Check Digit

Majority Rules

Steganography

Inrtoduwtion to Erorr Dwtetcion and Erorr Czrrectmon

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Introduction to Error Detection and Error Correction

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Burlington, June 19, 2019

Introduction	Check Digit	Majority Rules	Steganography
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Introduction	Check Digit	Majority Rules ooooooooooooooo	Steganography 00000
Cryptography	Basics		

Enough to send 0's and 1's: $\diamond A = 00000, B = 00001, C = 00010, \dots$ $Z = 11010, 0 = 11011, 1 = 11100, \dots$

Two major issues:

- Transmit message so only desired recipient can read.
- Ensure correct message received.

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Cryptography Ba	sies		

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Bit Error Dangers	: RSA		

If receive wrong bit in RSA, message completely different.



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Secret: p = 15217, q = 17569, d = 80998505. Public: N = pq = 267347473, e = 3141593. Note: $ed = 1 \mod (p - 1)(q - 1)$. Message: M = 195632041, send $M^e \mod N$ or X = 121209473. Decrypt: $X^d \mod N$ or 195632041.



If receive wrong bit in RSA, message completely different.

Secret: p = 15217, q = 17569, d = 80998505. Public: N = pq = 267347473, e = 3141593. Note: $ed = 1 \mod (p - 1)(q - 1)$. Message: M = 195632041, send $M^e \mod N$ or X = 121209473. Decrypt: $X^d \mod N$ or 195632041.

Imagine receive $\hat{X} = 121209483$. Message 195632041 Decrypts 121141028, only two digits are the same!

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Outline			

Will concentrate on Error Detection and Correction.

- Detection: Check Digit
- Correction: Majority Rules and Generalization

Introduction 0000	Check Digit ●○○	Majority Rules	Steganography

Check Digit

Introduction	Check Digit	Majority Rules	Steganography
0000	○●○		00000
Check Digit			

Introduction	Check Digit	Majority Rules	Steganography
0000	○●○		00000
Check Digit			

Think scanner at a supermarket....



Think scanner at a supermarket....



Last digit makes sum 0 mod 10 (or 0 mod 2).

Introduction	Check Digit	Majority Rules	Steganography
0000	○●○	০০০০০০০০০০০০০০	00000
Check Digit			

Think scanner at a supermarket....



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More involved methods detecting more: The Verhoeff algorithm catches single digit errors and flipping adjacent digits: https://en.wikipedia.org/wiki/ Verhoeff_algorithm.

Want to detect where the error is:



More involved methods detecting more: The Verhoeff algorithm catches single digit errors and flipping adjacent digits: https://en.wikipedia.org/wiki/ Verhoeff_algorithm.

Want to detect where the error is: Tell me twice!

(1) (1) OR (0) (0)

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		000000000000	

Majority Rules

Introduction	Check Digit	Majority Rules a●ooooooooooooo	Steganography
Tell Me Three	Times		

Tell Me Three Times detects and *probably* corrects (need probability of an error small).

Introduction	Check Digit	Majority Rules	Steganography
0000		○●○○○○○○○○○○	00000
Tell Me Three	Times		

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Introduction	Check Digit	Majority Rules	Steganography
0000		○○●○○○○○○○○○	00000
	Timos		

Crucially uses binary outcome: https://www. youtube.com/watch?v=RerJWv5vwxc and https:// www.youtube.com/watch?v=vWCGs27_xPI.

What is the problem with this method?

Introduction	Check Digit	Majority Rules	Steganography
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Tell Me Three	Times		

Crucially uses binary outcome: https://www. youtube.com/watch?v=RerJWv5vwxc and https:// www.youtube.com/watch?v=vWCGs27_xPI.

What is the problem with this method? Only one-third is information.

How can we do better?

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Tell Me <i>n</i> Times			



Tell Me Four Times: only 25% of message is data (general case just 1/n).

Want to correct errors but still send a lot of information.

What's a success?

Introduction	Check Digit	Majority Rules	Steganography
0000	000	○○○●○○○○○○○○	00000
Tell Me <i>n</i> Times			



Tell Me Four Times: only 25% of message is data (general case just 1/n).

Want to correct errors but still send a lot of information.

What's a success? Greater than 50% is data.

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Tell Me Three Times (revisited)

Let's revisit Tell Me Three Times:



How should we do two data points? How many check digits do you expect?

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Tell Me Three Times (revisited)

Let's revisit Tell Me Three Times:



How should we do two data points? How many check digits do you expect?



Introduction	Check Digit	Majority Rules	Steganography
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Two of Five			

This is better: 2 of 5 or 40% of message is data!



Unfortunately still below 50%.

How many data points should we try next: 3, 4, 5, ...?

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Three and Four Bits of Data





Which is better?

Introduction	

Check Digit

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Three and Four Bits of Data



Which is better? Both 50% but fewer needed with triangle.

What should we do next: 5, 6, 7, 8, 9, ...?



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Triangle and Square Numbers

$$T_n = n(n+1)/2$$
 and $S_n = n^2$.



Both give 60% of the message is data. Can we continue?

Data on exactly two lines, check bits on one.

Check Digit

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Triangle and Square Numbers

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Introduction	

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Triangle and Square Systems



Triangle: $T_n = n(n+1)/2$ data, n+1 check, so (n+2)(n+1)/2 bits total and n/(n+2) information.

Square: $S_n = n^2$ data, 2n check, so $n^2 + 2n$ bits total and n/(n+2) information.

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Triangle and Square Systems



Can get as high a percentage information as desire, at a cost of longer string (and thus more likely to have two errors).

Introduction	Check Digit	Majority Rules	Steganography
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Generalizations			

What is a better geometry to use?

Introduction	Check Digit	Majority Rules	Steganography
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Generalizations			





 $2\times 2\times 2$: 8 data points, 6 check bits (for planes): info is $8/14\approx 57\%.$

 $3\times3\times3$: 27 data points, 9 check bits (for planes): info is 27/36=75%.

For 6×6 data square info is 36/48 = 75%, for T_7 is $28/36 \approx 77.78\%$.

Introduction	Check Digit	Majority Rules	Steganography
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Generalizations			



 $4 \times 4 \times 4$: 64 data points, 12 check bits: info is $64/76 \approx 84.21\%$.

For 9×9 data square info is $81/99 \approx 81.82\%$.

For T_{11} triangle: 66 data points, info is $66/79 \approx 83.54\%$.

Introduction	Check Digit	Majority Rules	Steganography
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Generalizations			



 $n \times n \times n$: n^3 data points, 3n check bits: info is $n^2/(n^2+3)$.

Better percentage is information for large *n*; how should we generalize?

Introduction	Check Digit	Majority Rules	Steganography
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Other Approache	e		

Hamming Codes: Can send a message with 7 bits, 4 are data, and can correct one error: https://en.wikipedia.org/wiki/Hamming_code.

Extended binary Golay code: Can send a message with 24 bits, 12 are data, can correct any 3-bit errors and can detect some other errors: https://en.wikipedia.org/wiki/Binary_Golay_code.

Introduction 0000	Check D 000	igit	Majority Rules	000	Steganography 00000
Manhammi	ng				
	3 (D1)	5 (D2)	6 (D3)	7 (D4)	

1 (P1) 2 (P2) 4 (P3)

- If no errors, all correct.
- If only one color error, is P1, P2 or P3.
- If just blue and orange is D1.
- If just blue and green is D2.
- If just orange and green is D3
- If all wrong is D4.



Say want to transmit around $2^{12} = 4096$ bits of data.

Can do a square and cube; the Hamming code will do $2^{12} - 1 - 12$.

- Square: 4096 out of 4224 data: 96.9697%.
- Cube: 4096 out of 4144 data: 98.8417%.
- Hamming: 4083 out of 4095 data: 99.707%.

All converge to 100%, difference narrows as size increases.

Introduction	Check Digit	Majority Rules	Steganography
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Interleaving			

Say transmit

but a localized burst of noise, receive

Introduction	Check Digit	Majority Rules	Steganography
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Interleaving			

Transmit every fourth:

- 0100000001 \mapsto 0000000001
- 10111111111 \mapsto 1111111111
- 1100000001 → 1100000001
- 10111111110 \mapsto 11111111110



Steganography

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Check Dig

Majority Rules

Steganography

Can you see the cat in the tree?





How to transmit an image?

- Have an $L \times W$ grid with LW pixels.
- Each pixel a triple, maybe (Red, Green, Blue).
- Often each value in $\{0, 1, 2, 3, ..., 2^n 1\}$.
- n = 8 gives 256 choices for each, or 16,777,216 possibilities.

Introduction	Check Digit	Majority Rules	Steganography
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Steganography			

Steganography: Concealing a message in another message: https://en.wikipedia.org/wiki/ Steganography.



Steganography: Concealing a message in another message: https://en.wikipedia.org/wiki/ Steganography.

Take one of the colors, say red, a number from 0 to 255.

Write in binary: $r_7 2^7 + r_6 2^6 + \cdots + r_1 2 + r_0$.

If change just the last or last two digits, very minor change to image.

Can hide an image in another.

If just do last, can hide a black and white image easily....

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Can you see the cat in the tree?



Check Dig

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Can you see the cat in the tree?



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