquering The World

Fewest number of initial infections needed to get all...?



# Zombie Infection: Can $n - 1$ infect all on an $n \times n$ board?

# Zombie Infection: Can $n - 1$ infect all on an $n \times n$ board?

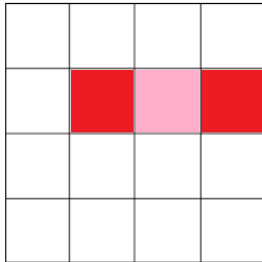
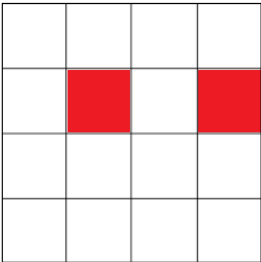


■	1	2
1	3	4
2	4	5

1	■	1
2	3	2
4	5	4

2	1	2
1	■	1
2	1	2

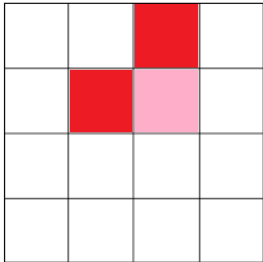
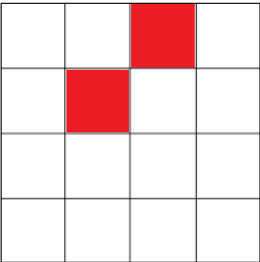
# Zombie Infection: Can $n - 1$ infect all on an $n \times n$ board?



Perimeter of infection unchanged.

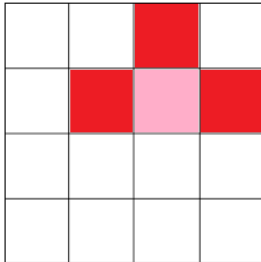
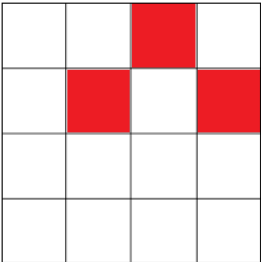


# Zombie Infection: Can $n - 1$ infect all on an $n \times n$ board?



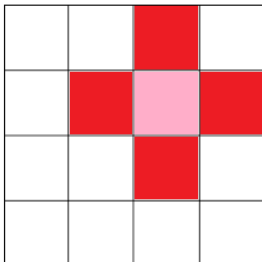
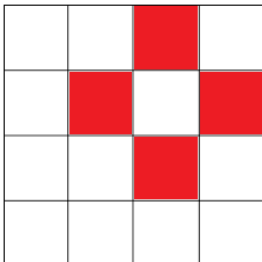
Perimeter of infection unchanged.

# Zombie Infection: Can $n - 1$ infect all on an $n \times n$ board?



Perimeter of infection decreases by 2.

# Zombie Infection: Can $n - 1$ infect all on an $n \times n$ board?



Perimeter of infection decreases by 4.

# Zombie Infection: $n - 1$ cannot infect all

- If  $n - 1$  infected, maximum perimeter is  $4(n - 1) = 4n - 4$ .

## Zombie Infection: $n - 1$ cannot infect all

- If  $n - 1$  infected, maximum perimeter is  $4(n - 1) = 4n - 4$ .
- **Mono-variant:** As time passes, perimeter of infection never increases.

## Zombie Infection: $n - 1$ cannot infect all

- If  $n - 1$  infected, maximum perimeter is  $4(n - 1) = 4n - 4$ .
- **Mono-variant:** As time passes, perimeter of infection never increases.
- Perimeter of  $n \times n$  square is  $4n$ , so at least 1 square safe!

## Zombie Infection: $n - 1$ cannot infect all

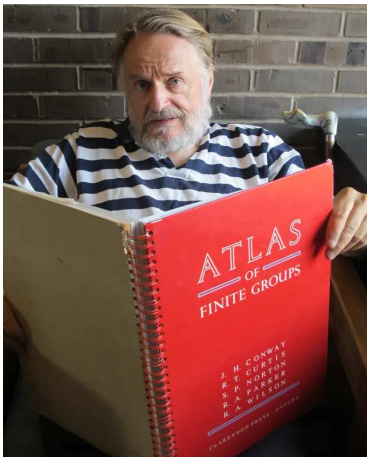
- If  $n - 1$  infected, maximum perimeter is  $4(n - 1) = 4n - 4$ .
- **Mono-variant:** As time passes, perimeter of infection never increases.
- Perimeter of  $n \times n$  square is  $4n$ , so at least 1 square safe!
- How many must be safe?
- Other questions?

## Zombie Infection: $n - 1$ cannot infect all

- If  $n - 1$  infected, maximum perimeter is  $4(n - 1) = 4n - 4$ .
- **Mono-variant:** As time passes, perimeter of infection never increases.
- Perimeter of  $n \times n$  square is  $4n$ , so at least 1 square safe!
- How many must be safe?
- Other questions: Is a row safe? Higher dimensions? Other regions (torus)?



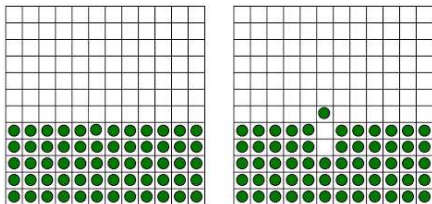
## Conway's Soldiers



**Figure:** John Horton Conway: Image from The Guardian.

## Conway's Soldiers / Checker Problem

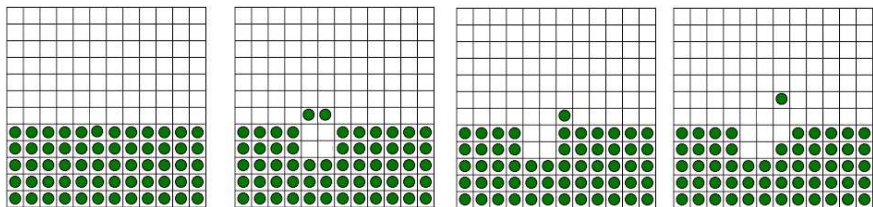
Problem: Infinite checkerboard, pieces at all  $(x, y)$  with  $y \leq 0$ .  
Using horizontal / vertical jumps (jumped piece gone forever),  
how high can you move a piece?



**Figure:** Left: A subset of the initial configuration. Right: moving a soldier / checker up 1.

## Conway's Soldiers

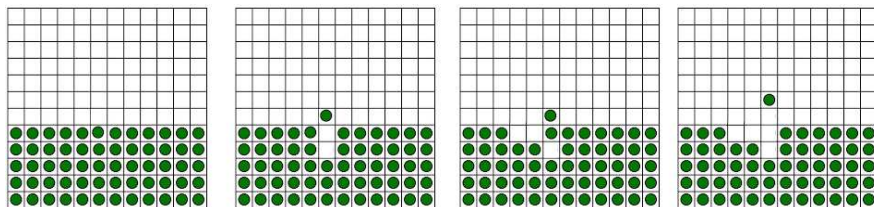
Problem: Infinite checkerboard, pieces at all  $(x, y)$  with  $y \leq 0$ . Using horizontal / vertical jumps (jumped piece gone forever), how high can you move a piece?



**Figure:** Left: A subset of the initial configuration. Right: moving a soldier / checker up 2. Can you do 3? 4? 5? Any height?

## Conway's Soldiers

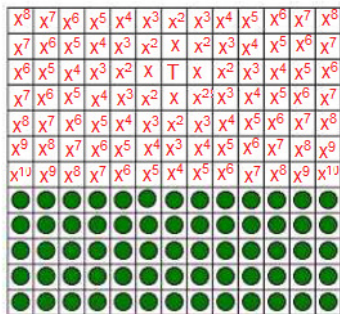
Problem: Infinite checkerboard, pieces at all  $(x, y)$  with  $y \leq 0$ .  
Using horizontal / vertical jumps (jumped piece gone forever),  
how high can you move a piece?



**Figure:** Left: A subset of the initial configuration. Right: Also moving a soldier / checker up 2. Can you do 3? 4? 5? Any height?

## Conway's Soldiers: The Monovariant: I

Problem: Infinite checkerboard, pieces at all  $(x, y)$  with  $y \leq 0$ .  
Using horizontal / vertical jumps (jumped piece gone forever),  
how high can you move a piece?



**Figure:** Conway's monovariant: What is it?

## Conway's Soldiers: The Monovariant: II

Choose target  $T = (0, 5)$ .

Fix  $x$  (to be determined later) and attach  $x^{i+j}$  to a point that is  $i$  units horizontally and  $j$  units vertically from  $T$ .

$x^8$	$x^7$	$x^6$	$x^5$	$x^4$	$x^3$	$x^2$	$x^3$	$x^4$	$x^5$	$x^6$	$x^7$	$x^8$
$x^7$	$x^6$	$x^5$	$x^4$	$x^3$	$x^2$	$x$	$x^2$	$x^3$	$x^4$	$x^5$	$x^6$	$x^7$
$x^6$	$x^5$	$x^4$	$x^3$	$x^2$	$x$	$T$	$x$	$x^2$	$x^3$	$x^4$	$x^5$	$x^6$
$x^7$	$x^6$	$x^5$	$x^4$	$x^3$	$x^2$	$x$	$x^2$	$x^3$	$x^4$	$x^5$	$x^6$	$x^7$
$x^8$	$x^7$	$x^6$	$x^5$	$x^4$	$x^3$	$x^2$	$x^3$	$x^4$	$x^5$	$x^6$	$x^7$	$x^8$
$x^9$	$x^8$	$x^7$	$x^6$	$x^5$	$x^4$	$x^3$	$x^4$	$x^5$	$x^6$	$x^7$	$x^8$	$x^9$
$x^{1+j}$	$x^9$	$x^8$	$x^7$	$x^6$	$x^5$	$x^4$	$x^5$	$x^6$	$x^7$	$x^8$	$x^9$	$x^{1+j}$
●	●	●	●	●	●	●	●	●	●	●	●	●
●	●	●	●	●	●	●	●	●	●	●	●	●
●	●	●	●	●	●	●	●	●	●	●	●	●
●	●	●	●	●	●	●	●	●	●	●	●	●
●	●	●	●	●	●	●	●	●	●	●	●	●

## Conway's Soldiers: The Monovariant: III

Choose a target point  $T$ ; for us it is a point of height 5 above the checkers:  $T = (0, 5)$ .

Fix  $x$  (to be determined later) and attach  $x^{i+j}$  to a point that is  $i$  units horizontally and  $j$  units vertically from  $T$ .

**What is the value of the initial board?**

- Zeroth row:  $\dots, x^7, x^6, x^5, x^6, x^7, \dots$ : sum is

$$x^5 + 2 \sum_{k=6}^{\infty} x^k = x^5 + \frac{2x^6}{1-x} = \frac{(1+x)x^5}{1-x}.$$

## Conway's Soldiers: The Monovariant: III

Choose a target point  $T$ ; for us it is a point of height 5 above the checkers:  $T = (0, 5)$ .

Fix  $x$  (to be determined later) and attach  $x^{i+j}$  to a point that is  $i$  units horizontally and  $j$  units vertically from  $T$ .

**What is the value of the initial board?**

- Zeroth row:  $\dots, x^7, x^6, x^5, x^6, x^7, \dots$ : sum is

$$x^5 + 2 \sum_{k=6}^{\infty} x^k = x^5 + \frac{2x^6}{1-x} = \frac{(1+x)x^5}{1-x}.$$

- Each row is  $x$  times previous: Thus initial board value is

$$\frac{(1+x)x^5}{1-x} \sum_{n=0}^{\infty} x^n = \frac{(1+x)x^5}{(1-x)^2}.$$



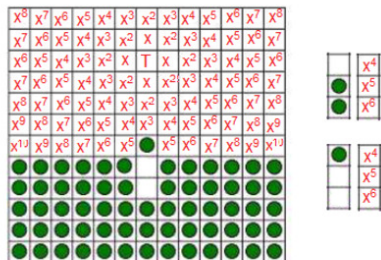
# Conway's Soldiers: The Monovariant: IV

Two moves: lose 2 pieces and add a piece further from  $T$ , or lose 2 pieces and add a piece closer to  $T$ .

First type of move clearly decreases value of board.

## Conway's Soldiers: The Monovariant: IV

Two moves: lose 2 pieces and add a piece further from  $T$ , or lose 2 pieces and add a piece closer to  $T$ .



**Figure:** Moving pieces on  $x^6$  and  $x^5$  to on  $x^4$ .

Change is  $x^4 - x^5 - x^6 = x^4(1 - x - x^2)$ , want this to be zero.

## Conway's Soldiers: The Monovariant: IV

Two moves: lose 2 pieces and add a piece further from  $T$ , or lose 2 pieces and add a piece closer to  $T$ .

Second type replaces  $x^{n+2}$  and  $x^{n+1}$  with an  $x^n$ : change is  $x^n - x^{n+1} - x^{n+2}$ . Choose  $x$  so that this change is zero.

Thus  $1 - x - x^2 = 0$  or  $x = (-1 \pm \sqrt{5})/2$ . Take positive root,  $(-1 + \sqrt{5})/2 = \varphi - 1$  ( $\varphi$  the golden mean).

**Monovariant: sum of the values of squares with checkers.**

## Conway's Soldiers: The Monovariant: V

Choose a target point  $T$ .

- Initial board value is

$$\frac{(1+x)x^5}{(1-x)^2} : \text{ when } x = \frac{\sqrt{5}-1}{2} \text{ get } 1.$$

- Target at  $(0, 4)$  contributes  $x = \frac{\sqrt{5}-1}{2} \approx 0.618034$ ; as less than 1 possible (and can be done).
- Target at  $(0, 5)$ , board's value at least 1. Moves never increase value: **IMPOSSIBLE IN FINITE TIME!**<sup>1</sup>

---

<sup>1</sup>Possible in "infinite" game: <https://tartarus.org/gareth/maths/stuff/solarmy.pdf>.

## New Results

**Conway Checkers  $m$ -game:** Start with  $m$  checkers on each gridpoint (original game is just 1), if jump over it lose one checker.

### SMALL 2024

Given a Conway Checkers  $m$ -game, the maximum row attainable,  $n_m$ , satisfies

$$\lfloor \log_{\varphi}(m) + 4.67 \rfloor \leq n_m \leq \lfloor \log_{\varphi}(m) + 5 \rfloor$$

for sufficiently large  $m$ , where  $\varphi$  is the golden ratio  $\frac{\sqrt{5}+1}{2}$ .

# Zeckendorf Minimality

## Introduction: Summand Minimality

Fibonacci:  $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, F_{n+2} = F_{n+1} + F_n$ .

### Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

Example:

$$2024 = 1597 + 377 + 34 + 13 + 3 = F_{16} + F_{13} + F_8 + F_6 + F_3.$$

## Introduction: Summand Minimality

Fibonacci:  $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, F_{n+2} = F_{n+1} + F_n$ .

### Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

**Example:**

$$2024 = 1597 + 377 + 34 + 13 + 3 = F_{16} + F_{13} + F_8 + F_6 + F_3.$$

Conversely, we can construct the Fibonacci sequence using this property:

1



## Introduction: Summand Minimality

Fibonacci:  $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, F_{n+2} = F_{n+1} + F_n$ .

### Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

**Example:**

$$2024 = 1597 + 377 + 34 + 13 + 3 = F_{16} + F_{13} + F_8 + F_6 + F_3.$$

Conversely, we can construct the Fibonacci sequence using this property:

1, 2

## Introduction: Summand Minimality

Fibonacci:  $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, F_{n+2} = F_{n+1} + F_n$ .

### Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

**Example:**

$$2024 = 1597 + 377 + 34 + 13 + 3 = F_{16} + F_{13} + F_8 + F_6 + F_3.$$

Conversely, we can construct the Fibonacci sequence using this property:

1, 2, 3

## Introduction: Summand Minimality

Fibonacci:  $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, F_{n+2} = F_{n+1} + F_n$ .

### Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

**Example:**

$$2024 = 1597 + 377 + 34 + 13 + 3 = F_{16} + F_{13} + F_8 + F_6 + F_3.$$

Conversely, we can construct the Fibonacci sequence using this property:

1, 2, 3, 5

## Introduction: Summand Minimality

Fibonacci:  $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, F_{n+2} = F_{n+1} + F_n$ .

### Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

**Example:**

$$2024 = 1597 + 377 + 34 + 13 + 3 = F_{16} + F_{13} + F_8 + F_6 + F_3.$$

Conversely, we can construct the Fibonacci sequence using this property:

1, 2, 3, 5, 8

## Introduction: Summand Minimality

Fibonacci:  $F_1 = 1, F_2 = 2, F_3 = 3, F_4 = 5, F_{n+2} = F_{n+1} + F_n$ .

### Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

**Example:**

$$2024 = 1597 + 377 + 34 + 13 + 3 = F_{16} + F_{13} + F_8 + F_6 + F_3.$$

Conversely, we can construct the Fibonacci sequence using this property:

1, 2, 3, 5, 8, 13...

## Summand Minimality

### Example

- $18 = 13 + 5 = F_6 + F_4$ , legal decomposition, two summands.
- $18 = 13 + 3 + 2 = F_6 + F_3 + F_2$ , non-legal decomposition, three summands.

### Theorem

*The Zeckendorf decomposition is **summand minimal**.*

### Overall Question

What other recurrences are summand minimal?

# Zeckendorf Decomposition is Minimal

## Theorem

*The Zeckendorf decomposition is **summand minimal**: no decomposition as a sum of Fibonacci numbers (1, 2, 3, 5, ...) has fewer summands than it.*

## Zeckendorf Decomposition is Minimal

### Theorem

*The Zeckendorf decomposition is **summand minimal**: no decomposition as a sum of Fibonacci numbers (1, 2, 3, 5, ...) has fewer summands than it.*

If  $n = \sum_k a_k F_k$  (with  $a_k$  non-negative integers), define the weighted index attached to this decomposition  $\mathcal{D}$  to be

$$\text{Index}(\mathcal{D}) = \sum_k a_k \sqrt{k}.$$

More natural  $\sum_k a_k k$  but square-root makes strictly decreasing.



## Zeckendorf Decomposition is Minimal

### Theorem

*The Zeckendorf decomposition is **summand minimal**: no decomposition as a sum of Fibonacci numbers (1, 2, 3, 5, ...) has fewer summands than it.*

If  $n = \sum_k a_k F_k$  (with  $a_k$  non-negative integers), define the weighted index attached to this decomposition  $\mathcal{D}$  to be

$$\text{Index}(\mathcal{D}) = \sum_k a_k \sqrt{k}.$$

More natural  $\sum_k a_k k$  but square-root makes strictly decreasing.

Bounded process: For fixed  $n$ , only indices up to certain point used, and  $a_k \leq n$ .

## Zeckendorf Decomposition is Minimal: Proof

Show  $\text{Index}(\mathcal{D})$  is a mono-variant, end in the Zeckendorf decomposition, number summands never increased.

## Zeckendorf Decomposition is Minimal: Proof

Show  $\text{Index}(\mathcal{D})$  is a mono-variant, end in the Zeckendorf decomposition, number summands never increased.

If  $\mathcal{D}$  is not the Zeckendorf, have  $2F_k$  or  $F_k \wedge F_{k+1}$ .

## Zeckendorf Decomposition is Minimal: Proof

Show  $\text{Index}(\mathcal{D})$  is a mono-variant, end in the Zeckendorf decomposition, number summands never increased.

If  $\mathcal{D}$  is not the Zeckendorf, have  $2F_k$  or  $F_k \wedge F_{k+1}$ .

$$F_k \wedge F_{k+1} \rightarrow F_{k+2}:$$

- $\sqrt{k} + \sqrt{k+1} > \sqrt{k+2}$ .

$$2F_k \rightarrow F_{k-2} + F_{k+1}:$$

- $k \geq 3: 2\sqrt{k} > \sqrt{k-2} + \sqrt{k+1}$

- $k = 2: 2\sqrt{2} > \sqrt{1} + \sqrt{3}$

- $k = 1: 2\sqrt{1} > \sqrt{2}$

Only finitely many values, each move lowers, continue till hit Zeckendorf, number of summands never increased.

## Positive Linear Recurrence Sequences

### Definition

A **positive linear recurrence sequence (PLRS)** is a sequence given by a recurrence  $\{a_n\}$  with

$$a_n := c_1 a_{n-1} + \cdots + c_t a_{n-t}$$

and each  $c_i \geq 0$  and  $c_1, c_t > 0$ . We use **ideal initial conditions**  $a_{-(n-1)} = 0, \dots, a_{-1} = 0, a_0 = 1$  and call  $(c_1, \dots, c_t)$  the **signature of the sequence**.

### Theorem (Cordwell, Hlavacek, Huynh, M., Peterson, Vu)

*For a PLRS with signature  $(c_1, c_2, \dots, c_t)$ , the Generalized Zeckendorf Decompositions are summand minimal if and only if*

$$c_1 \geq c_2 \geq \cdots \geq c_t.$$

## Zeckendorf Games

# Fibonacci Game: Rules

- Two player game, alternate turns, last to move wins.

## Fibonacci Game: Rules

- Two player game, alternate turns, last to move wins.
- Bins  $F_1, F_2, F_3, \dots$ , start with  $N$  pieces in  $F_1$  and others empty.



## Fibonacci Game: Rules

- Two player game, alternate turns, last to move wins.
- Bins  $F_1, F_2, F_3, \dots$ , start with  $N$  pieces in  $F_1$  and others empty.
- A turn is one of the following moves:
  - ◇ If have two pieces on  $F_k$  can remove and put one piece at  $F_{k+1}$  and one at  $F_{k-2}$   
(if  $k = 1$  then  $2F_1$  becomes  $1F_2$ )
  - ◇ If pieces at  $F_k$  and  $F_{k+1}$  remove and add one at  $F_{k+2}$ .

## Fibonacci Game: Rules

- Two player game, alternate turns, last to move wins.
- Bins  $F_1, F_2, F_3, \dots$ , start with  $N$  pieces in  $F_1$  and others empty.
- A turn is one of the following moves:
  - ◇ If have two pieces on  $F_k$  can remove and put one piece at  $F_{k+1}$  and one at  $F_{k-2}$   
(if  $k = 1$  then  $2F_1$  becomes  $1F_2$ )
  - ◇ If pieces at  $F_k$  and  $F_{k+1}$  remove and add one at  $F_{k+2}$ .

### Questions:

- Does the game end? How long?
- For each  $N$  who has the winning strategy?
- What is the winning strategy?

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

10	0	0	0	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 1:  $F_1 + F_1 = F_2$

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

8	1	0	0	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 2:  $F_1 + F_1 = F_2$

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

6	2	0	0	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 1:  $2F_2 = F_3 + F_1$

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

7	0	1	0	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 2:  $F_1 + F_1 = F_2$

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

5	1	1	0	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 1:  $F_2 + F_3 = F_4$ .

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

5	0	0	1	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 2:  $F_1 + F_1 = F_2$ .



## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

3	1	0	1	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 1:  $F_1 + F_1 = F_2$ .

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

1	2	0	1	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 2:  $F_1 + F_2 = F_3$ .

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

0	1	1	1	0
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

Next move: Player 1:  $F_3 + F_4 = F_5$ .

## Sample Game

Start with 10 pieces at  $F_1$ , rest empty.

---

0	1	0	0	1
$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

---

No moves left, Player One wins.

## Sample Game

Player One won in 9 moves.

(1)	10	0	0	0	0
(2)	8	1	0	0	0
(1)	6	2	0	0	0
(2)	7	0	1	0	0
(1)	5	1	1	0	0
(2)	5	0	0	1	0
(1)	3	1	0	1	0
(2)	1	2	0	1	0
(1)	0	1	1	1	0
	0	1	0	0	1
	$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

## Sample Game

Player Two won in 10 moves.

(1)	10	0	0	0	0
(2)	8	1	0	0	0
(1)	6	2	0	0	0
(2)	7	0	1	0	0
(1)	5	1	1	0	0
(2)	5	0	0	1	0
(1)	3	1	0	1	0
(2)	1	2	0	1	0
(1)	2	0	1	1	0
(2)	0	1	1	1	0
	0	1	0	0	1
	$[F_1 = 1]$	$[F_2 = 2]$	$[F_3 = 3]$	$[F_4 = 5]$	$[F_5 = 8]$

## Games end

### Theorem

*All games end in finitely many moves.*

**Proof:** The sum of the square roots of the indices is a strict monovariant.

- Adding consecutive:  $(\sqrt{k} + \sqrt{k+1}) - \sqrt{k+2} > 0$ .
- Splitting:  $2\sqrt{k} - (\sqrt{k+1} + \sqrt{k+1}) > 0$ .
- Spitting 1's:  $2\sqrt{1} - \sqrt{2} > 0$ .
- Splitting 2's:  $2\sqrt{2} - (\sqrt{3} + \sqrt{1}) > 0$ .

## Games Lengths: I

**Upper bound:** At most  $3n - 3Z(n) - I(n) + 1$  moves

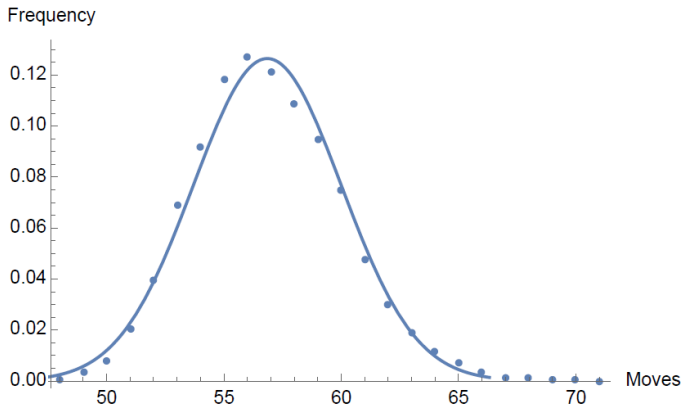
- $Z(n)$  is the number of terms in the Zeckendorf decomposition,
- $I(n)$  is the sum of the indices.

**Fastest game:**  $n - Z(n)$  moves ( $Z(n)$  is the number of summands in  $n$ 's Zeckendorf decomposition).

From always moving on the largest summand possible (deterministic).



## Games Lengths: II



**Figure:** Frequency graph of the number of moves in 9,999 simulations of the Zeckendorf Game with random moves when  $n = 60$  vs a Gaussian. **Natural conjecture....**

## Winning Strategy

### Theorem

*Player Two Has a Winning Strategy*

Idea is to show if not, Player Two could steal Player One's strategy.

**Non-constructive!**

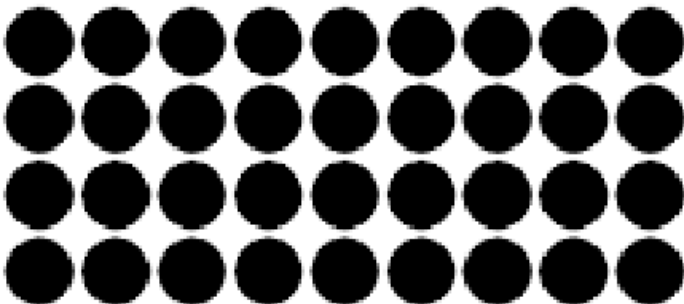
Will highlight idea with a simpler game.

## Winning Strategy: Intuition from Dot Game

Two players, alternate. Turn is choosing a dot at  $(i, j)$  and coloring every dot  $(m, n)$  with  $i \leq m$  and  $j \leq n$ .

Once all dots colored game ends; whomever goes last loses.

Prove Player 1 has a winning strategy!

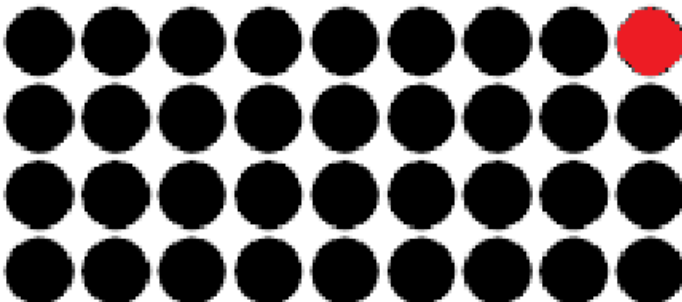


## Winning Strategy: Intuition from Dot Game

Two players, alternate. Turn is choosing a dot at  $(i, j)$  and coloring every dot  $(m, n)$  with  $i \leq m$  and  $j \leq n$ .

Once all dots colored game ends; whomever goes last loses.

**Proof Player 1 has a winning strategy.** If have, play; if not, steal.

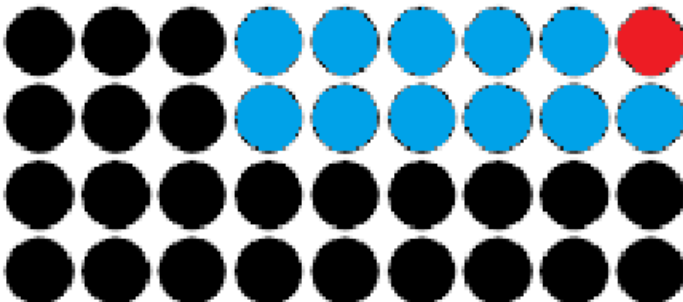


## Winning Strategy: Intuition from Dot Game

Two players, alternate. Turn is choosing a dot at  $(i, j)$  and coloring every dot  $(m, n)$  with  $i \leq m$  and  $j \leq n$ .

Once all dots colored game ends; whomever goes last loses.

**Proof Player 1 has a winning strategy.** If have, play; if not, steal.

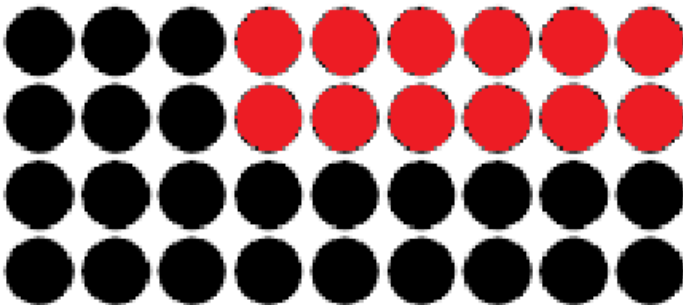


## Winning Strategy: Intuition from Dot Game

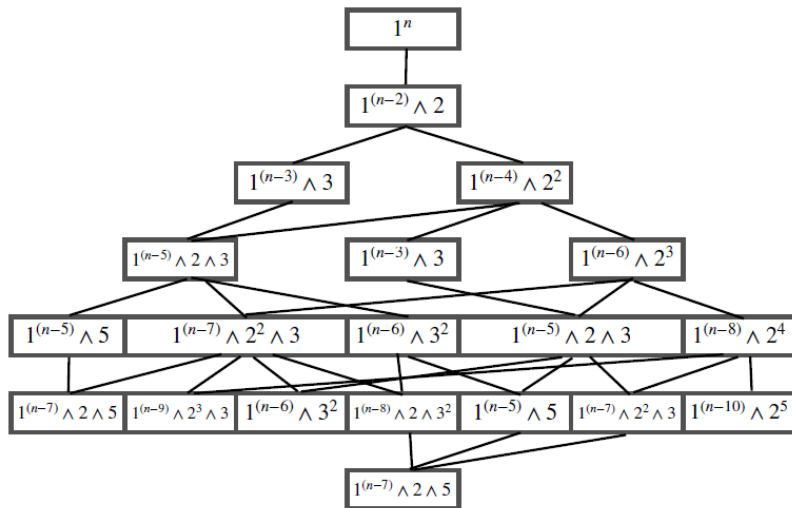
Two players, alternate. Turn is choosing a dot at  $(i, j)$  and coloring every dot  $(m, n)$  with  $i \leq m$  and  $j \leq n$ .

Once all dots colored game ends; whomever goes last loses.

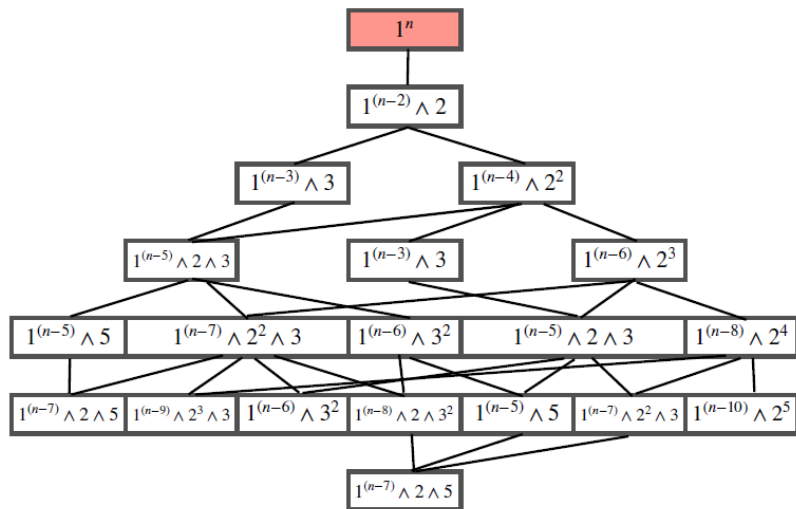
**Proof Player 1 has a winning strategy.** If have, play; if not, steal.



## Sketch of Proof for Player Two's Winning Strategy

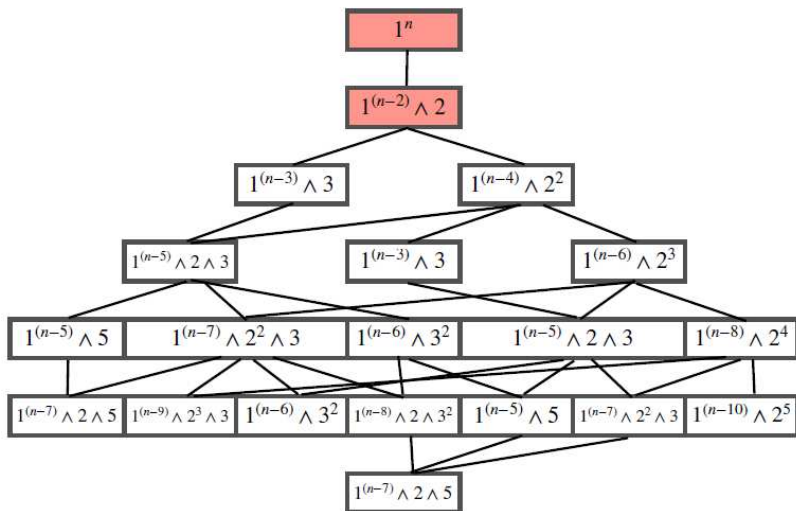


## Sketch of Proof for Player Two's Winning Strategy

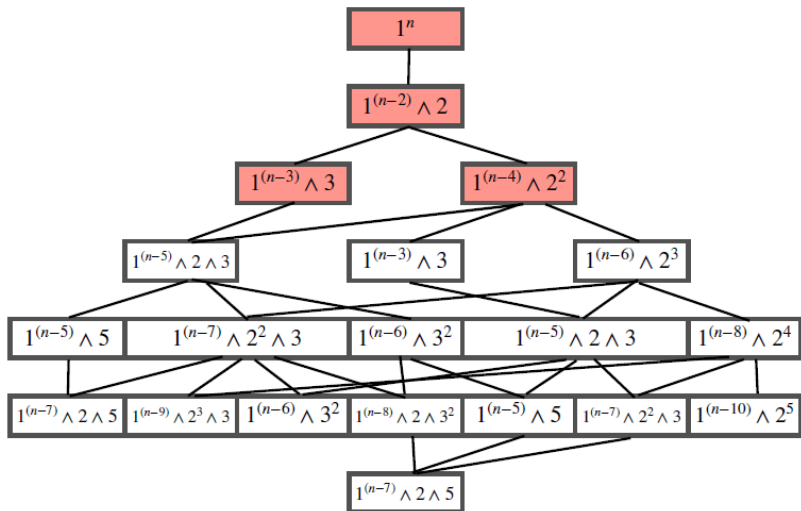




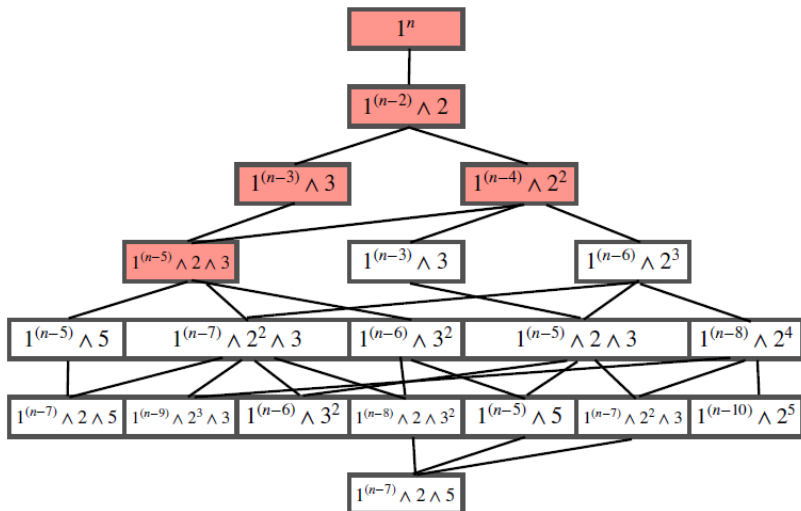
## Sketch of Proof for Player Two's Winning Strategy



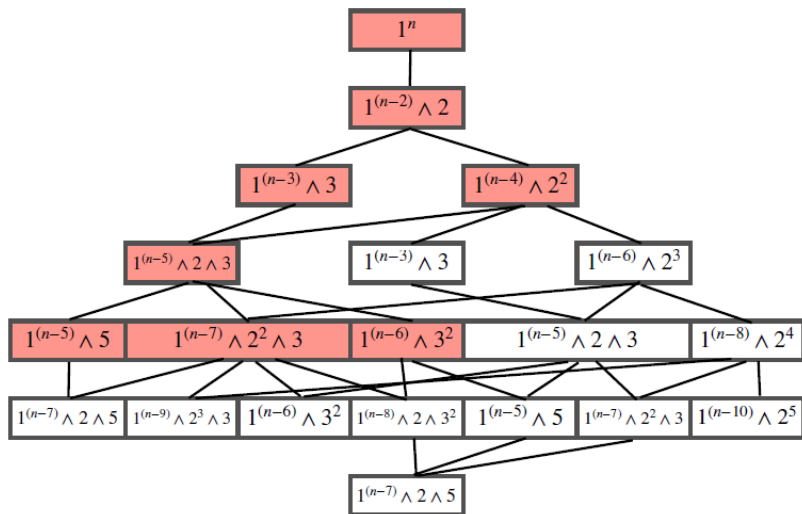
## Sketch of Proof for Player Two's Winning Strategy



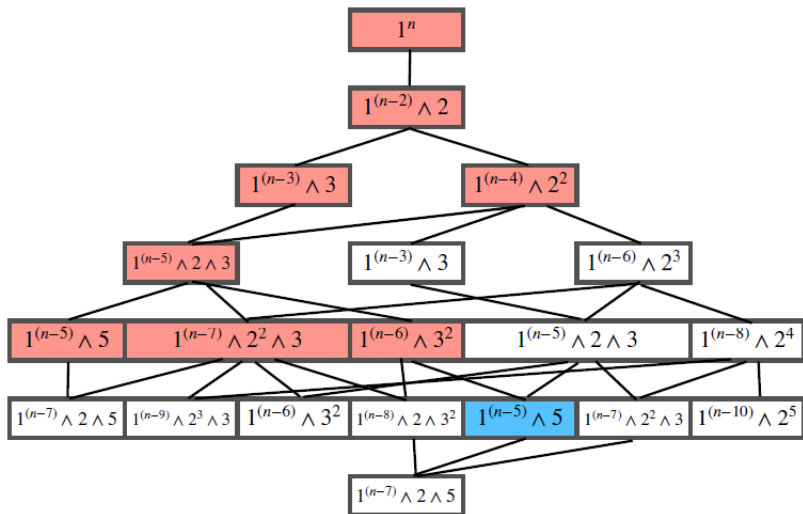
## Sketch of Proof for Player Two's Winning Strategy



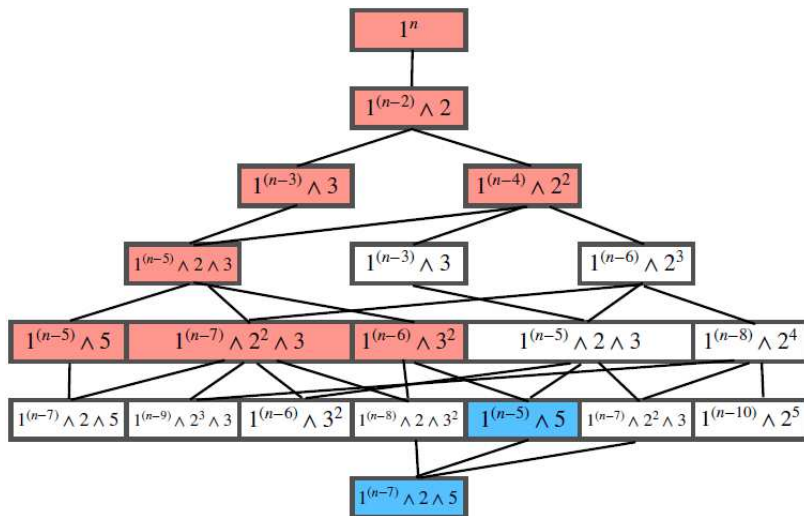
## Sketch of Proof for Player Two's Winning Strategy



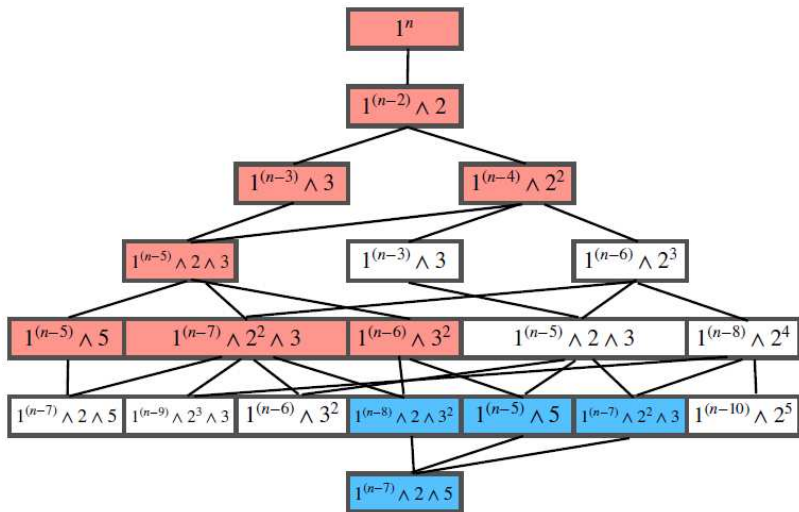
## Sketch of Proof for Player Two's Winning Strategy



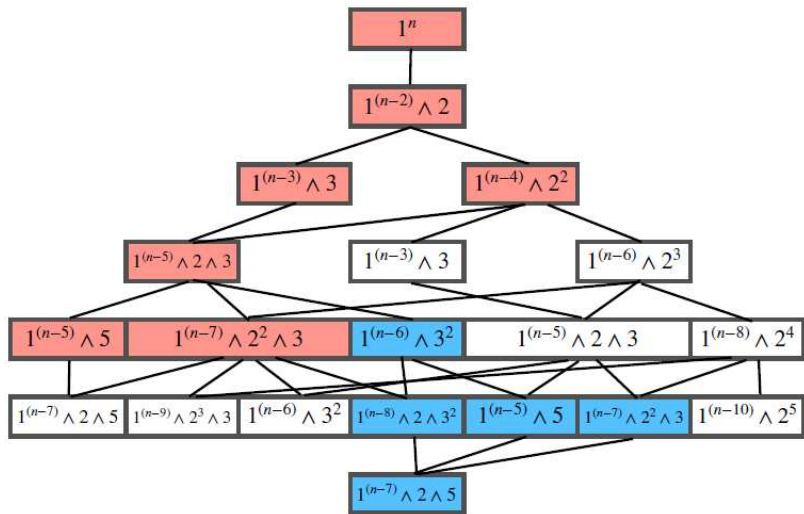
## Sketch of Proof for Player Two's Winning Strategy



## Sketch of Proof for Player Two's Winning Strategy



## Sketch of Proof for Player Two's Winning Strategy





## The Bergman Game

### Definition

The **Bergman Game** is played with the standard split/combine moves from the Zeckendorf game, but on a two-sided infinite tape instead of a one-sided infinite tape.

It produces base- $\varphi$  decompositions ( $\varphi = (1 + \sqrt{5})/2$ ).

### Example

0	0	4	0	0
1	0	2	1	0
1	0	1	0	1

$$4 = \varphi^{-2} + \varphi^0 + \varphi^2.$$

## The Bergman Game

**Theorem (Baily, Dell, Durmić, Fleischmann, Jackson, Mijares, M., Pesikoff, Reifenberg, Reina, Yang)**

*The longest Bergman Game with  $n$  summands terminates in  $\Theta(n^2)$  time regardless of where the summands are placed. The shortest possible Bergman Game terminates in  $\Theta(n)$  time.*

**Natural Question: Who has the winning strategy?**

- Not currently known.
- Game tree explodes, escaping a strategy steal.

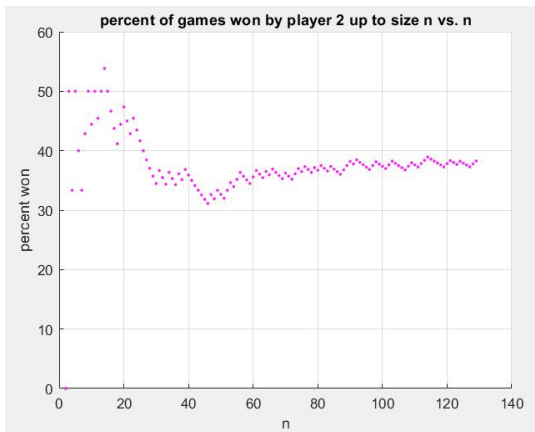
## The Frodnekcez Game (Reverse Zeckendorf Game): Rules

- The Zeckendorf game in reverse, last to move wins.
- Bins  $F_1, F_2, F_3, \dots$ , for some natural number  $N$ , start with one piece in bin  $F_k$  if  $F_k$  is in the Zeckendorf decomposition of  $N$ , and have other bins empty.
- A turn is one of the following moves:
  - ◇ If one piece at  $F_{k+1}$  and one at  $F_{k-2}$ , can remove and add two pieces on  $F_k$ .
  - ◇ If piece at  $F_{k+2}$ , remove and add one piece at both  $F_k$  and  $F_{k+1}$ .

( $F_1$  and  $F_3$  becomes  $2F_2$ , and  $F_2$  becomes  $2F_1$ )

Problem created and analyzed by PANTHERs 2023 from the 2023 SMALL REU: Zoe Batterman, Aditya Jambhale, Akash Narayanan, Kishan Sharma, Andrew Yang, Chris Yao.

## Winning Strategy?



**Figure:** In the forward Zeckendorf game, Player 2 wins for all  $N > 2$ . The reverse game is more interesting. **Natural conjecture...**

## Current / Future Work

- What if  $p \geq 3$  people play the Fibonacci game? Some multi-player results.
- Does the number of moves in random games converge to a Gaussian? Evidence....
- How long do games take? Proved closed interval.
- Accelerated games: do as many of one move as wish....
- What of other recurrences?

**\$500 Prize: Determine the winning strategy.**

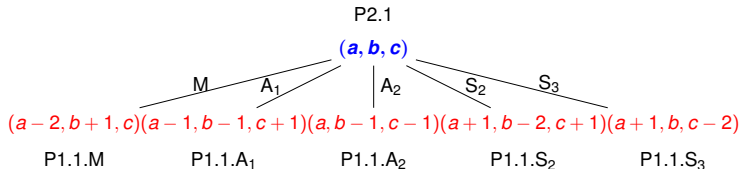
# Black Hole Zeckendorf Game (Ongoing Work: SMALL 2024)

How can we simplify the game?

## $F_m$ Black Hole Variation

Any pieces placed in a column  $F_i$  for  $i \geq m$  are permanently removed from gameplay.

For the  $F_4$  case, this allows for the following moves:



Thanks / References

## Thanks

Work supported by multiple NSF REU grants, the Eureka Program, the Finnerty Fund, **Amherst**, **Carnegie Mellon**, **Michigan**, **Williams** and **Yale**.

Many thanks to all my co-authors, including Benjamin Baily, Paul Baird-Smith, Ela Boldyriew, Katherine Cordwell, Anna Cusenza, Linglong Dai, Justine Dell, Pei Ding, Aidan Dunkelberg, Irfan Durmic, Alyssa Epstein, Henry Fleischmann, Kristen Flint, Diego Garcia-Fernandezsesma, John Haviland, Max Hlavacek, Kate Huffman, Chi Huynh, Faye Jackson, Dianhui Ke, Daniel Kleber, Jason Kuretski, Phuc Lam, John Lentfer, Ruoci Li, Xiaonan Li, Tianhao Luo, Micah McClatchey, Isaac Mijares, Clayton Mizgerd, Alexandra Newlon, Ethan Pesikoff, Carsten Peterson, Thomas Rascon, Luke Reifenberg, Alicia Smith Reina, Eliel Sosis, Chenyang Sun, Vashisth Tiwari, Fernando Trejos Suarez, Risa Vandegrift, Yen Nhi Truong Vu, Dong Xia, Ajmain Yamin, Yingzi Yang, Jingkai Ye, Yunhao Zhang, Xiaoyan Zheng, Zhyi Zhou, Weiduo Zhu.



# Papers

- The Zeckendorf Game (with Paul Baird-Smith, Alyssa Epstein and Kristen Flint), Proceedings of CANT 2018. [https://web.williams.edu/Mathematics/sjmiller/public\\_html/math/papers/ZeckGameCANT10.pdf](https://web.williams.edu/Mathematics/sjmiller/public_html/math/papers/ZeckGameCANT10.pdf)
- On Summand Minimality of Generalized Zeckendorf Decompositions (with Katherine Cordwell, Max Hlavacek, Chi Huynh, Carsten Peterson, and Yen Nhi Truong Vu), Research in Number Theory (4 (2018), no. 43, <https://doi.org/10.1007/s40993-018-0137-7>)/ [https://web.williams.edu/Mathematics/sjmiller/public\\_html/math/papers/ZeckMinimalSummands61.pdf](https://web.williams.edu/Mathematics/sjmiller/public_html/math/papers/ZeckMinimalSummands61.pdf)
- The Generalized Zeckendorf Game (with Paul Baird-Smith, Alyssa Epstein and Kristen Flint), Fibonacci Quarterly (57 (2019) no. 5, 1-14). [https://web.williams.edu/Mathematics/sjmiller/public\\_html/math/papers/ZeckGameGeneral\\_FibQ10.pdf](https://web.williams.edu/Mathematics/sjmiller/public_html/math/papers/ZeckGameGeneral_FibQ10.pdf)
- The Fibonacci Quilt Game (with Alexandra Newlon), Fibonacci Quarterly (2 (2020), 157-168). [https://web.williams.edu/Mathematics/sjmiller/public\\_html/math/papers/FQgame30.pdf](https://web.williams.edu/Mathematics/sjmiller/public_html/math/papers/FQgame30.pdf)
- Extending Zeckendorf's Theorem to a Non-constant Recurrence and the Zeckendorf Game on this Non-constant Recurrence Relation (with Ela Boldyriew, Anna Cusenza, Linglong Dai, Pei Ding, Aidan Dunkelberg, John Haviland, Kate Huffman, Dianhui Ke, Daniel Kleber, Jason Kuretski, John Lentfer, Tianhao Luo, Clayton Mizgerd, Vashisth Tiwari, Jingkai Ye, Yunhao Zhang, Xiaoyan Zheng, and Weiduo Zhu), Fibonacci Quarterly. (5 (2020), 55-76). [https://web.williams.edu/Mathematics/sjmiller/public\\_html/math/papers/ZeckExtendingZeckNonConstantGame40.pdf](https://web.williams.edu/Mathematics/sjmiller/public_html/math/papers/ZeckExtendingZeckNonConstantGame40.pdf)
- Deterministic Zeckendorf Games (with Ruoci Li, Xiaonan Li, Clay Mizgerd, Chenyang Sun, Dong Xia, and Zhiyi Zhou), Fibonacci Quarterly. (58 (2020), no. 5, 152-160). [https://web.williams.edu/Mathematics/sjmiller/public\\_html/math/papers/zeckgamedeterministic31.pdf](https://web.williams.edu/Mathematics/sjmiller/public_html/math/papers/zeckgamedeterministic31.pdf)

## Papers

- Winning Strategy for the Multiplayer and Multialliance Zeckendorf Games (with Anna Cusenza, Aiden Dunkelberg, Kate Huffman, Dianhui Ke, Daniel Kleber, Clayton Mizgerd, Vashisth Tiwari, Jingkai Ye, Xiaoyan Zheng), *Fibonacci Quarterly*. (59 (2021), 308–318). [https://web.williams.edu/Mathematics/sjmillier/public\\_html/math/papers/ZeckGameWinningAlliancePolymath20.pdf](https://web.williams.edu/Mathematics/sjmillier/public_html/math/papers/ZeckGameWinningAlliancePolymath20.pdf)
- Bounds on Zeckendorf Games (with Anna Cusenza, Aiden Dunkelberg, Kate Huffman, Dianhui Ke, Daniel Kleber, Micah McClatchey, Clayton Mizgerd, Vashisth Tiwari, Jingkai Ye, Xiaoyan Zheng), *Fibonacci Quarterly* (1 (2022), no. 1, 57–71). [https://web.williams.edu/Mathematics/sjmillier/public\\_html/math/papers/ZeckGame\\_BoundLength\\_2020polyreu\\_10.pdf](https://web.williams.edu/Mathematics/sjmillier/public_html/math/papers/ZeckGame_BoundLength_2020polyreu_10.pdf)
- Completeness of Positive Linear Recurrence Sequences (with Ela Boldyriew, John Haviland, Phuc Lam, John Lentfer, Fernando Trejos Suarez), submitted to CANT Conference Proceedings. [https://web.williams.edu/Mathematics/sjmillier/public\\_html/math/papers/ZeckCompletePLRS02.pdf](https://web.williams.edu/Mathematics/sjmillier/public_html/math/papers/ZeckCompletePLRS02.pdf)
- The Generalized Bergman Game (with Benjamin Baily, Justine Dell, Irfan Durmic, Henry Fleischmann, Faye Jackson, Isaac Mijares, Ethan Pesikoff, Luke Reifenberg, Alicia Smith Reina, Yingzi Yang). <https://arxiv.org/abs/2109.00117>
- Winning Strategies for the Generalized Zeckendorf Game (Steven J. Miller, Eliel Sosis, Jingkai Ye), submitted to the *Fibonacci Quarterly*. [https://web.williams.edu/Mathematics/sjmillier/public\\_html/math/papers/ZeckGame\\_WinStrategies\\_Polymath2022.pdf](https://web.williams.edu/Mathematics/sjmillier/public_html/math/papers/ZeckGame_WinStrategies_Polymath2022.pdf)
- Accelerated Zeckendorf Games (with Diego Garcia-Fernandezsesma, Thomas Rascon, Risa Vandegrift, Ajmain Yamin), preprint.

Thank you!

## The Cookie Problem and Zeckendorf's Theorem

## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

**Cookie Monster** eats  $P - 1$  cookies:

## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

**Cookie Monster** eats  $P - 1$  cookies:  $\binom{C+P-1}{P-1}$  ways to do.

## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

**Cookie Monster** eats  $P - 1$  cookies:  $\binom{C+P-1}{P-1}$  ways to do.

Divides the cookies into  $P$  sets.



## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

**Cookie Monster** eats  $P - 1$  cookies:  $\binom{C+P-1}{P-1}$  ways to do.

Divides the cookies into  $P$  sets.

**Example:** 8 cookies and 5 people ( $C = 8$ ,  $P = 5$ ):

## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

**Cookie Monster** eats  $P - 1$  cookies:  $\binom{C+P-1}{P-1}$  ways to do.

Divides the cookies into  $P$  sets.

**Example:** 8 cookies and 5 people ( $C = 8$ ,  $P = 5$ ):



## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

**Cookie Monster** eats  $P - 1$  cookies:  $\binom{C+P-1}{P-1}$  ways to do.

Divides the cookies into  $P$  sets.

**Example:** 8 cookies and 5 people ( $C = 8, P = 5$ ):



## The Cookie Problem

### The Cookie Problem

The number of ways of dividing  $C$  identical cookies among  $P$  distinct people is  $\binom{C+P-1}{P-1}$ .

*Proof:* Consider  $C + P - 1$  cookies in a line.

**Cookie Monster** eats  $P - 1$  cookies:  $\binom{C+P-1}{P-1}$  ways to do.

Divides the cookies into  $P$  sets.

**Example:** 8 cookies and 5 people ( $C = 8, P = 5$ ):



## Preliminaries: The Cookie Problem: Reinterpretation

### Reinterpreting the Cookie Problem

The number of solutions to  $x_1 + \dots + x_p = C$  with  $x_i \geq 0$  is  $\binom{C+p-1}{p-1}$ .

Let  $p_{n,k} = \# \{N \in [F_n, F_{n+1}) : \text{the Zeckendorf decomposition of } N \text{ has exactly } k \text{ summands}\}$ .

For  $N \in [F_n, F_{n+1})$ , the **largest summand is  $F_n$** .

$$N = F_{i_1} + F_{i_2} + \dots + F_{i_{k-1}} + F_n,$$

$$1 \leq i_1 < i_2 < \dots < i_{k-1} < i_k = n, \quad i_j - i_{j-1} \geq 2.$$

$$d_1 := i_1 - 1, \quad d_j := i_j - i_{j-1} - 2 \quad (j > 1).$$

$$d_1 + d_2 + \dots + d_k = n - 2k + 1, \quad d_j \geq 0.$$

Cookie counting  $\Rightarrow p_{n,k} = \binom{n-2k+1-k-1}{k-1} = \binom{n-k}{k-1}$ .