# HOMEWORK ASSIGNMENT # 1

### MATH 251, FALL 2006, WILLIAMS COLLEGE

### 1. Problem One

Construct the truth table of  $q \longleftrightarrow ((\neg p) \lor (\neg q))$ .

1.1. Solution. To fit the table on the page, we use x to denote  $(\neg p) \lor (\neg q)$ .

p	q	$\neg p$	$\neg q$	$(\neg p) \lor (\neg q) = x$	$q \longrightarrow x$	$x \longrightarrow q$	$q \longleftrightarrow x$
T	T	F	F	F	F	Т	F
T	F	F	T	T	T	F	F
F	T	T	F	T	T	T	T
F	F	T	T	T	T	F	F

## 2. Problem Two

Suppose that p and r are false statements and that q and s are true statements. Find the truth values of the following statements.

$$\begin{array}{ll} (1) & (p \longrightarrow q) \longrightarrow r \\ (2) & (s \longrightarrow (p \land (\neg r))) \land ((p \longrightarrow (r \lor q)) \land s) \end{array}$$

2.1. Solution. We construct just the relevant row of the truth table. For part (1):

We continue in part (2) with more of the row we need. Pay attention to the new symbols used to simplify things.

and we continue

$$\begin{array}{c|c} z \land s & (s \longrightarrow y) \land (z \land s) \\ \hline T & F \end{array}$$

So we see that both statements are **false**.

Date: September 19, 2006.

#### 3. PROBLEM THREE

Consider the statement  $\mathcal{A}$ : "If n is an integer, then  $\frac{n}{n+1}$  is not an integer."

- (1) Is  $\mathcal{A}$  true or false? Either prove it is true, or give a counterexample.
- (2) Write the converse, contrapositive and negation of  $\mathcal{A}$ .

3.1. Solution. In part one, the statement is false. My counterexample is n = 0. That is note the truth of the statements: "n = 0 is an integer" (so the hypothesis is true), and " $\frac{0}{1} = 0$  is an integer" (so the conclusion is false).

For part two: The converse of  $\mathcal{A}$  is "If  $\frac{n}{n+1}$  is not an integer, then n is an integer." The contrapositive is "If  $\frac{n}{n+1}$  is an integer, then n is not an integer." Finally, the negation is "n is an integer and  $\frac{n}{n+1}$  is an integer."

Notice that our counterexample is a number which makes the negation true. That is what a counterexample should be.

### 4. PROBLEM FOUR

Given the premises  $p \longrightarrow (\neg r)$  and  $r \lor q$ , write down a valid conclusion that involves p and q only and is not a tautology, or show that no such conclusion is possible.

4.1. Solution. I constructed a big truth table to figure out what was going on. Here it is, in all its glory.

p	q	r	$ \neg r$	$a = (p \longrightarrow (\neg r))$	$b = (r \lor q)$	$a \wedge b$	x	$a \wedge b \longrightarrow x$
T	T	T	F	F	T	F	1	Т
T	T	F	T	T	T	T	1	T
T	F	T	F	F	T	F	2	T
T	F	F	T	T	F	F	2	T
F	T	T	F	T	T	T	3	T
F	T	F	T	T	T	T	3	T
F	F	T	F	T	T	T	4	T
F	F	F	T	T	F	F	4	T

We need to find a statement x involving only p and q (but possibly being some horrible compound statement), which is not a tautology

and makes the statement

$$(p \longrightarrow (\neg r)) \land (r \lor q) \longrightarrow x$$

into a tautology. (Recall that this second requirement is the way to recognize a valid argument.)

Now, look at the table. First, notice that the truth value of x in the first two rows must agree. This is because these two rows are the same from the point of view of p and q. We label the value of these rows as 1.

The same can be said about the other pairs of rows, and we label their values as 2, 3 and 4.

If we consult the standard truth table for an implication, we see that there is no choice but to put T as the value of 1, 3 and 4. This is because the conjunction of the hypotheses is true in all of these rows.

In order to avoid x being a tautology, we have to have value 2 equal to F.

So, if we now focus just on the atomic statements p and q and the desired x we get a table that looks like the following.

$$\begin{array}{c|ccc} p & q & x \\ \hline T & T & T \\ T & F & F \\ F & T & T \\ F & F & T \end{array}$$

Fortunately, we can recognize this as the table of  $p \longrightarrow q$ . So the simplest answer is the statement  $p \longrightarrow q$ .

### 5. Problem Five

Decide if the following arguments are valid.

If I stay up late at night, then I will be tired in the morning.

(1) I stayed up late last night. I am tired this morning.

If I stay up late at night, then I will be tired in the morning.

(2) I am not tired this morning. I stayed up late last night.

If I stay up late at night, then I will be tired in the morning.

(3) I am not tired this morning. I did not stay up late last night.

If I stay up late at night, then I will be tired in the morning.

(4) I did not stay up late last night. I am not tired this morning.

If I like mathematics, then I will study.

- I will not study.
- (5) Either I like mathematics or I like football. I like football.

Either I study or I like football.

(6) If I like football, then I like mathematics. If I don't study, then I like mathematics.

5.1. Solution. Argument (2) is invalid: you might have gone to bed early just to be well rested. Argument (4) is invalid: A person can be tired for other reasons, like illness.

All the other arguments are valid. It is possible to test all of these using truth tables, but I can't bear to typeset another one after the last problem.

4