Generalized Ramanujan Primes

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Prime Numbers

Introduction

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- Exact spacing of primes unknown but roughly growing logarithmically (Prime Number Theorem).
- We expect linearly increasing intervals to contain an increasing amount of primes.

Historical Introduction

Introduction

Bertrand's Postulate (1845)

For all integers x > 2, there exists at least one prime in (x/2, x].

Definition

The *n*-th Ramanujan prime R_n : smallest integer such that for any $x > R_n$, at least *n* primes in any (x/2, x].

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Ramanujan Primes

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- Sondow: $R_n \sim p_{2n}$.

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Theorem

- Ramanujan: For each integer n, R_n exists.
- Sondow: $R_n \sim p_{2n}$.
- Sondow: As $n \to \infty$, 50% of primes are Ramanujan.

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• For each *c* and integer *n*, does $R_{c,n}$ exist? Yes!

Definition

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What percent of primes are c-Ramanujan?

Given $R_{c,n} \sim p_{\frac{n}{1-c}}$, it follows $R_{c,N(1-c)} \sim p_N$.

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$$R_{c,n} \sim p_{\frac{n}{1-c}}$$
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1-c

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Preliminaries

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The Prime Number Theorem states:

$$\lim_{x\to\infty}\frac{\pi(x)}{x/\log(x)}=1.$$

The logarithmic integral function Li(x) is defined by

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The Prime Number Theorem

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The Prime Number Theorem gives us

$$\pi(x) = \operatorname{Li}(x) + O\left(\frac{x}{\log^2 x}\right),$$

i.e., there is a C > 0 such that for all x sufficiently large

$$-C\frac{x}{\log x} \leq \pi(x) - \operatorname{Li}(x) \leq C\frac{x}{\log x}.$$

Theorem (Amersi, Beckwith, Ronan, 2011)

For all $n \in \mathbb{Z}$ and all $c \in (0, 1)$, the n-th c-Ramanujan prime $R_{c,n}$ exists.

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Proof:

• The number of primes in (cx, x] is $\pi(x) - \pi(cx)$.

Existence of $R_{c,n}$

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- The number of primes in (cx, x] is $\pi(x) \pi(cx)$.
- Using the Prime Number Theorem and Mean Value Theorem:

$$\pi(x) - \pi(cx) = \operatorname{Li}(x) - \operatorname{Li}(cx) + O(x \log^{-2} x)$$

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$$\pi(x) - \pi(cx) = \operatorname{Li}(x) - \operatorname{Li}(cx) + O(x \log^{-2} x)$$
$$= \operatorname{Li}'(y_c)(x - cx) + O(x \log^{-2} x)$$
$$\text{with } y_c(x) \in [cx, x].$$

•
$$\pi(x) - \pi(cx) = \frac{(1-c)x}{\log y_c} + O(x \log^{-2} x).$$

Existence of $R_{c,n}$

Introduction

- $\pi(x) \pi(cx) = \frac{(1-c)x}{\log x} + O(x \log^{-2} x).$
- Since $\log y_c = \log x b_c$, for $b_c \in [0, -\log c]$,

$$\pi(x) - \pi(cx) = \frac{(1-c)x}{\log x - b_c} + O\left(\frac{x}{\log^2 x}\right).$$

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$$\pi(x) - \pi(cx) = \frac{(1-c)x}{\log x - b_c} + O\left(\frac{x}{\log^2 x}\right).$$

• For sufficiently large x, $\pi(x) - \pi(cx)$ is strictly increasing and $\pi(x) - \pi(cx) \ge n$, for all integers n.

Bounds on $R_{c,n}$

• $R_{c,n} \geq p_n$

Bounds on $\overline{R_{c,n}}$

• $R_{c,n} \ge p_n \ge n \log n$.

Bounds on $R_{c,n}$

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- $R_{c,n} \geq p_n \geq n \log n$.
- Want $R_{c,n} \leq \alpha_c n \log(\alpha_c n)$ for large n.

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- Conjecture $\alpha_c = \frac{2}{1-c}$.

- $R_{c,n} > p_n > n \log n$.
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- To prove, must show $\pi(x) \pi(cx) \ge n \ \forall x \ge \alpha_c n \log(\alpha_c n)$.

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- $n_{c,n} \geq p_n \geq m \log n$.
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- Conjecture $\alpha_c = \frac{2}{1-c}$.
- To prove, must show $\pi(x) \pi(cx) \ge n \ \forall x \ge \alpha_c n \log(\alpha_c n)$.
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- To prove, must show $\pi(x) - \pi(cx) > n \ \forall x > \alpha_c n \log(\alpha_c n).$
- We then obtain bounds on $\log R_{c,n}$:

$$\left(1 - \frac{\beta_c \log \log n}{\log n}\right) \log n \leq \log R_{c,n} \leq \left(1 + \frac{\beta_c \log \log n}{\log n}\right) \log n.$$

Results

Asymptotic Behavior

Introduction

Theorem (Amersi, Beckwith, Ronan, 2011)

For any fixed $c \in (0, 1)$, the *n*th *c*-Ramanujan prime is asymptotic to the $\frac{n}{1-c}$ th prime as $n \to \infty$

$$\begin{aligned} \left| R_{c,n} - p_{\frac{n}{1-c}} \right| & \leq \left| \left| R_{c,n} - \frac{n}{1-c} \log R_{c,n} \right| + \left| \frac{n}{1-c} \log R_{c,n} - \frac{n}{1-c} \log n \right| \\ & + \left| \frac{n}{1-c} \log n - \frac{n}{1-c} \log \frac{n}{1-c} \right| \\ & + \left| \frac{n}{1-c} \log n - p_{\frac{n}{1-c}} \right| \end{aligned}$$

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For any fixed $c \in (0,1)$, the nth c-Ramanujan prime is asymptotic to the $\frac{n}{1-c}$ th prime as $n \to \infty$

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Since
$$\frac{n \log \log n}{n_0} \to 0$$
 as $n \to \infty$

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Since
$$\frac{n \log \log n}{p_n} \to 0$$
 as $n \to \infty \Rightarrow R_{c,n} \sim p_{\frac{n}{1-c}}$

Results

Frequency of c-Ramanujan Primes

Theorem (Amersi, Beckwith, Ronan, 2011)

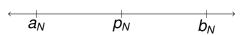
In the limit, the probability of a generic prime being a c-Ramanujan prime is 1 - c.

• Define $N = \lfloor \frac{n}{1-c} \rfloor$.

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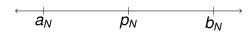


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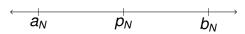
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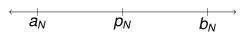


- Worst cases:
 - $R_{c,n} = a_N$ and every prime in $(a_N, p_N]$ is *c*-Ramanujan,

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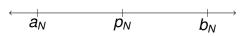


- Worst cases:
 - $R_{c,n} = a_N$ and every prime in $(a_N, p_N]$ is c-Ramanujan,
 - $R_{c,n} = b_N$ and every prime in $[p_N, b_N)$ is c-Ramanujan.

Theorem (Amersi, Beckwith, Ronan, 2011)

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- Worst cases:
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 - $R_{c,n} = b_N$ and every prime in $[p_N, b_N)$ is *c*-Ramanujan.
- Goal: $\frac{\pi(b_N)-\pi(a_N)}{\pi(p_N)} \to 0$ as $N \to \infty$.

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Prime Number Theorem

$$\Rightarrow \frac{\pi(b_N) - \pi(a_N)}{\pi(p_N)} \le \xi \frac{\log \log N}{\log N} \to 0 \text{ as } N \to \infty$$

Prime Numbers

101 103 107 113 127 131 137 139 149 157 163 167 173 179 181 193 197 199 211 223 227

Ramanujan Primes

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We define

Introduction

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Results

Coin Flipping Model (Variation on Cramer Model)

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- N, the number of trials,

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$$\mathbb{E}[L_N] \quad \approx \quad \frac{\log N}{\log(1/P)} - \left(\frac{1}{2} - \frac{\log(1-P) + \gamma}{\log(1/P)}\right)$$

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- $\gamma = 0.5772...$, the Euler-Mascheroni constant,
- P, the probability of Heads,
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$$\mathbb{E}[L_N] \approx \frac{\log N}{\log(1/P)} - \left(\frac{1}{2} - \frac{\log(1-P) + \gamma}{\log(1/P)}\right)$$

$$\operatorname{Var}[L_N] \approx \frac{\pi^2}{6\log^2(1/P)} + \frac{1}{12}$$

What is P?

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- As $N \to \infty$, $P_c = 1 c$,
- For finite intervals [a, b], P_c is a function of a and b,
- Choose $a = 10^5$, $b = 10^6$.

Distribution of generalized Ramanujan primes

	Length of the longest run in [10 ⁵ , 10 ⁶] of					
	<i>c</i> -Ramanujan primes		Non-c-Ramanujan primes			
С	Expected	Actual	Expected	Actual		
0.50	14	20	16	36		

	Length of the longest run in [10 ⁵ , 10 ⁶] of					
	<i>c</i> -Ramanujan primes		Non-c-Ramanujan primes			
С	Expected	Actual	Expected	Actual		
0.10	70	58	5	3		
0.20	38	36	7	7		
0.30	25	25	10	12		
0.40	18	21	13	16		
0.50	14	20	16	36		
0.60	11	17	22	42		
0.70	9	14	30	78		
0.80	7	9	46	154		
0.90	5	11	91	345		

Open Problems

Introduction

1 Laishram and Sondow: $p_{2n} < R_n < p_{3n}$ for n > 1. Can we find good choices of a_c and b_c such that $p_{a_cn} \leq R_{c,n} \leq p_{b_cn}$ for all n?

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Results

2 For a given prime p, for what values of c is p a c-Ramanujan prime?

- **1** Laishram and Sondow: $p_{2n} < R_n < p_{3n}$ for n > 1. Can we find good choices of a_c and b_c such that $p_{a_0n} < R_{c,n} < p_{b_0n}$ for all n?
- 2 For a given prime p, for what values of c is p a c-Ramanujan prime?
- Is there any explanation for the unexpected distribution of c-Ramanujan primes amongst the primes?

Acknowledgements

Introduction

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We would like to thank Jonathan Sondow and our colleagues from the 2011 REU at Williams College.





