# Why I love Monovariants: From Zombies to Conway's Soldiers to Fibonacci Games 

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## Invariants / Monovariants

Invariant: a quantity that is unchanged throughout the process / operations. (Big application: Noether's theorem).

Monovariant: a quantity that only changes in one direction throughout the process / operations. See

$$
\begin{gathered}
\text { https://howardhalim.com/math/Invariants\%20and\% } \\
\text { 20Monovariants.pdf }
\end{gathered}
$$

for a nice collection of problems.
Often a challenge to find a useful monovariant.

## Zombies

## Zombie Infection: Rules

- If share walls with 2 or more infected, become infected.
- Once infected, always infected.


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Initial Configuration
One moment Cater

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Initial Configuration One moment later


Iwo moments later

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Threemoments later

## Zombie Infection: Conquering The World

Easiest initial state that ensures all eventually infected is...?


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## Zombie Infection: Can $n-1$ infect all on an $n \times n$ board?



Perimeter of infection decreases by 4 .

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- If $n-1$ infected, maximum perimeter is $4(n-1)=4 n-4$.


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- How many must be safe?
- Other questions?


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- Perimeter of $n \times n$ square is $4 n$, so at least 1 square safe!
- How many must be safe?
- Other questions: Is a row safe? Higher dimensions? Other regions (torus)?


## Conway's Soldiers



Figure: John Horton Conway: Image from The Guardian.

## Conway's Soldiers / Checker Problem

Problem: Infinite checkerboard, pieces at all $(x, y)$ with $y \leq 0$. Using horizontal / vertical jumps (jumped piece gone forever), how high can you move a piece?


Figure: Left: A subset of the initial configuration. Right: moving a soldier / checker up 1.

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Figure: Left: A subset of the initial configuration. Right: moving a soldier / checker up 2. Can you do 3? 4? 5? Any height?

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Figure: Left: A subset of the initial configuration. Right: Also moving a soldier / checker up 2. Can you do 3? 4? 5? Any height?

## Conway's Soldiers: The Monovariant: I

Problem: Infinite checkerboard, pieces at all $(x, y)$ with $y \leq 0$. Using horizontal / vertical jumps (jumped piece gone forever), how high can you move a piece?


Figure: Conway's monovariant: What is it?

## Conway's Soldiers: The Monovariant: II

Choose target $T=(0,5)$.
Fix $x$ (to be determined later) and attach $x^{i+j}$ to a point that is $i$ units horizontally and $j$ units vertically from $T$.

|  | $\mathrm{x}^{7} \mathrm{x}^{6}$ |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $x^{6} x^{5}$ | $x^{4} \times$ |  | $x^{2}$ - | x | $x^{2} x^{3}$ | ${ }^{\text {x }}{ }^{\text {x }}{ }^{4}$ | ${ }^{5}$ | $x^{5} x^{6}$ |  |
|  | ${ }^{4}$ | $\mathrm{x}^{3}$ | $\mathrm{x}^{2}$ | $\times 1$ | T |  | $x^{2} x^{3}$ | ${ }^{3}{ }^{4}$ | $x^{4} \times$ |  |
| $\mathrm{x}^{7} \mathrm{x}^{6}$ | $x^{6} x^{5}$ | ${ }^{4}{ }^{4}$ | ${ }^{3}{ }^{3}$ | $x^{2}$ | $x{ }^{\text {x }}$ | $x^{2} x^{3}$ |  | ${ }^{4}{ }^{\text {x }}$ | ${ }^{5}$ |  |
| x | $x^{7} \chi^{6}$ | $\times^{5}$ | ${ }^{4}$ | ${ }^{3} \times$ | $x^{2} \times$ | $x^{3}{ }^{4}$ | $\mathrm{x}^{4} \mathrm{X}^{5}$ | ${ }^{5} x^{6}$ | ${ }^{6}{ }^{7}$ |  |
|  | $x^{8} \mathrm{x}^{7}$ | ${ }^{6} \times$ | ${ }^{5}$ | ${ }^{4} \times{ }^{3}$ | -3 | $x^{4} x^{5}$ | $x^{5} x^{6}$ | ${ }^{6}$ x ${ }^{7}$ |  |  |
|  |  |  |  |  |  |  |  |  |  |  |
|  | $x^{9} x^{8}$ |  |  |  | $\mathrm{x}^{4}$ x |  | ${ }^{6}$ | ${ }^{\text {x }}$ |  |  |
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## Conway's Soldiers: The Monovariant: III

Choose a target point $T$; for us it is a point of height 5 above the checkers: $T=(0,5)$.

Fix $x$ (to be determined later) and attach $x^{i+j}$ to a point that is $i$ units horizontally and $j$ units vertically from $T$.

What is the value of the initial board?

- Zeroth row: $\ldots, x^{7}, x^{6}, x^{5}, x^{6}, x^{7}, \ldots$ : sum is

$$
x^{5}+2 \sum_{k=6}^{\infty} x^{k}=x^{5}+\frac{2 x^{6}}{1-x}=\frac{(1+x) x^{5}}{1-x}
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$$

- Each row is $x$ times previous: Thus initial board value is

$$
\frac{(1+x) x^{5}}{1-x} \sum_{n=0}^{\infty} x^{n}=\frac{(1+x) x^{5}}{(1-x)^{2}}
$$

## Conway's Soldiers: The Monovariant: IV

Two moves: lose 2 pieces and add a piece further from $T$, or lose 2 pieces and add a piece closer to $T$.

First type of move clearly decreases value of board.

## Conway's Soldiers: The Monovariant: IV

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Figure: Moving pieces on $x^{6}$ and $x^{5}$ to on $x^{4}$.
Change is $x^{4}-x^{5}-x^{6}=x^{4}\left(1-x-x^{2}\right)$, want this to be zero.

## Conway's Soldiers: The Monovariant: IV

Two moves: lose 2 pieces and add a piece further from $T$, or lose 2 pieces and add a piece closer to $T$.

Second type replaces $x^{n+2}$ and $x^{n+1}$ with an $x^{n}$ : change is $x^{n}-x^{n+1}-x^{n+2}$. Choose $x$ so that this change is zero.

Thus $1-x-x^{2}=0$ or $x=(-1 \pm \sqrt{5}) / 2$. Take positive root, $(-1+\sqrt{5}) / 2=\varphi-1(\varphi$ the golden mean $)$.

Monovariant: sum of the values of squares with checkers.

## Conway's Soldiers: The Monovariant: V

Choose a target point $T$.

- Initial board value is

$$
\frac{(1+x) x^{5}}{(1-x)^{2}}: \text { when } x=\frac{\sqrt{5}-1}{2} \text { get } 1 .
$$

- Target at $(0,4)$ contributes $x=\frac{\sqrt{5}-1}{2} \approx 0.618034$; as less than 1 possible (and can be done).
- Target at $(0,5)$, board's value at least 1 . Moves never increase value: IMPOSSIBLE IN FINITE TIME!¹


## New Results

Conway Checkers m-game: Start with $m$ checkers on each gridpoint (original game is just 1), if jump over it lose one checker.

## SMALL 2024

Given a Conway Checkers $m$-game, the maximum row attainable, $n_{m}$, satisfies

$$
\left\lfloor\log _{\varphi}(m)+4.67\right\rfloor \leq n_{m} \leq\left\lfloor\log _{\varphi}(m)+5\right\rfloor
$$

for sufficiently large $m$, where $\varphi$ is the golden ratio $\frac{\sqrt{5}+1}{2}$.

## Zeckendorf Minimality

## Introduction: Summand Minimality

Fibonaccis: $F_{1}=1, F_{2}=2, F_{3}=3, F_{4}=5, F_{n+2}=F_{n+1}+F_{n}$.

## Zeckendorf's Theorem

Every positive integer can be written uniquely as a sum of one or more non-consecutive Fibonacci numbers.

Example: $2024=1597+377+34+13+3=F_{16}+F_{13}+F_{8}+F_{6}+F_{3}$.

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Conversely, we can construct the Fibonacci sequence using this property:
$1,2,3,5,8,13 \ldots$

## Summand Minimality

## Example

- $18=13+5=F_{6}+F_{4}$, legal decomposition, two summands.
- $18=13+3+2=F_{6}+F_{3}+F_{2}$, non-legal decomposition, three summands.


## Theorem

The Zeckendorf decomposition is summand minimal.

## Overall Question

What other recurrences are summand minimal?

## Zeckendorf Decomposition is Minimal

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If $n=\sum_{k} a_{k} F_{k}$ (with $a_{k}$ non-negative integers), define the weighted index attached to this decomposition $\mathcal{D}$ to be $\operatorname{Index}(\mathcal{D})=\sum_{k} a_{k} \sqrt{k}$.

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More natural $\sum_{k} a_{k} k$ but square-root makes strictly decreasing.
Bounded process: For fixed $n$, only indices up to certain point used, and $a_{k} \leq n$.

## Zeckendorf Decomposition is Minimal: Proof

Show $\operatorname{Index}(\mathcal{D})$ is a mono-variant, end in the Zeckendorf decomposition, number summands never increased.

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$$
\begin{aligned}
& F_{k} \wedge F_{k+1} \rightarrow F_{k+2}: \\
& \quad \bullet \sqrt{k}+\sqrt{k+1}>\sqrt{k+2}
\end{aligned}
$$

$2 F_{k} \rightarrow F_{k-2}+F_{k+1}:$

- $k \geq 3: 2 \sqrt{k}>\sqrt{k-2}+\sqrt{k+1}$
- $k=2: 2 \sqrt{2}>\sqrt{1}+\sqrt{3}$
- $k=1: 2 \sqrt{1}>\sqrt{2}$

Only finitely many values, each move lowers, continue till hit Zeckendorf, number of summands never increased.

## Positive Linear Recurrence Sequences

## Definition

A positive linear recurrence sequence (PLRS) is a sequence given by a recurrence $\left\{a_{n}\right\}$ with

$$
a_{n}:=c_{1} a_{n-1}+\cdots+c_{t} a_{n-t}
$$

and each $c_{i} \geq 0$ and $c_{1}, c_{t}>0$. We use ideal initial conditions $a_{-(n-1)}=0, \ldots, a_{-1}=0, a_{0}=1$ and call $\left(c_{1}, \ldots, c_{t}\right)$ the signature of the sequence.

## Theorem (Cordwell, Hlavacek, Huynh, M., Peterson, Vu)

For a PLRS with signature $\left(c_{1}, c_{2}, \ldots, c_{t}\right)$, the Generalized Zeckendorf Decompositions are summand minimal if and only if

$$
c_{1} \geq c_{2} \geq \cdots \geq c_{t}
$$

## Zeckendorf Games

## Fibonacci Game: Rules

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- A turn is one of the following moves:
$\diamond$ If have two pieces on $F_{k}$ can remove and put one piece at $F_{k+1}$ and one at $F_{k-2}$
(if $k=1$ then $2 F_{1}$ becomes $1 F_{2}$ )
$\diamond$ If pieces at $F_{k}$ and $F_{k+1}$ remove and add one at $F_{k+2}$.


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$\diamond$ If pieces at $F_{k}$ and $F_{k+1}$ remove and add one at $F_{k+2}$.

Questions:

- Does the game end? How long?
- For each $N$ who has the winning strategy?
- What is the winning strategy?


## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 10 | 0 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 1: $F_{1}+F_{1}=F_{2}$

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 8 | 1 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 2: $F_{1}+F_{1}=F_{2}$

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 6 | 2 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 1: $2 F_{2}=F_{3}+F_{1}$

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 7 | 0 | 1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 2: $F_{1}+F_{1}=F_{2}$

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 5 | 1 | 1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 1: $F_{2}+F_{3}=F_{4}$.

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 5 | 0 | 0 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 2: $F_{1}+F_{1}=F_{2}$.

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 3 | 1 | 0 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 1: $F_{1}+F_{1}=F_{2}$.

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 1 | 2 | 0 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 2: $F_{1}+F_{2}=F_{3}$.

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 0 | 1 | 1 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

Next move: Player 1: $F_{3}+F_{4}=F_{5}$.

## Sample Game

Start with 10 pieces at $F_{1}$, rest empty.

| 0 | 1 | 0 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: |
| $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

No moves left, Player One wins.

## Sample Game

Player One won in 9 moves.

| $(1)$ | 10 | 0 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $(2)$ | 8 | 1 | 0 | 0 | 0 |
| $(1)$ | 6 | 2 | 0 | 0 | 0 |
| $(2)$ | 7 | 0 | 1 | 0 | 0 |
| $(1)$ | 5 | 1 | 1 | 0 | 0 |
| $(2)$ | 5 | 0 | 0 | 1 | 0 |
| $(1)$ | 3 | 1 | 0 | 1 | 0 |
| $(2)$ | 1 | 2 | 0 | 1 | 0 |
|  |  |  |  |  |  |
| $(1)$ | 0 | 1 | 1 | 1 | 0 |
|  | 0 | 1 | 0 | 0 | 1 |
|  | $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

## Sample Game

Player Two won in 10 moves.

| $(1)$ | 10 | 0 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $(2)$ | 8 | 1 | 0 | 0 | 0 |
| $(1)$ | 6 | 2 | 0 | 0 | 0 |
| $(2)$ | 7 | 0 | 1 | 0 | 0 |
| $(1)$ | 5 | 1 | 1 | 0 | 0 |
| $(2)$ | 5 | 0 | 0 | 1 | 0 |
| $(1)$ | 3 | 1 | 0 | 1 | 0 |
| $(2)$ | 1 | 2 | 0 | 1 | 0 |
| $(1)$ | 2 | 0 | 1 | 1 | 0 |
| $(2)$ | 0 | 1 | 1 | 1 | 0 |
|  | 0 | 1 | 0 | 0 | 1 |
|  | $\left[F_{1}=1\right]$ | $\left[F_{2}=2\right]$ | $\left[F_{3}=3\right]$ | $\left[F_{4}=5\right]$ | $\left[F_{5}=8\right]$ |

## Games end

## Theorem

All games end in finitely many moves.

Proof: The sum of the square roots of the indices is a strict monovariant.

- Adding consecutive: $(\sqrt{k}+\sqrt{k+1})-\sqrt{k+2}>0$.
- Splitting: $2 \sqrt{k}-(\sqrt{k+1}+\sqrt{k+1})>0$.
- Spitting 1 's: $2 \sqrt{1}-\sqrt{2}>0$.
- Splitting 2's: $2 \sqrt{2}-(\sqrt{3}+\sqrt{1})>0$.


## Games Lengths: I

Upper bound: At most $3 n-3 Z(n)-I(n)+1$ moves

- $Z(n)$ is the number of terms in the Zeckendorf decomposition,
- $I(n)$ is the sum of the indices.

Fastest game: $n-Z(n)$ moves $(Z(n)$ is the number of summands in $n$ 's Zeckendorf decomposition).

From always moving on the largest summand possible (deterministic).

## Games Lengths: II



Figure: Frequency graph of the number of moves in 9,999 simulations of the Zeckendorf Game with random moves when $n=60$ vs a Gaussian. Natural conjecture....

## Winning Strategy

## Theorem

Player Two Has a Winning Strategy

Idea is to show if not, Player Two could steal Player One's strategy.

Non-constructive!

Will highlight idea with a simpler game.

## Winning Strategy: Intuition from Dot Game

Two players, alternate. Turn is choosing a dot at $(i, j)$ and coloring every dot $(m, n)$ with $i \leq m$ and $j \leq n$.

Once all dots colored game ends; whomever goes last loses.
Prove Player 1 has a winning strategy!


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## Sketch of Proof for Player Two's Winning Strategy



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## The Bergman Game

## Definition

The Bergman Game is played with the standard split/combine moves from the Zeckendorf game, but on a two-sided infinite tape instead of a one-sided infinite tape. It produces base- $\varphi$ decompositions $(\varphi=(1+\sqrt{5}) / 2)$.

## Example

| 0 | 0 | 4 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- |
| 1 | 0 | 2 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 |

$4=\varphi^{-2}+\varphi^{0}+\varphi^{2}$.

## The Bergman Game

## Theorem (Baily, Dell, Durmić, Fleischmann, Jackson, Mijares, M., Pesikoff, Reifenberg, Reina, Yang)

The longest Bergman Game with $n$ summands terminates in $\Theta\left(n^{2}\right)$ time regardless of where the summands are placed. The shortest possible Bergman Game terminates in $\Theta(n)$ time.

Natural Question: Who has the winning strategy?

- Not currently known.
- Game tree explodes, escaping a strategy steal.


## The Frodnekcez Game (Reverse Zeckendorf Game): Rules

- The Zeckendorf game in reverse, last to move wins.
- Bins $F_{1}, F_{2}, F_{3}, \ldots$, for some natural number $N$, start with one piece in bin $F_{k}$ if $F_{k}$ is in the Zeckendorf decomposition of $N$, and have other bins empty.
- A turn is one of the following moves:
$\diamond$ If one piece at $F_{k+1}$ and one at $F_{k-2}$, can remove and add two pieces on $F_{k}$.
$\diamond$ If piece at $F_{k+2}$, remove and add one piece at both $F_{k}$ and $F_{k+1}$.
( $F_{1}$ and $F_{3}$ becomes $2 F_{2}$, and $F_{2}$ becomes $2 F_{1}$ )
Problem created and analyzed by PANTHers 2023 from the 2023 SMALL REU: Zoe Batterman, Aditya Jambhale, Akash Narayanan, Kishan Sharma, Andrew Yang, Chris Yao.


## Winning Strategy?



Figure: In the forward Zeckendorf game, Player 2 wins for all $N>2$. The reverse game is more interesting. Natural conjecture...

## Current / Future Work

- What if $p \geq 3$ people play the Fibonacci game? Some multi-player results.
- Does the number of moves in random games converge to a Gaussian? Evidence....
- How long do games take? Proved closed interval.
- Accelerated games: do as many of one move as wish....
- What of other recurrences?
\$500 Prize: Determine the winning strategy.


## Thanks / References

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## Thank you!

## The Cookie Problem and Zeckendorf's Theorem

## The Cookie Problem

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The number of ways of dividing $C$ identical cookies among $P$ distinct people is $\binom{C+P-1}{P-1}$.

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Example: 8 cookies and 5 people ( $C=8, P=5$ ):

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## Preliminaries: The Cookie Problem: Reinterpretation

## Reinterpreting the Cookie Problem

The number of solutions to $x_{1}+\cdots+x_{P}=C$ with $x_{i} \geq 0$ is $\binom{C+P-1}{P-1}$.

Let $p_{n, k}=\#\left\{N \in\left[F_{n}, F_{n+1}\right)\right.$ : the Zeckendorf decomposition of $N$ has exactly $k$ summands $\}$.

For $N \in\left[F_{n}, F_{n+1}\right)$, the largest summand is $F_{n}$.

$$
\begin{gathered}
N=F_{i_{1}}+F_{i_{1}}+\cdots+F_{i_{k-1}}+F_{n}, \\
1 \leq i_{1}<i_{2}<\cdots<i_{k-1}<i_{k}=n, i_{j}-i_{j-1} \geq 2 . \\
d_{1}:=i_{1}-1, d_{j}:=i_{j}-i_{j-1}-2(j>1) . \\
d_{1}+d_{2}+\cdots+d_{k}=n-2 k+1, d_{j} \geq 0 .
\end{gathered}
$$

Cookie counting $\Rightarrow p_{n, k}=\binom{n-2 k+1-k-1}{k-1}=\binom{n-k}{k-1}$.

