# Inrtoduwtion to Erorr Dwtetcion and Erorr Czrrectmon 

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# Introduction to Error Detection and Error Correction 

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## Introduction

## Cryptography Basics

Enough to send 0's and 1's:

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\begin{aligned}
\diamond A=00000, & B=00001, \quad C=00010, \ldots \\
Z=11010, & 0=11011, \quad 1=11100, \ldots .
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Two major issues:
$\diamond$ Transmit message so only desired recipient can read.
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Secret: $p=15217, q=17569, d=80998505$.
Public: $N=p q=267347473, e=3141593$.
Note: ed $=1 \bmod (p-1)(q-1)$.
Message: $M=195632041$, send $M^{e} \bmod N$ or $X=121209473$.
Decrypt: $X^{d} \bmod N$ or 195632041.

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Message: $M=195632041$, send $M^{e} \bmod N$ or $X=121209473$.
Decrypt: $X^{d} \bmod N$ or 195632041.
Imagine receive $\widetilde{X}=121209483$.
Message 195632041
Decrypts 121141028, only two digits are the same!

## Outline

Will concentrate on Error Detection and Correction.

- Detection: Check Digit
- Correction: Majority Rules and Generalization


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## Next Steps

More involved methods detecting more: The Verhoeff algorithm catches single digit errors and flipping adjacent digits: https://en.wikipedia.org/wiki/
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Verhoeff_algorithm.
Want to detect where the error is: Tell me twice!

## © ©ゃ○๑

## Majority Rules

## Tell Me Three Times

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Crucially uses binary outcome: https://www . youtube.com/watch?v=RerJWv5vwxc and https:// www. youtube.com/watch?v=vWCGs27_xPI.

What is the problem with this method?

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What is the problem with this method? Only one-third is information.

How can we do better?

## Tell Me $n$ Times

## (1) (1) (1) (1) (0) (2) ๑

Tell Me Four Times: only $25 \%$ of message is data (general case just $1 / n$ ).

Want to correct errors but still send a lot of information.
What's a success?

## Tell Me $n$ Times

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Tell Me Four Times: only $25 \%$ of message is data (general case just $1 / n$ ).

Want to correct errors but still send a lot of information.
What's a success? Greater than $50 \%$ is data.

## Tell Me Three Times (revisited)

## Let's revisit Tell Me Three Times:



How should we do two data points?
How many check digits do you expect?

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## Two of Five

This is better: 2 of 5 or $40 \%$ of message is data!


Unfortunately still below 50\%.
How many data points should we try next: $3,4,5, \ldots$ ?

## Three and Four Bits of Data



Which is better?

## Three and Four Bits of Data



Which is better? Both $50 \%$ but fewer needed with triangle.
What should we do next: $5,6,7,8,9, \ldots$ ?

## Triangle and Square Numbers

$$
T_{n}=n(n+1) / 2 \text { and } S_{n}=n^{2} .
$$



Both give $60 \%$ of the message is data. Can we continue?
Data on exactly two lines, check bits on one.

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## Triangle and Square Systems



Triangle: $T_{n}=n(n+1) / 2$ data, $n+1$ check, so $(n+2)(n+1) / 2$ bits total and $n /(n+2)$ information.

Square: $S_{n}=n^{2}$ data, $2 n$ check, so $n^{2}+2 n$ bits total and $n /(n+2)$ information.

## Triangle and Square Systems



Can get as high a percentage information as desire, at a cost of longer string (and thus more likely to have two errors).

## Generalizations

What is a better geometry to use?

## Generalizations

$2 \times 2 \times 2$ : 8 data points, 6 check bits (for planes): info is $8 / 14 \approx 57 \%$.
$3 \times 3 \times 3$ : 27 data points, 9 check bits (for planes): info is $27 / 36=75 \%$.

For $6 \times 6$ data square info is $36 / 48=75 \%$, for $T_{7}$ is $28 / 36 \approx 77.78 \%$.

## Generalizations


$4 \times 4 \times 4$ : 64 data points, 12 check bits: info is $64 / 76 \approx 84.21 \%$.

For $9 \times 9$ data square info is $81 / 99 \approx 81.82 \%$.
For $T_{11}$ triangle: 66 data points, info is $66 / 79 \approx 83.54 \%$.

## Generalizations


$n \times n \times n: n^{3}$ data points, $3 n$ check bits: info is $n^{2} /\left(n^{2}+3\right)$.
Better percentage is information for large $n$; how should we generalize?

## Other Approaches

Hamming Codes: Can send a message with 7 bits, 4 are data, and can correct one error: https://en. wikipedia.org/wiki/Hamming_code.

Extended binary Golay code: Can send a message with 24 bits, 12 are data, can correct any 3-bit errors and can detect some other errors: https://en.wikipedia. org/wiki/Binary_Golay_code.

## Manhamming



- If no errors, all correct.
- If only one color error, is P1, P2 or P3.
- If just blue and orange is D1.
- If just blue and green is D2.
- If just orange and green is D3
- If all wrong is D4.


## Comparison

Say want to transmit around $2^{12}=4096$ bits of data.
Can do a square and cube; the Hamming code will do $2^{12}-1-12$.

- Square: 4096 out of 4224 data: $96.9697 \%$.
- Cube: 4096 out of 4144 data: 98.8417\%.
- Hamming: 4083 out of 4095 data: 99.707\%.

All converge to $100 \%$, difference narrows as size increases.

## Interleaving

Say transmit
01111010010101010101010101010101010101011110...
but a localized burst of noise, receive
01110111010101010101010101010101010101011110...

## Interleaving

Transmit every fourth:

- $01000000001 \mapsto 00000000001$
- $1011111111 \rightarrow 1111111111$
- $11000000001 \mapsto 11000000001$
- $10111111110 \mapsto 11111111110$


## Steganography

## Can you see the cat in the tree?



## Transmitting Images

How to transmit an image?

- Have an $L \times W$ grid with $L W$ pixels.
- Each pixel a triple, maybe (Red, Green, Blue).
- Often each value in $\left\{0,1,2,3, \ldots, 2^{n}-1\right\}$.
- $n=8$ gives 256 choices for each, or 16,777,216 possibilities.


## Steganography

Steganography: Concealing a message in another message: https://en.wikipedia.org/wiki/
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Take one of the colors, say red, a number from 0 to 255.
Write in binary: $r_{7} 2^{7}+r_{6} 2^{6}+\cdots+r_{1} 2+r_{0}$.
If change just the last or last two digits, very minor change to image.

Can hide an image in another.
If just do last, can hide a black and white image easily....

## Can you see the cat in the tree?



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